

RISC-V Debug Specification  
Version 1.0-STABLE  
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# Preface

**This stable specification might change before being accepted as standard. We're trying to only make backwards-compatible changes, but sometimes it's unavoidable to change a corner case. If you implement this document, please keep an eye on [the github repository](#) for changes that might affect your implementation.**

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# Chapter 1

## Introduction

When a design progresses from simulation to hardware implementation, a user's control and understanding of the system's current state drops dramatically. To help bring up and debug low level software and hardware, it is critical to have good debugging support built into the hardware. When a robust OS is running on a core, software can handle many debugging tasks. However, in many scenarios, hardware support is essential.

This document outlines a standard architecture for debug support on RISC-V hardware platforms. This architecture allows a variety of implementations and tradeoffs, which is complementary to the wide range of RISC-V implementations. At the same time, this specification defines common interfaces to allow debugging tools and components to target a variety of hardware platforms based on the RISC-V ISA.

System designers may choose to add additional hardware debug support, but this specification defines a standard interface for common functionality.

### 1.1 Terminology

#### **advanced feature**

An advanced feature for advanced users. Most users will not be able to take advantage of it.

#### **AMO**

Atomic Memory Operation.

#### **BYPASS**

JTAG instruction that selects a single bit data register, also called BYPASS.

#### **component**

A RISC-V core, or other part of a hardware platform. Typically all components will be connected to a single system bus.

#### **CSR**

Control and Status Register.

**DM** Debug Module (see Chapter [3](#)).

**DMI**

Debug Module Interface (see Section 3.1).

**DR** JTAG Data Register.

**DTM**

Debug Transport Module (see Section 6).

**DXLEN**

Debug XLEN, which is the widest XLEN a hart supports, ignoring the current value of MXL in `misa`.

**essential feature**

An essential feature must be present in order for debug to work correctly.

**GPR**

General Purpose Register.

**hardware platform**

A single system consisting of one or more *components*.

**hart**

A hardware thread in a RISC-V core.

**IDCODE**

32-bit Identification CODE, and a JTAG instruction that returns the IDCODE value.

**IR** JTAG Instruction Register.

**JTAG**

Refers to work done by IEEE's Joint Test Action Group, described in IEEE 1149.1.

**legacy feature**

A legacy feature should only be implemented to support legacy hardware that is present in a system.

**Minimal RISC-V Debug Specification**

A subset of the full Debug Specification that allows for very small implementations. See Chapter 3.

**NAPOT**

Naturally Aligned Power-Of-Two.

**NMI**

Non-Maskable Interrupt.

**physical address**

An address that is directly usable on the system bus.

**recommended feature**

A recommended feature is not required for debug to work correctly, but it is so useful that it should not be omitted without good reason.



**SBA**

System Bus Access (see Section 3.10).

**specialized feature**

A specialized feature, that only makes sense in the context of some specific hardware.

**TAP**

Test Access Port, defined in IEEE 1149.1.

**TM** Trigger Module (see Section 5).

**virtual address**

An address as a hart sees it. If the hart is using address translation this may be different from the physical address. If there is no translation then it will be the same.

**xepc**

The exception program counter CSR (e.g. `mepc`) that is appropriate for the mode being trapped to.

## 1.2 Context

This specification attempts to support all RISC-V ISA extensions that have, roughly, been ratified through the first half of 2023. In particular, though, this specification specifically addresses features in the following extensions:

1. A
2. C
3. D
4. F
5. H
6. Sm1p13
7. Ss1p13
8. Smstateen
9. V
10. Zawrs
11. Zcmp
12. Zicbom
13. Zicboz
14. Zicbop

### 1.2.1 Versions

Version 0.13 of this document was ratified by the RISC-V Foundation's board. Versions 0.13.x are bug fix releases to that ratified specification.

Version 0.14 was a working version that was never officially ratified.

Version 1.0 is almost entirely forwards and backwards compatible with Version 0.13.

### 1.2.1.1 Bugfixes from 0.13 to 1.0

Changes that fix a bug in the spec:

1. Fix order of operations described in `sbddata0`. #392
2. Resume ack is set after resume, in Section 3.5. #400
3. `sselect` applies to `svalue`. #402
4. `mte` only applies when `action=0`. #411
5. `aamsize` does not affect Argument Width. #420
6. Clarify that harts halt out of reset if `haltreq = 1`. #419

### 1.2.1.2 Incompatible Changes from 0.13 to 1.0

Changes that are not backwards-compatible. Debuggers or hardware implementations that implement 0.13 will have to change something in order to implement 1.0:

1. Make `haltsum0` optional if there is only one hart. #505
2. System bus autoincrement only happens if an access actually takes place. (`sbddata0`) #507
3. Bump `version` to 3. #512
4. Require debugger to poll `dmactive` after lowering it. #566
5. Add `pending` to `icount`. #574
6. When a selected trigger is disabled, `tdata2` and `tdata3` can be written with any value supported by any of the types this trigger supports. #721
7. `tcontrol` fields only apply to breakpoint traps, not any trap. #723
8. If `version` is greater than 0, then `hit0` (previously called `mcontrol6.hit`) now contains 0 when a trigger fires more than one instruction after the instruction that matched. (This information is now reflected in `hit1`.) #795
9. If `version` is greater than 0, then bit 20 of `mcontrol6` is no longer used for timing information. (Previously the bit was called `mcontrol6.timing`.) #807
10. If `version` is greater than 0, then the encodings of `size` for sizes greater than 64 bit have changed. #807

### 1.2.1.3 Minor Changes from 0.13 to 1.0

Changes that slightly modify defined behavior. Technically backwards incompatible, but unlikely to be noticeable:

1. `stopcount` only applies to hart-local counters. #405
2. `version` may be invalid when `dmactive = 0`. #414
3. Address triggers (`mcontrol1`) may fire on any accessed address. #421
4. All trigger registers (Section 5.3) are optional. #431

5. When extending IR, `bypass` still is all ones. #437
6. `ebreaks` and `ebreaku` are WARL. #458
7. NMIs are disabled by `stepie`. #465
8. R/W1C fields should be cleared by writing every bit high. #472
9. Specify trigger priorities in Table 5.2 relative to exceptions. #478
10. Time may pass before `dmactive` becomes high. #500
11. Clear MPRV when resuming into lower privilege mode. #503
12. Halt state may not be preserved across reset. #504
13. Hardware should clear trigger action when `dmode` is cleared and action is 1. #501
14. Change quick access exceptions to halt the target in Section 3.7.1.2. #585
15. Writing 0 to `tdata1` forces a state where `tdata2` and `tdata3` are writable. #598
16. Solutions to deal with reentrancy in Section 5.4 prevent triggers from *matching*, not merely *firing*. This primarily affects `icount` behavior. #722
17. Attempts to access an unimplemented CSR raise an illegal instruction exception. #791

#### 1.2.1.4 New Features from 0.13 to 1.0

New backwards-compatible feature that did not exist before:

1. Add halt groups and external triggers in Section 3.6. #404
2. Reserve some DMI space for non-standard use. See `custom`, and `custom0` through `custom15`. #406
3. Reserve trigger `type` values for non-standard use. #417
4. Add `nmi` bit to `itrigger`. #408 and #709
5. Recommend matching on every accessed address. #449
6. Add resume groups in Section 3.6. #506
7. Add `relaxedpriv`. #536
8. Move `scontext`, renaming original to `mscontext`, and create `hcontext`. #535
9. Add `mcontrol6`, deprecating `mcontrol`. #538
10. Add hypervisor support: `ebreakvs`, `ebreakvu`, `v`, `hcontext`, `mcontrol`, `mcontrol6`, and `priv`. #549
11. Optionally make `anyunavail` and `allunavail` sticky, controlled by `stickyunavail`. #520
12. Add `tmexttrigger` to support trigger module external trigger inputs. #543
13. Describe `mcontrol` and `mcontrol6` behavior with atomic instructions. #561
14. Trigger hit bits must be set on fire, may be set on match. #593
15. Add `sbytemask` and `sbytemask` to `textra32` and `textra64`. #588
16. Allow debugger to request harts stay alive with keepalive bit in Section 3.14.2. #592
17. Add `ndmresetpending` to allow a debugger to determine when `ndmreset` is complete. #594
18. Add `intctl` to support triggers from an interrupt controller. #599

#### 1.2.1.5 Incompatible Changes During 1.0 Stable

Backwards-incompatible changes between two versions that are both called 1.0 stable.

1. `nmi` was moved from `etrigger` to `itrigger`, and is now subject to the mode bits in that trigger.
2. [#728](#) introduced Message Registers, which were later removed in [#878](#).
3. It may not be possible to read the contents of the Program Buffer using the `progbuf` registers. [#731](#)
4. `tcontrol` fields apply to all traps, not just breakpoint traps. This reverts [#723](#). [#880](#)

## 1.3 About This Document

### 1.3.1 Structure

This document contains two parts. The main part of the document is the specification, which is given in the numbered chapters. The second part of the document is a set of appendices. The information in the appendices is intended to clarify and provide examples, but is not part of the actual specification.

### 1.3.2 ISA vs. non-ISA

This specification contains both ISA and non-ISA parts. The ISA parts define self-contained ISA extensions. The other parts of the document describe the non-ISA external debug extension. Chapters whose contents are solely one or the other are labeled as such in their title. Chapters without such a label apply to both ISA and non-ISA.

### 1.3.3 Register Definition Format

All register definitions in this document follow the format shown below. A simple graphic shows which fields are in the register. The upper and lower bit indices are shown to the top left and top right of each field. The total number of bits in the field are shown below it.

After the graphic follows a table which for each field lists its name, description, allowed accesses, and reset value. The allowed accesses are listed in [Table 1.2](#). The reset value is either a constant or “Preset.” The latter means it is an implementation-specific legal value.

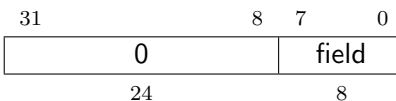
Parts of the register which are currently unused are labeled with the number 0. Software must only write 0 to those fields, and ignore their value while reading. Hardware must return 0 when those fields are read, and ignore the value written to them.

---

*This behavior enables us to use those fields later without having to increase the values in the version fields.*

Names of registers and their fields are hyperlinks to their definition, and are also listed in the index on page [117](#).

### 1.3.3.1 Long Name (shortname, at 0x123)



Field	Description	Access	Reset
field	Description of what this field is used for.	R/W	15

Table 1.2: Register Access Abbreviations

R	Read-only.
R/W	Read/Write.
R/W1C	Read/Write Ones to Clear. Writing 0 to every bit has no effect. Writing 1 to every bit clears the field. The result of other writes is undefined.
WARZ	Write any, read zero. A debugger may write any value. When read this field returns 0.
W1	Write-only. Only writing 1 has an effect. When read the returned value should be 0.
WARL	Write any, read legal. A debugger may write any value. If a value is unsupported, the implementation converts the value to one that is supported.

## 1.4 Background

There are several use cases for dedicated debugging hardware, both in native debug and external debug. Native debug (sometimes called self-hosted debug) refers to debug software running on a RISC-V platform which debugs the same platform. The optional Trigger Module provides features that are useful for native debug. External debug refers to debug software running somewhere else, debugging the RISC-V platform via a debug transport like JTAG. The entire document provides features that are useful for external debug.

This specification addresses the use cases listed below. Implementations can choose not to implement every feature, which means some use cases might not be supported.

- Accessing hardware on a hardware platform without a working CPU. (External debug.)
- Bootstrapping a hardware platform to test, configure, and program components before there is any executable code path in the hardware platform. (External debug.)
- Debugging low-level software in the absence of an OS or other software. (External debug.)
- Debugging issues in the OS itself. (External or native debug.)
- Debugging processes running on an OS. (Native or external debug.)

## 1.5 Supported Features

The debug interface described in this specification supports the following features:

1. All hart registers (including CSRs) can be read/written.
2. Memory can be accessed either from the hart's point of view, through the system bus directly, or both.
3. RV32, RV64, and future RV128 are all supported.
4. Any hart in the hardware platform can be independently debugged.
5. A debugger can discover almost<sup>1</sup> everything it needs to know itself, without user configuration.
6. Each hart can be debugged from the very first instruction executed.
7. A RISC-V hart can be halted when a software breakpoint instruction is executed.
8. Hardware single-step can execute one instruction at a time.
9. Debug functionality is independent of the debug transport used.
10. The debugger does not need to know anything about the microarchitecture of the harts it is debugging.
11. Arbitrary subsets of harts can be halted and resumed simultaneously. (Optional)
12. Arbitrary instructions can be executed on a halted hart. That means no new debug functionality is needed when a core has additional or custom instructions or state, as long as there exist programs that can move that state into GPRs. (Optional)
13. Registers can be accessed without halting. (Optional)
14. A running hart can be directed to execute a short sequence of instructions, with little overhead. (Optional)
15. A system bus manager allows memory access without involving any hart. (Optional)
16. A RISC-V hart can be halted when a trigger matches the PC, read/write address/data, or an instruction opcode. (Optional)
17. Harts can be grouped, and harts in the same group will all halt when any of them halts. These groups can also react to or notify external triggers. (Optional)

This document does not suggest a strategy or implementation for hardware test, debugging or error detection techniques. Scan, built-in self test (BIST), etc. are out of scope of this specification, but this specification does not intend to limit their use in RISC-V systems.

It is possible to debug code that uses software threads, but there is no special debug support for it.

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<sup>1</sup>Notable exceptions include information about the memory map and peripherals.

## Chapter 2

# System Overview

Figure 2.1 shows the main components of Debug Support. Blocks shown in dotted lines are optional.

The user interacts with the Debug Host (e.g. laptop), which is running a debugger (e.g. gdb). The debugger communicates with a Debug Translator (e.g. OpenOCD, which may include a hardware driver) to communicate with Debug Transport Hardware (e.g. Olimex USB-JTAG adapter). The Debug Transport Hardware connects the Debug Host to the hardware platform's Debug Transport Module (DTM). The DTM provides access to one or more Debug Modules (DMs) using the Debug Module Interface (DMI).

Each hart in the hardware platform is controlled by exactly one DM. Harts may be heterogeneous. There is no further limit on the hart-DM mapping, but usually all harts in a single core are controlled by the same DM. In most hardware platforms there will only be one DM that controls all the harts in the hardware platform.

DMs provide run control of their harts in the hardware platform. Abstract commands provide access to GPRs. Additional registers are accessible through abstract commands or by writing programs to the optional Program Buffer.

The Program Buffer allows the debugger to execute arbitrary instructions on a hart. This mechanism can also be used to access memory. An optional system bus access block allows memory accesses without using a RISC-V hart to perform the access.

Each RISC-V hart may implement a Trigger Module. When trigger conditions are met, harts will halt and inform the debug module that they have halted.

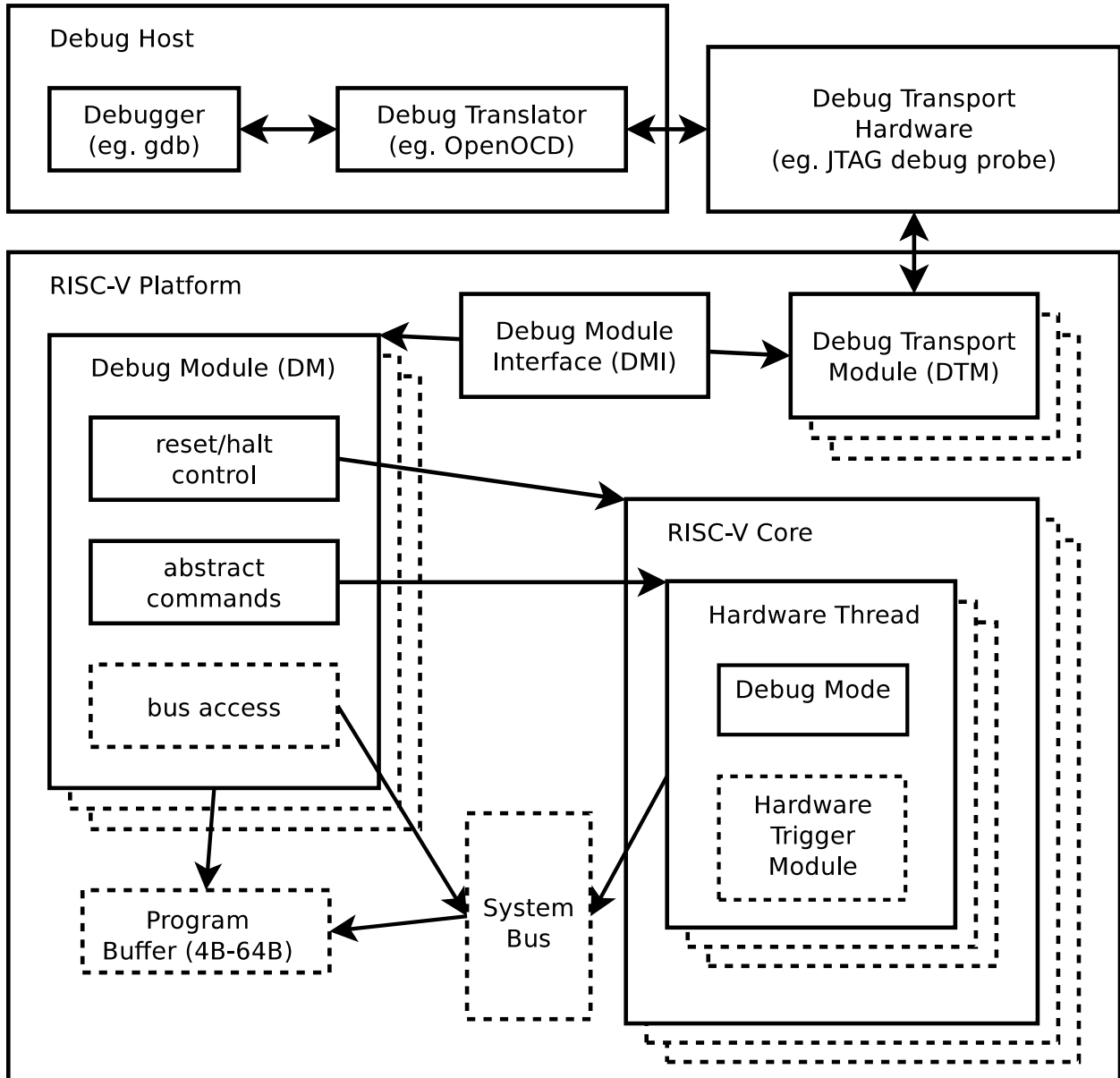


Figure 2.1: RISC-V Debug System Overview



## Chapter 3

# Debug Module (DM) (non-ISA extension)

The Debug Module implements a translation interface between abstract debug operations and their specific implementation. It might support the following operations:

1. Give the debugger necessary information about the implementation. (Required)
2. Allow any individual hart to be halted and resumed. (Required)
3. Provide status on which harts are halted. (Required)
4. Provide abstract read and write access to a halted hart's GPRs. (Required)
5. Provide access to a reset signal that allows debugging from the very first instruction after reset. (Required)
6. Provide a mechanism to allow debugging harts immediately out of reset (regardless of the reset cause). (Optional)
7. Provide abstract access to non-GPR hart registers. (Optional)
8. Provide a Program Buffer to force the hart to execute arbitrary instructions. (Optional)
9. Allow multiple harts to be halted, resumed, and/or reset at the same time. (Optional)
10. Allow memory access from a hart's point of view. (Optional)
11. Allow direct System Bus Access. (Optional)
12. Group harts. When any hart in the group halts, they all halt. (Optional)
13. Respond to external triggers by halting each hart in a configured group. (Optional)
14. Signal an external trigger when a hart in a group halts. (Optional)

In order to be compatible with this specification an implementation must:

1. Implement all the required features listed above.
2. Implement at least one of Program Buffer, System Bus Access, or Abstract Access Memory command mechanisms.
3. Do at least one of:
  - (a) Implement the Program Buffer.
  - (b) Implement abstract access to all registers that are visible to software running on the hart including all the registers that are present on the hart and listed in Table 3.3.

- (c) Implement abstract access to at least all GPRs, `dcsr`, and `dpc`, and advertise the implementation as conforming to the “Minimal RISC-V Debug Specification 1.0-STABLE”, instead of the “RISC-V Debug Specification 1.0-STABLE”.

A single DM can debug up to  $2^{20}$  harts.

### 3.1 Debug Module Interface (DMI)

Debug Modules are subordinates on a bus called the Debug Module Interface (DMI). The bus manager is the Debug Transport Module(s). The Debug Module Interface can be a trivial bus with one manager and one subordinate (see [A.3](#)), or use a more full-featured bus like TileLink or the AMBA Advanced Peripheral Bus. The details are left to the system designer.

The DMI uses between 7 and 32 address bits. Each address points at a single 32-bit register that can be read or written. The bottom of the address space is used for the first (and usually only) DM. Extra space can be used for custom debug devices, other cores, additional DMs, etc. If there are additional DMs on this DMI, the base address of the next DM in the DMI address space is given in `nextdm`.

The Debug Module is controlled via register accesses to its DMI address space.

### 3.2 Reset Control

There are two methods that allow a debugger to reset harts. `ndmreset` resets all the harts in the hardware platform, as well as all other parts of the hardware platform except for the Debug Modules, Debug Transport Modules, and Debug Module Interface. Exactly what is affected by this reset is implementation dependent, but it must be possible to debug programs from the first instruction executed. `hartreset` resets all the currently selected harts. In this case an implementation may reset more harts than just the ones that are selected. The debugger can discover which other harts are reset (if any) by selecting them and checking `anyhavereset` and `allhavereset`.

To perform either of these resets, the debugger first asserts the bit, and then clears it. The actual reset may start as soon as the bit is asserted, but may start an arbitrarily long time after the bit is deasserted. The reset itself may also take an arbitrarily long time. While the reset is on-going, harts are either in the running state, indicating it’s possible to perform some abstract commands during this time, or in the unavailable state, indicating it’s not possible to perform any abstract commands during this time. Once a hart’s reset is complete, `havereset` becomes set. When a hart comes out of reset and `haltreq` or `resethaltreq` are set, the hart will immediately enter Debug Mode (halted state). Otherwise, if the hart was initially running it will execute normally (running state) and if the hart was initially halted it should now be running but may be halted.

---

*There is no general, reliable way for the debugger to know when reset has actually begun.*

The Debug Module’s own state and registers should only be reset at power-up and while `dmactive` in `dmcontrol` is 0. If there is another mechanism to reset the DM, this mechanism must also reset

all the harts accessible to the DM.

Due to clock and power domain crossing issues, it might not be possible to perform arbitrary DMI accesses across hardware platform reset. While `ndmreset` or any external reset is asserted, the only supported DM operations are reading and writing `dmcontrol`. The behavior of other accesses is undefined.

When harts have been reset, they must set a sticky `havereset` state bit. The conceptual `havereset` state bits can be read for selected harts in `anyhavereset` and `allhavereset` in `dmstatus`. These bits must be set regardless of the cause of the reset. The `havereset` bits for the selected harts can be cleared by writing 1 to `ackhavereset` in `dmcontrol`. The `havereset` bits might or might not be cleared when `dmactive` is low.

### 3.3 Selecting Harts

Up to  $2^{20}$  harts can be connected to a single DM. Commands issued to the DM only apply to the currently selected harts.

To enumerate all the harts, a debugger must first determine `HARTSELLEN` by writing all ones to `hartsel` (assuming the maximum size) and reading back the value to see which bits were actually set. Then it selects each hart starting from 0 until either `anynonexistent` in `dmstatus` is 1, or the highest index (depending on `HARTSELLEN`) is reached.

The debugger can discover the mapping between hart indices and `mhartid` by using the interface to read `mhartid`, or by reading the hardware platform's configuration structure.

#### 3.3.1 Selecting a Single Hart

All debug modules must support selecting a single hart. The debugger can select a hart by writing its index to `hartsel`. Hart indexes start at 0 and are contiguous until the final index.

#### 3.3.2 Selecting Multiple Harts

Debug Modules may implement a Hart Array Mask register to allow selecting multiple harts at once. The  $n$ th bit in the Hart Array Mask register applies to the hart with index  $n$ . If the bit is 1 then the hart is selected. Usually a DM will have a Hart Array Mask register exactly wide enough to select all the harts it supports, but it's allowed to tie any of these bits to 0.

The debugger can set bits in the hart array mask register using `hawindowssel` and `hawindow`, then apply actions to all selected harts by setting `hasel`. If this feature is supported, multiple harts can be halted, resumed, and reset simultaneously. The state of the hart array mask register is not affected by setting or clearing `hasel`.

Execution of Abstract Commands ignores this mechanism and only applies to the hart selected by `hartsel`.

### 3.4 Hart DM States

Every hart that can be selected is in exactly one of the following four DM states: non-existent, unavailable, running, or halted. Which state the selected harts are in is reflected by [allnonexistent](#), [anyononexistent](#), [allunavail](#), [anyunavail](#), [allrunning](#), [anyrunning](#), [allhalted](#), and [anyhalted](#).

Harts are nonexistent if they will never be part of this hardware platform, no matter how long a user waits. E.g. in a simple single-hart hardware platform only one hart exists, and all others are nonexistent. Debuggers may assume that a hardware platform has no harts with indexes higher than the first nonexistent one.

Harts are unavailable if they might exist/become available at a later time, or if there are other harts with higher indexes than this one. Harts may be unavailable for a variety of reasons including being reset, temporarily powered down, and not being plugged into the hardware platform. That means harts might become available or unavailable at any time, although these events should be rare in hardware platforms built to be easily debugged. There are no guarantees about the state of the hart when it becomes available.

Hardware platforms with very large number of harts may permanently disable some during manufacturing, leaving holes in the otherwise continuous hart index space. In order to let the debugger discover all harts, they must show up as unavailable even if there is no chance of them ever becoming available.

Harts are running when they are executing normally, as if no debugger was attached. This includes being in a low power mode or waiting for an interrupt, as long as a halt request will result in the hart being halted.

Harts are halted when they are in Debug Mode, only performing tasks on behalf of the debugger.

Which states a hart that is reset goes through is implementation dependent. Harts may be unavailable while reset is asserted, and some time after reset is deasserted. They might transition to running for some time after reset is deasserted. Finally they end up either running or halted, depending on [haltreq](#) and [resethaltreq](#).

### 3.5 Run Control

For every hart, the Debug Module tracks 4 conceptual bits of state: halt request, resume ack, halt-on-reset request, and hart reset. (The hart reset and halt-on-reset request bits are optional.) These 4 bits reset to 0, except for resume ack, which may reset to either 0 or 1. The DM receives halted, running, and havereset signals from each hart. The debugger can observe the state of resume ack in [allresumeack](#) and [anyresumeack](#), and the state of halted, running, and havereset signals in [allhalted](#), [anyhalted](#), [allrunning](#), [anyrunning](#), [allhavereset](#), and [anyhavereset](#). The state of the other bits cannot be observed directly.

When a debugger writes 1 to [haltreq](#), each selected hart's halt request bit is set. When a running hart, or a hart just coming out of reset, sees its halt request bit high, it responds by halting, deasserting its running signal, and asserting its halted signal. Halted harts ignore their halt request

bit.

When a debugger writes 1 to `resumereq`, each selected hart's resume ack bit is cleared and each selected, halted hart is sent a resume request. Harts respond by resuming, clearing their halted signal, and asserting their running signal. At the end of this process the resume ack bit is set. These status signals of all selected harts are reflected in `allresumeack`, `anyresumeack`, `allrunning`, and `anyrunning`. Resume requests are ignored by running harts.

When halt or resume is requested, a hart must respond in less than one second, unless it is unavailable. (How this is implemented is not further specified. A few clock cycles will be a more typical latency).

The DM can implement optional halt-on-reset bits for each hart, which it indicates by setting `hasresethaltreq` to 1. This means the DM implements the `setresethaltreq` and `clrresethaltreq` bits. Writing 1 to `setresethaltreq` sets the halt-on-reset request bit for each selected hart. When a hart's halt-on-reset request bit is set, the hart will immediately enter debug mode on the next deassertion of its reset. This is true regardless of the reset's cause. The hart's halt-on-reset request bit remains set until cleared by the debugger writing 1 to `clrresethaltreq` while the hart is selected, or by DM reset.

If the DM is reset while a hart is halted, it is UNSPECIFIED whether that hart resumes. Debuggers should use `resumereq` to explicitly resume harts before clearing `dmactive` and disconnecting.

### 3.6 Halt Groups, Resume Groups, and External Triggers

An optional feature allows a debugger to place harts into two kinds of groups: halt groups and resume groups. It is also possible to add external triggers to a halt and resume groups. At any given time, each hart and each trigger is a member of exactly one halt group and exactly one resume group.

In both halt and resume groups, group 0 is special. Harts in group 0 halt/resume as if groups aren't implemented at all.

When any hart in a halt group halts:

1. That hart halts normally, with `cause` reflecting the original cause of the halt.
2. All the other harts in the halt group that are running will quickly halt. `cause` for those harts should be set to 6, but may be set to 3. Other harts in the halt group that are halted but have started the process of resuming must also quickly become halted, even if they do resume briefly.
3. Any external triggers in that group are notified.

Adding a hart to a halt group does not automatically halt that hart, even if other harts in the group are already halted.

When an external trigger that's a member of the halt group fires:

1. All the harts in the halt group that are running will quickly halt. [cause](#) for those harts should be set to 6, but may be set to 3. Other harts in the halt group that are halted but have started the process of resuming must also quickly become halted, even if they do resume briefly.

When any hart in a resume group resumes:

1. All the other harts in that group that are halted will quickly resume as soon as any currently executing abstract commands have completed. Each hart in the group sets its resume ack bit as soon as it has resumed. Harts that are in the process of halting should complete that process and stay halted.
2. Any external triggers in that group are notified.

Adding a hart to a resume group does not automatically resume that hart, even if other harts in the group are currently running.

When an external trigger that's a member of the resume group fires:

1. All the harts in that group that are halted will quickly resume as soon as any currently executing abstract commands have completed. Each hart in the group sets its resume ack bit as soon as it has resumed. Harts that are in the process of halting should complete that process and stay halted.

External triggers are abstract concepts that can signal the DM and/or receive signals from the DM. This configuration is done through [dmcs2](#), where external triggers are referred to by a number. Commonly, external triggers are capable of sending a signal from the hardware platform into the DM, as well as receiving a signal from the DM to take their own action on. It is also allowable for an external trigger to be input-only or output-only. By convention external triggers 0–7 are bidirectional, triggers 8–11 are input-only, and triggers 12–15 are output-only but this is not required.

---

*External triggers could be used to implement near simultaneous halting/resuming of all cores in a hardware platform, when not all cores are RISC-V cores.*

When the DM is reset, all harts must be placed in the lowest-numbered halt and resume groups that they can be in. (This will usually be group 0.)

Some designs may choose to hardcode hart groups to a group other than group 0, meaning it is never possible to halt or resume just a single hart. This is explicitly allowed. In that case it must be possible to discover the groups by using [dmcs2](#) even if it's not possible to change the configuration.

### 3.7 Abstract Commands

The DM supports a set of abstract commands, most of which are optional. Depending on the implementation, the debugger may be able to perform some abstract commands even when the selected hart is not halted. Debuggers can only determine which abstract commands are supported by a given hart in a given state (running, halted, or held in reset) by attempting them and then looking at [cmderr](#) in [abstractcs](#) to see if they were successful. Commands may be supported with

some options set, but not with other options set. If a command has unsupported options set or if bits that are defined as 0 aren't 0, then the DM must set `cmderr` to 2 (not supported).

---

*Example: Every DM must support the Access Register command, but might not support accessing CSRs. If the debugger requests to read a CSR in that case, the command will return “not supported.”*

Debuggers execute abstract commands by writing them to `command`. They can determine whether an abstract command is complete by reading `busy` in `abstractcs`. If the debugger starts a new command while `busy` is set, `cmderr` becomes 1 (busy), the currently executing command still gets to run to completion, but any error generated by the currently executing command is lost. After completion, `cmderr` indicates whether the command was successful or not. Commands may fail because a hart is not halted, not running, unavailable, or because they encounter an error during execution.

If the command takes arguments, the debugger must write them to the `data` registers before writing to `command`. If a command returns results, the Debug Module must ensure they are placed in the `data` registers before `busy` is cleared. Which `data` registers are used for the arguments is described in Table 3.1. In all cases the least-significant word is placed in the lowest-numbered `data` register. The argument width depends on the command being executed, and is `DXLEN` where not explicitly specified.

Table 3.1: Use of Data Registers

Argument Width	arg0/return value	arg1	arg2
32	<code>data0</code>	<code>data1</code>	<code>data2</code>
64	<code>data0</code> , <code>data1</code>	<code>data2</code> , <code>data3</code>	<code>data4</code> , <code>data5</code>
128	<code>data0</code> – <code>data3</code>	<code>data4</code> – <code>data7</code>	<code>data8</code> – <code>data11</code>

---

*The Abstract Command interface is designed to allow a debugger to write commands as fast as possible, and then later check whether they completed without error. In the common case the debugger will be much slower than the target and commands succeed, which allows for maximum throughput. If there is a failure, the interface ensures that no commands execute after the failing one. To discover which command failed, the debugger has to look at the state of the DM (e.g. contents of `data0`) or hart (e.g. contents of a register modified by a Program Buffer program) to determine which one failed.*

Before starting an abstract command, a debugger must ensure that `haltreq`, `resumereq`, and `ackhavereset` are all 0.

While an abstract command is executing (`busy` in `abstractcs` is high), a debugger must not change `hartsel`, and must not write 1 to `haltreq`, `resumereq`, `ackhavereset`, `setresethaltreq`, or `clrresethaltreq`.

If an abstract command does not complete in the expected time and appears to be hung, the debugger can try to reset the hart (using `hartreset` or `ndmreset`). If that doesn't clear `busy`, then it can try resetting the Debug Module (using `dmactive`).

If an abstract command is started while the selected hart is unavailable or if a hart becomes unavailable while executing an abstract command, then the Debug Module may terminate the abstract command, setting `busy` low, and `cmderr` to 4 (halt/resume). Alternatively, the command could just appear to be hung (`busy` never goes low).

### 3.7.1 Abstract Command Listing

This section describes each of the different abstract commands and how their fields should be interpreted when they are written to `command`.

Each abstract command is a 32-bit value. The top 8 bits contain `cmdtype` which determines the kind of command. Table 3.2 lists all commands.

Table 3.2: Meaning of `cmdtype`

<code>cmdtype</code>	Command	Page
0	Access Register Command	18
1	Quick Access	20
2	Access Memory Command	20

#### 3.7.1.1 Access Register

This command gives the debugger access to CPU registers and allows it to execute the Program Buffer. It performs the following sequence of operations:

1. If `write` is clear and `transfer` is set, then copy data from the register specified by `regno` into the `arg0` region of `data`, and perform any side effects that occur when this register is read from M-mode.
2. If `write` is set and `transfer` is set, then copy data from the `arg0` region of `data` into the register specified by `regno`, and perform any side effects that occur when this register is written from M-mode.
3. If `aarpostincrement` and `transfer` are set, increment `regno`. `regno` may also be incremented if `aarpostincrement` is set and `transfer` is clear.
4. Execute the Program Buffer, if `postexec` is set.

If any of these operations fail, `cmderr` is set and none of the remaining steps are executed. An implementation may detect an upcoming failure early, and fail the overall command before it reaches the step that would cause failure. If the failure is that the requested register does not exist in the hart, `cmderr` must be set to 3 (exception).

Debug Modules must implement this command and must support read and write access to all GPRs when the selected hart is halted. Debug Modules may optionally support accessing other registers, or accessing registers when the hart is running. It is recommended that if one register in a group is accessible, then all registers in that group are accessible, but each individual register (aside from GPRs) may be supported differently across read, write, and halt status.

Registers might not be accessible if they wouldn't be accessible by M mode code currently running. (E.g. `fflags` might not be accessible when `mstatus.FS` is 0.) If this is the case, the debugger is responsible for changing state to make the registers accessible. The Core Debug Registers (Section 4.10) should be accessible if abstract CSR access is implemented.

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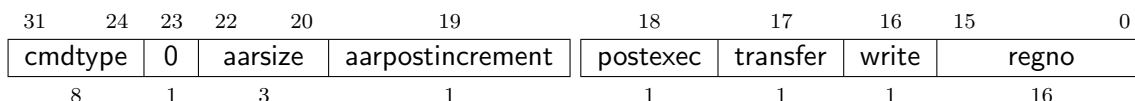
*The encoding of `aarsize` was chosen to match `sbaccess` in `sbc`.*



Table 3.3: Abstract Register Numbers

Numbers	Group Description
0x0000 – 0x0fff	CSRs. The “PC” can be accessed here through <a href="#">dpc</a> .
0x1000 – 0x101f	GPRs
0x1020 – 0x103f	Floating point registers
0xc000 – 0xffff	Reserved for non-standard extensions and internal use.

This command modifies `arg0` only when a register is read. The other `data` registers are not changed.



Field	Description
<code>cmdtype</code>	This is 0 to indicate Access Register Command.
<code>aarsize</code>	2 (32bit): Access the lowest 32 bits of the register. 3 (64bit): Access the lowest 64 bits of the register. 4 (128bit): Access the lowest 128 bits of the register. If <code>aarsize</code> specifies a size larger than the register’s actual size, then the access must fail. If a register is accessible, then reads of <code>aarsize</code> less than or equal to the register’s actual size must be supported. Writing less than the full register may be supported, but what happens to the high bits in that case is UNSPECIFIED. This field controls the Argument Width as referenced in Table 3.1.
<code>aarpostincrement</code>	0 (disabled): No effect. This variant must be supported. 1 (enabled): After a successful register access, <code>regno</code> is incremented. Incrementing past the highest supported value causes <code>regno</code> to become UNSPECIFIED. Supporting this variant is optional. It is undefined whether the increment happens when <code>transfer</code> is 0.
<code>postexec</code>	0 (disabled): No effect. This variant must be supported, and is the only supported one if <code>progbuFSIZE</code> is 0. 1 (enabled): Execute the program in the Program Buffer exactly once after performing the transfer, if any. Supporting this variant is optional.

*Continued on next page*

Field	Description
transfer	0 (disabled): Don't do the operation specified by <a href="#">write</a> . 1 (enabled): Do the operation specified by <a href="#">write</a> . This bit can be used to just execute the Program Buffer without having to worry about placing valid values into <a href="#">aarsize</a> or <a href="#">regno</a> .
write	When <a href="#">transfer</a> is set: 0 (arg0): Copy data from the specified register into <a href="#">arg0</a> portion of <a href="#">data</a> . 1 (register): Copy data from <a href="#">arg0</a> portion of <a href="#">data</a> into the specified register.
regno	Number of the register to access, as described in <a href="#">Table 3.3</a> . <a href="#">dpc</a> may be used as an alias for PC if this command is supported on a non-halted hart.

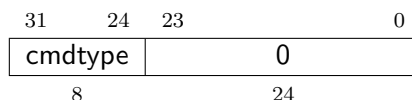
### 3.7.1.2 Quick Access

Perform the following sequence of operations:

1. If the hart is halted, the command sets [cmderr](#) to “halt/resume” and does not continue.
2. Halt the hart. If the hart halts for some other reason (e.g. breakpoint), the command sets [cmderr](#) to “halt/resume” and does not continue.
3. Execute the Program Buffer. If an exception occurs, [cmderr](#) is set to “exception,” the Program Buffer execution ends, and the hart is halted with [cause](#) set to 3.
4. If the Program Buffer executed without an exception, then resume the hart.

Implementing this command is optional.

This command does not touch the [data](#) registers.



Field	Description
cmdtype	This is 1 to indicate Quick Access command.

### 3.7.1.3 Access Memory

This command lets the debugger perform memory accesses, with the exact same memory view and permissions as the selected hart has. This includes access to hart-local memory-mapped registers, etc. The command performs the following sequence of operations:

1. Copy data from the memory location specified in `arg1` into the `arg0` portion of `data`, if `write` is clear.
2. Copy data from the `arg0` portion of `data` into the memory location specified in `arg1`, if `write` is set.
3. If `aampostincrement` is set, increment `arg1`.

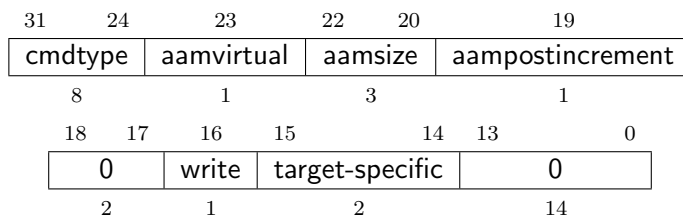
If any of these operations fail, `cmderr` is set and none of the remaining steps are executed. An access may only fail if the hart, running M-mode code, might encounter that same failure when it attempts the same access. An implementation may detect an upcoming failure early, and fail the overall command before it reaches the step that would cause failure.

Debug Modules may optionally implement this command and may support read and write access to memory locations when the selected hart is running or halted. If this command supports memory accesses while the hart is running, it must also support memory accesses while the hart is halted.

---

*The encoding of `aamsize` was chosen to match `sbaccess` in `sbc`s.*

This command modifies `arg0` only when memory is read. It modifies `arg1` only if `aampostincrement` is set. The other `data` registers are not changed.



Field	Description
<code>cmdtype</code>	This is 2 to indicate Access Memory Command.
<code>aamvirtual</code>	<p>An implementation does not have to implement both virtual and physical accesses, but it must fail accesses that it doesn't support.</p> <p>0 (physical): Addresses are physical (to the hart they are performed on).</p> <p>1 (virtual): Addresses are virtual, and translated the way they would be from M-mode, with MPRV set.</p> <p>Debug Modules on systems without address translation (i.e. virtual addresses equal physical) may optionally allow <code>aamvirtual</code> set to 1, which would produce the same result as that same abstract command with <code>aamvirtual</code> cleared.</p>

*Continued on next page*

Field	Description
aamsize	0 (8bit): Access the lowest 8 bits of the memory location. 1 (16bit): Access the lowest 16 bits of the memory location. 2 (32bit): Access the lowest 32 bits of the memory location. 3 (64bit): Access the lowest 64 bits of the memory location. 4 (128bit): Access the lowest 128 bits of the memory location.
aampostincrement	After a memory access has completed, if this bit is 1, increment <code>arg1</code> (which contains the address used) by the number of bytes encoded in <code>aamsize</code> . Supporting this variant is optional, but highly recommended for performance reasons.
write	0 ( <code>arg0</code> ): Copy data from the memory location specified in <code>arg1</code> into the low bits of <code>arg0</code> . The value of the remaining bits of <code>arg0</code> are UNSPECIFIED. 1 (memory): Copy data from the low bits of <code>arg0</code> into the memory location specified in <code>arg1</code> .
target-specific	These bits are reserved for target-specific uses.

### 3.8 Program Buffer

To support executing arbitrary instructions on a halted hart, a Debug Module can include a Program Buffer that a debugger can write small programs to. DMs that support all necessary functionality using abstract commands only may choose to omit the Program Buffer.

A debugger can write a small program to the Program Buffer, and then execute it exactly once with the Access Register Abstract Command, setting the `postexec` bit in `command`. The debugger can write whatever program it likes (including jumps out of the Program Buffer), but the program must end with `ebreak` or `c.ebreak`. An implementation may support an implicit `ebreak` that is executed when a hart runs off the end of the Program Buffer. This is indicated by `impebreak`. With this feature, a Program Buffer of just 2 32-bit words can offer efficient debugging.

While these programs are executed, the hart does not leave Debug Mode (see Section 4.1). If an exception is encountered during execution of the Program Buffer, no more instructions are executed, the hart remains in Debug Mode, and `cmderr` is set to 3 (exception error). If the debugger executes a program that doesn't terminate with an `ebreak` instruction, the hart will remain in Debug Mode and the debugger will lose control of the hart.

If `progbufsize` is 1 then the following apply:

1. `impebreak` must be 1.

2. If the debugger writes a compressed instruction into the Program Buffer, it must be placed into the lower 16 bits and accompanied by a compressed `nop` in the upper 16 bits.

---

*This requirement on the debugger for the case of `progbufsize` equal to 1 is to accommodate hardware designs that prefer to stuff instructions directly into the pipeline when halted, instead of having the Program Buffer exist in the address space somewhere.*

The Program Buffer may be implemented as RAM which is accessible to the hart. A debugger can determine if this is the case by executing small programs that attempt to write and read back relative to `pc` while executing from the Program Buffer. If so, the debugger has more flexibility in what it can do with the program buffer.

### 3.9 Overview of Hart Debug States

Figure 3.1 shows a conceptual view of the states passed through by a hart during run/halt debugging as influenced by the different fields of `dmcontrol`, `abstractcs`, `abstractauto`, and `command`.

### 3.10 System Bus Access

A debugger can access memory from a hart's point of view using a Program Buffer or the Abstract Access Memory command. (Both these features are optional.) A Debug Module may also include a System Bus Access block to provide memory access without involving a hart, regardless of whether Program Buffer is implemented. The System Bus Access block uses physical addresses.

The System Bus Access block may support 8-, 16-, 32-, 64-, and 128-bit accesses. Table 3.7 shows which bits in `sbddata` are used for each access size.

Table 3.7: System Bus Data Bits

Access Size	Data Bits
8	<code>sbddata0</code> bits 7:0
16	<code>sbddata0</code> bits 15:0
32	<code>sbddata0</code>
64	<code>sbddata1</code> , <code>sbddata0</code>
128	<code>sbddata3</code> , <code>sbddata2</code> , <code>sbddata1</code> , <code>sbddata0</code>

Depending on the microarchitecture, data accessed through System Bus Access might not always be coherent with that observed by each hart. It is up to the debugger to enforce coherency if the implementation does not. This specification does not define a standard way to do this. Possibilities may include writing to special memory-mapped locations, or executing special instructions via the Program Buffer.

---

*Implementing a System Bus Access block has several benefits even when a Debug Module also implements a Program Buffer. First, it is possible to access memory in a running system with minimal impact. Second, it may improve performance when accessing memory. Third, it may provide access to devices that a hart does not have access to.*

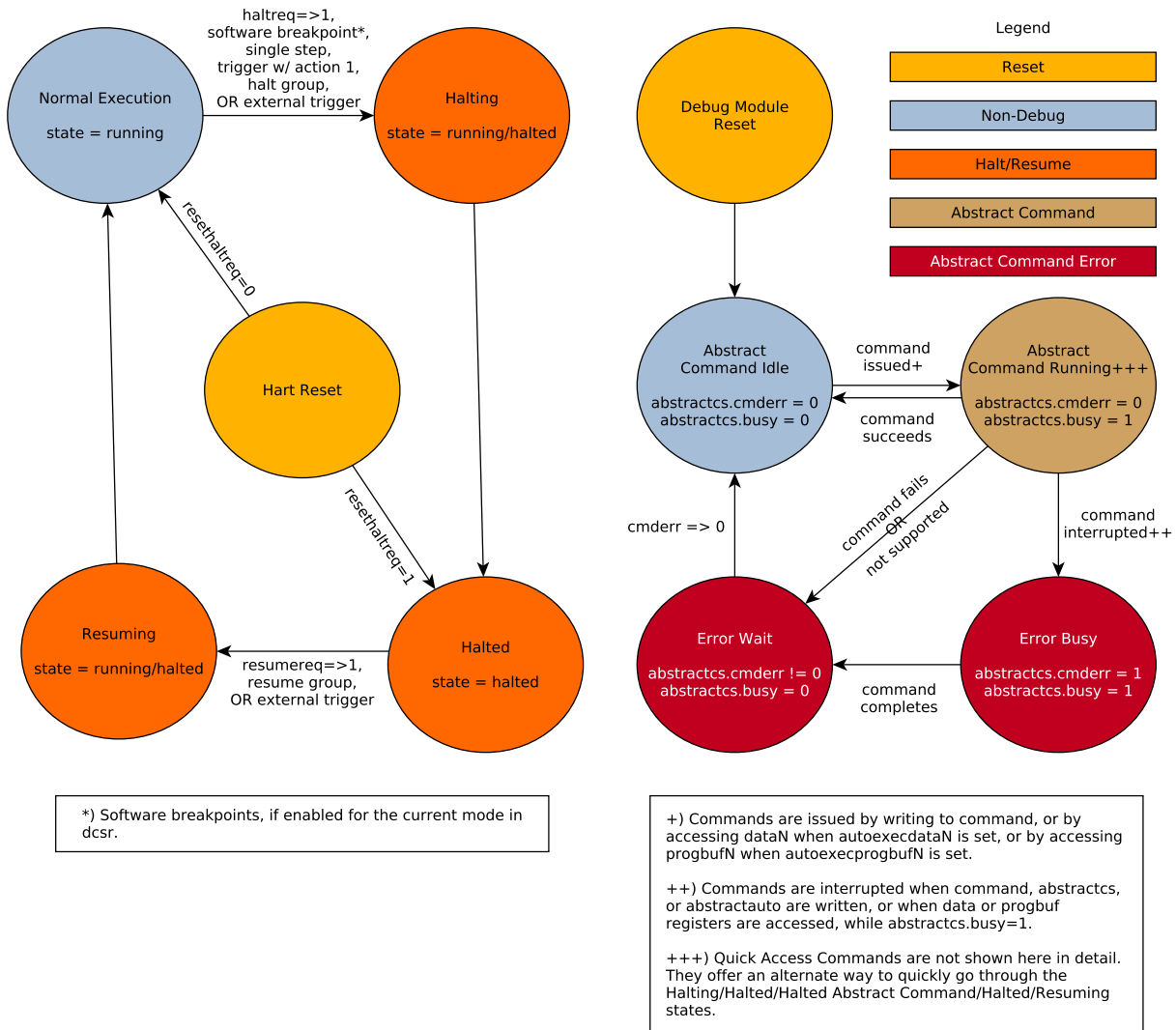


Figure 3.1: Run/Halt Debug State Machine for single-hart hardware platforms. As only a small amount of state is visible to the debugger, the states and transitions are conceptual.

### 3.11 Minimally Intrusive Debugging

Depending on the task it is performing, some harts can only be halted very briefly. There are several mechanisms that allow accessing resources in such a running system with a minimal impact on the running hart.

First, an implementation may allow some abstract commands to execute without halting the hart.

Second, the Quick Access abstract command can be used to halt a hart, quickly execute the contents of the Program Buffer, and let the hart run again. Combined with instructions that allow Program Buffer code to access the `data` registers, as described in [hartinfo](#), this can be used to quickly perform a memory or register access. For some hardware platforms this will be too intrusive, but many hardware platforms that can't be halted can bear an occasional hiccup of a hundred or less cycles.

Third, if the System Bus Access block is implemented, it can be used while a hart is running to access system memory.

### 3.12 Security

To protect intellectual property it may be desirable to lock access to the Debug Module. To allow access during a manufacturing process and not afterwards, a reasonable solution could be to add a fuse bit to the Debug Module that can be used to be permanently disable it. Since this is technology specific, it is not further addressed in this spec.

Another option is to allow the DM to be unlocked only by users who have an access key. Between [authenticated](#), [authbusy](#), and [authdata](#) arbitrarily complex authentication mechanism can be supported. When [authenticated](#) is clear, the DM must not interact with the rest of the hardware platform, nor expose details about the harts connected to the DM. All DM registers should read 0, while writes should be ignored, with the following mandatory exceptions:

1. [authenticated](#) in [dmstatus](#) is readable.
2. [authbusy](#) in [dmstatus](#) is readable.
3. [version](#) in [dmstatus](#) is readable.
4. [dmactive](#) in [dmcontrol](#) is readable and writable.
5. [authdata](#) is readable and writable.

Implementations where it's not possible to unlock the DM by using [authdata](#) should not implement that register.

### 3.13 Version Detection

To detect the version of the Debug Module with a minimum of side effects, use the following procedure:

1. Read `dmcontrol`.
2. If `dmactive` is 0 or `ndmreset` is 1:
  - (a) Write `dmcontrol`, preserving `hartreset`, `hasel`, `hartsello`, and `hartselhi` from the value that was read, setting `dmactive`, and clearing all the other bits.
  - (b) Read `dmcontrol` until `dmactive` is high.
3. Read `dmstatus`, which contains `version`.

If it was necessary to clear `ndmreset`, this might have the following unavoidable side effects:

1. `haltreq` is cleared, potentially preventing a halt request made by a previous debugger from taking effect.
2. `resumereq` is cleared, potentially preventing a resume request made by a previous debugger from taking effect.
3. `ndmreset` is deasserted, releasing the hardware platform from reset if a previous debugger had set it.
4. `dmactive` is asserted, releasing the DM from reset. This in itself is not observable by any harts.

This procedure is guaranteed to work in future versions of this spec. The meaning of the `dmcontrol` bits where `hartreset`, `hasel`, `hartsello`, and `hartselhi` currently reside might change, but preserving them will have no side effects. Clearing the bits of `dmcontrol` not explicitly mentioned here will have no side effects beyond the ones mentioned above.

### 3.14 Debug Module Registers

The registers described in this section are accessed over the DMI bus. Each DM has a base address (which is 0 for the first DM). The register addresses below are offsets from this base address.

Debug Module DMI Registers that are unimplemented or not mentioned in the table below return 0 when read. Writing them has no effect.

Table 3.8: Debug Module Debug Bus Registers

Address	Name	Page
0x04	Abstract Data 0 ( <code>data0</code> )	39
0x05	Abstract Data 1 ( <code>data1</code> )	
0x06	Abstract Data 2 ( <code>data2</code> )	
0x07	Abstract Data 3 ( <code>data3</code> )	
0x08	Abstract Data 4 ( <code>data4</code> )	
0x09	Abstract Data 5 ( <code>data5</code> )	
0x0a	Abstract Data 6 ( <code>data6</code> )	
0x0b	Abstract Data 7 ( <code>data7</code> )	
0x0c	Abstract Data 8 ( <code>data8</code> )	
0x0d	Abstract Data 9 ( <code>data9</code> )	
0x0e	Abstract Data 10 ( <code>data10</code> )	

*Continued on next page*



Table 3.8: Debug Module Debug Bus Registers

Address	Name	Page
0x0f	Abstract Data 11 ( <code>data11</code> )	
0x10	Debug Module Control ( <code>dmcontrol</code> )	30
0x11	Debug Module Status ( <code>dmstatus</code> )	28
0x12	Hart Info ( <code>hartinfo</code> )	33
0x13	Halt Summary 1 ( <code>haltsum1</code> )	42
0x14	Hart Array Window Select ( <code>hawindowselect</code> )	34
0x15	Hart Array Window ( <code>hawindow</code> )	35
0x16	Abstract Control and Status ( <code>abstractcs</code> )	35
0x17	Abstract Command ( <code>command</code> )	36
0x18	Abstract Command Autoexec ( <code>abstractauto</code> )	37
0x19	Configuration Structure Pointer 0 ( <code>confstrptr0</code> )	38
0x1a	Configuration Structure Pointer 1 ( <code>confstrptr1</code> )	38
0x1b	Configuration Structure Pointer 2 ( <code>confstrptr2</code> )	38
0x1c	Configuration Structure Pointer 3 ( <code>confstrptr3</code> )	39
0x1d	Next Debug Module ( <code>nextdm</code> )	39
0x1f	Custom Features ( <code>custom</code> )	48
0x20	Program Buffer 0 ( <code>progbuf0</code> )	40
0x21	Program Buffer 1 ( <code>progbuf1</code> )	
0x22	Program Buffer 2 ( <code>progbuf2</code> )	
0x23	Program Buffer 3 ( <code>progbuf3</code> )	
0x24	Program Buffer 4 ( <code>progbuf4</code> )	
0x25	Program Buffer 5 ( <code>progbuf5</code> )	
0x26	Program Buffer 6 ( <code>progbuf6</code> )	
0x27	Program Buffer 7 ( <code>progbuf7</code> )	
0x28	Program Buffer 8 ( <code>progbuf8</code> )	
0x29	Program Buffer 9 ( <code>progbuf9</code> )	
0x2a	Program Buffer 10 ( <code>progbuf10</code> )	
0x2b	Program Buffer 11 ( <code>progbuf11</code> )	
0x2c	Program Buffer 12 ( <code>progbuf12</code> )	
0x2d	Program Buffer 13 ( <code>progbuf13</code> )	
0x2e	Program Buffer 14 ( <code>progbuf14</code> )	
0x2f	Program Buffer 15 ( <code>progbuf15</code> )	
0x30	Authentication Data ( <code>authdata</code> )	40
0x32	Debug Module Control and Status 2 ( <code>dmcs2</code> )	40
0x34	Halt Summary 2 ( <code>haltsum2</code> )	42
0x35	Halt Summary 3 ( <code>haltsum3</code> )	43
0x37	System Bus Address 127:96 ( <code>sbaddress3</code> )	46
0x38	System Bus Access Control and Status ( <code>sbcsc</code> )	43
0x39	System Bus Address 31:0 ( <code>sbaddress0</code> )	45
0x3a	System Bus Address 63:32 ( <code>sbaddress1</code> )	45
0x3b	System Bus Address 95:64 ( <code>sbaddress2</code> )	46
0x3c	System Bus Data 31:0 ( <code>sbdata0</code> )	46
0x3d	System Bus Data 63:32 ( <code>sbdata1</code> )	47
0x3e	System Bus Data 95:64 ( <code>sbdata2</code> )	48

*Continued on next page*

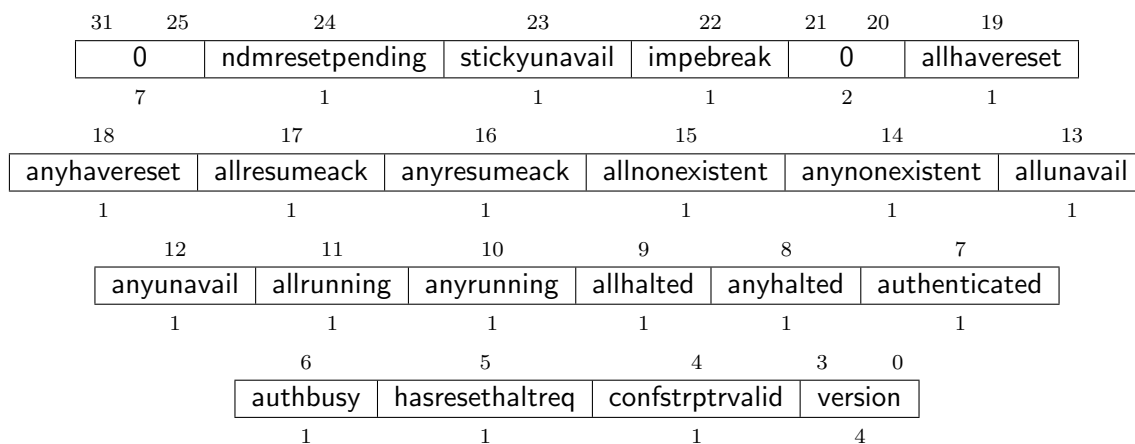
Table 3.8: Debug Module Debug Bus Registers

Address	Name	Page
0x3f	System Bus Data 127:96 ( <i>sbdata3</i> )	48
0x40	Halt Summary 0 ( <i>haltsum0</i> )	41
0x70	Custom Features 0 ( <i>custom0</i> )	48
0x71	Custom Features 1 ( <i>custom1</i> )	
0x72	Custom Features 2 ( <i>custom2</i> )	
0x73	Custom Features 3 ( <i>custom3</i> )	
0x74	Custom Features 4 ( <i>custom4</i> )	
0x75	Custom Features 5 ( <i>custom5</i> )	
0x76	Custom Features 6 ( <i>custom6</i> )	
0x77	Custom Features 7 ( <i>custom7</i> )	
0x78	Custom Features 8 ( <i>custom8</i> )	
0x79	Custom Features 9 ( <i>custom9</i> )	
0x7a	Custom Features 10 ( <i>custom10</i> )	
0x7b	Custom Features 11 ( <i>custom11</i> )	
0x7c	Custom Features 12 ( <i>custom12</i> )	
0x7d	Custom Features 13 ( <i>custom13</i> )	
0x7e	Custom Features 14 ( <i>custom14</i> )	
0x7f	Custom Features 15 ( <i>custom15</i> )	

### 3.14.1 Debug Module Status (*dmstatus*, at 0x11)

This register reports status for the overall Debug Module as well as the currently selected harts, as defined in [hasel](#). Its address will not change in the future, because it contains [version](#).

This entire register is read-only.



Field	Description	Access	Reset
ndmresetpending	0 (false): Unimplemented, or <code>ndmreset</code> is zero and no <code>ndmreset</code> is currently in progress. 1 (true): <code>ndmreset</code> is currently nonzero, or there is an <code>ndmreset</code> in progress.	R	-
stickyunavail	0 (current): The per-hart <code>unavail</code> bits reflect the current state of the hart. 1 (sticky): The per-hart <code>unavail</code> bits are sticky. Once they are set, they will not clear until the debugger acknowledges them using <code>ackunavail</code> .	R	Preset
impebreak	If 1, then there is an implicit <code>ebreak</code> instruction at the non-existent word immediately after the Program Buffer. This saves the debugger from having to write the <code>ebreak</code> itself, and allows the Program Buffer to be one word smaller. This must be 1 when <code>progbufsize</code> is 1.	R	Preset
allhavereset	This field is 1 when all currently selected harts have been reset and reset has not been acknowledged for any of them.	R	-
anyhavereset	This field is 1 when at least one currently selected hart has been reset and reset has not been acknowledged for that hart.	R	-
allresumeack	This field is 1 when all currently selected harts have their resume ack bit set.	R	-
anyresumeack	This field is 1 when any currently selected hart has its resume ack bit set.	R	-
allnonexistent	This field is 1 when all currently selected harts do not exist in this hardware platform.	R	-
anynonexistent	This field is 1 when any currently selected hart does not exist in this hardware platform.	R	-
allunavail	This field is 1 when all currently selected harts are unavailable, or (if <code>stickyunavail</code> is 1) were unavailable without that being acknowledged.	R	-
anyunavail	This field is 1 when any currently selected hart is unavailable, or (if <code>stickyunavail</code> is 1) was unavailable without that being acknowledged.	R	-
allrunning	This field is 1 when all currently selected harts are running.	R	-
anyrunning	This field is 1 when any currently selected hart is running.	R	-
allhalted	This field is 1 when all currently selected harts are halted.	R	-
anyhalted	This field is 1 when any currently selected hart is halted.	R	-

*Continued on next page*

Field	Description	Access	Reset
authenticated	0 (false): Authentication is required before using the DM. 1 (true): The authentication check has passed. On components that don't implement authentication, this bit must be preset as 1.	R	Preset
authbusy	0 (ready): The authentication module is ready to process the next read/write to <code>authdata</code> . 1 (busy): The authentication module is busy. Accessing <code>authdata</code> results in unspecified behavior. <code>authbusy</code> only becomes set in immediate response to an access to <code>authdata</code> .	R	0
hasresethaltreq	1 if this Debug Module supports halt-on-reset functionality controllable by the <code>setresethaltreq</code> and <code>clrresethaltreq</code> bits. 0 otherwise.	R	Preset
confstrpvalid	0 (invalid): <code>confstrptr0–confstrptr3</code> hold information which is not relevant to the configuration structure. 1 (valid): <code>confstrptr0–confstrptr3</code> hold the address of the configuration structure.	R	Preset
version	0 (none): There is no Debug Module present. 1 (0.11): There is a Debug Module and it conforms to version 0.11 of this specification. 2 (0.13): There is a Debug Module and it conforms to version 0.13 of this specification. 3 (1.0): There is a Debug Module and it conforms to version 1.0 of this specification. 15 (custom): There is a Debug Module but it does not conform to any available version of this spec.	R	3

### 3.14.2 Debug Module Control (`dmcontrol`, at `0x10`)

This register controls the overall Debug Module as well as the currently selected harts, as defined in [hasel](#).

Throughout this document we refer to `hartsel`, which is `hartselhi` combined with `hartsello`. While the spec allows for 20 `hartsel` bits, an implementation may choose to implement fewer than that. The actual width of `hartsel` is called `HARTSELLEN`. It must be at least 0 and at most 20. A debugger should discover `HARTSELLEN` by writing all ones to `hartsel` (assuming the maximum size) and reading back the value to see which bits were actually set. Debuggers must not change `hartsel` while an abstract command is executing.

---

*There are separate `setresethaltreq` and `clrresethaltreq` bits so that it is possible to write `dmcontrol` without changing the halt-on-reset request bit for each selected hart, when not all selected harts have the same configuration.*

On any given write, a debugger may only write 1 to at most one of the following bits: `resumereq`, `hartreset`, `ackhavereset`, `setresethaltreq`, and `clrresethaltreq`. The others must be written 0.

`resethaltreq` is an optional internal bit of per-hart state that cannot be read, but can be written with `setresethaltreq` and `clrresethaltreq`.

`keepalive` is an optional internal bit of per-hart state. When it is set, it suggests that the hardware should attempt to keep the hart available for the debugger, e.g. by keeping it from entering a low-power state once powered on. Even if the bit is implemented, hardware might not be able to keep a hart available. The bit is written through `setkeepalive` and `clrkeepalive`.

For forward compatibility, `version` will always be readable when bit 1 (`ndmreset`) is 0 and bit 0 (`dmactive`) is 1.

31	30	29	28	27					
haltreq	resumereq	hartreset	ackhavereset	ackunavail					
1	1	1	1	1					
26	25	16	15	6	5				
hasel	hartsello	hartselhi	setkeepalive	clrkeepalive					
1	10	10	1	1					
3	2		1	0					
setresethaltreq		clrresethaltreq		ndmreset	dmactive				
1	1		1	1					

Field	Description	Access	Reset
<code>haltreq</code>	Writing 0 clears the halt request bit for all currently selected harts. This may cancel outstanding halt requests for those harts. Writing 1 sets the halt request bit for all currently selected harts. Running harts will halt whenever their halt request bit is set. Writes apply to the new value of <code>hartsel</code> and <code>hasel</code> .	WARZ	-
<code>resumereq</code>	Writing 1 causes the currently selected harts to resume once, if they are halted when the write occurs. It also clears the resume ack bit for those harts. <code>resumereq</code> is ignored if <code>haltreq</code> is set. Writes apply to the new value of <code>hartsel</code> and <code>hasel</code> .	W1	-
<code>hartreset</code>	This optional field writes the reset bit for all the currently selected harts. To perform a reset the debugger writes 1, and then writes 0 to deassert the reset signal. While this bit is 1, the debugger must not change which harts are selected. If this feature is not implemented, the bit always stays 0, so after writing 1 the debugger can read the register back to see if the feature is supported. Writes apply to the new value of <code>hartsel</code> and <code>hasel</code> .	WARL	0

*Continued on next page*

Field	Description	Access	Reset
ackhavereset	0 (nop): No effect. 1 (ack): Clears <b>havereset</b> for any selected harts. Writes apply to the new value of <b>hartsel</b> and <b>hasel</b> .	W1	-
ackunavail	0 (nop): No effect. 1 (ack): Clears <b>unavail</b> for any selected harts that are currently available. Writes apply to the new value of <b>hartsel</b> and <b>hasel</b> .	W1	-
hasel	Selects the definition of currently selected harts. 0 (single): There is a single currently selected hart, that is selected by <b>hartsel</b> . 1 (multiple): There may be multiple currently selected harts – the hart selected by <b>hartsel</b> , plus those selected by the hart array mask register. An implementation which does not implement the hart array mask register must tie this field to 0. A debugger which wishes to use the hart array mask register feature should set this bit and read back to see if the functionality is supported.	WARL	0
hartsello	The low 10 bits of <b>hartsel</b> : the DM-specific index of the hart to select. This hart is always part of the currently selected harts.	WARL	0
hartselhi	The high 10 bits of <b>hartsel</b> : the DM-specific index of the hart to select. This hart is always part of the currently selected harts.	WARL	0
setkeepalive	This optional field sets <b>keepalive</b> for all currently selected harts, unless <b>clrkeepalive</b> is simultaneously set to 1. Writes apply to the new value of <b>hartsel</b> and <b>hasel</b> .	W1	-
clrkeepalive	This optional field clears <b>keepalive</b> for all currently selected harts. Writes apply to the new value of <b>hartsel</b> and <b>hasel</b> .	W1	-
setresethaltreq	This optional field writes the halt-on-reset request bit for all currently selected harts, unless <b>clrresethaltreq</b> is simultaneously set to 1. When set to 1, each selected hart will halt upon the next deassertion of its reset. The halt-on-reset request bit is not automatically cleared. The debugger must write to <b>clrresethaltreq</b> to clear it. Writes apply to the new value of <b>hartsel</b> and <b>hasel</b> . If <b>hasresethaltreq</b> is 0, this field is not implemented.	W1	-
clrresethaltreq	This optional field clears the halt-on-reset request bit for all currently selected harts. Writes apply to the new value of <b>hartsel</b> and <b>hasel</b> .	W1	-

*Continued on next page*

Field	Description	Access	Reset
ndmreset	This bit controls the reset signal from the DM to the rest of the hardware platform. The signal should reset every part of the hardware platform, including every hart, except for the DM and any logic required to access the DM. To perform a hardware platform reset the debugger writes 1, and then writes 0 to deassert the reset.	R/W	0
dmactive	This bit serves as a reset signal for the Debug Module itself. After changing the value of this bit, the debugger must poll <code>dmcontrol</code> until <code>dmactive</code> has taken the requested value before performing any action that assumes the requested <code>dmactive</code> state change has completed. Hardware may take an arbitrarily long time to complete activation or deactivation and will indicate completion by setting <code>dmactive</code> to the requested value. 0 (inactive): The module's state, including authentication mechanism, takes its reset values (the <code>dmactive</code> bit is the only bit which can be written to something other than its reset value). Any accesses to the module may fail. Specifically, <code>version</code> might not return correct data. 1 (active): The module functions normally. No other mechanism should exist that may result in resetting the Debug Module after power up. To place the Debug Module into a known state, a debugger may write 0 to <code>dmactive</code> , poll until <code>dmactive</code> is observed 0, write 1 to <code>dmactive</code> , and poll until <code>dmactive</code> is observed 1. Implementations may pay attention to this bit to further aid debugging, for example by preventing the Debug Module from being power gated while debugging is active.	R/W	0

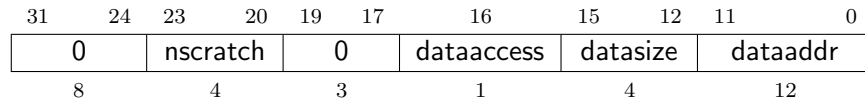
### 3.14.3 Hart Info (`hartinfo`, at `0x12`)

This register gives information about the hart currently selected by `hartsel`.

This register is optional. If it is not present it should read all-zero.

If this register is included, the debugger can do more with the Program Buffer by writing programs which explicitly access the `data` and/or `dscratch` registers.

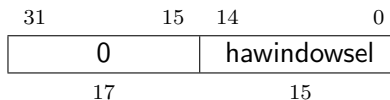
This entire register is read-only.



Field	Description	Access	Reset
nscratch	Number of <code>dscratch</code> registers available for the debugger to use during program buffer execution, starting from <code>dscratch0</code> . The debugger can make no assumptions about the contents of these registers between commands.	R	Preset
dataaccess	0 (csr): The <code>data</code> registers are shadowed in the hart by CSRs. Each CSR is <code>DXLEN</code> bits in size, and corresponds to a single argument, per Table 3.1. 1 (memory): The <code>data</code> registers are shadowed in the hart's memory map. Each register takes up 4 bytes in the memory map.	R	Preset
datasize	If <code>dataaccess</code> is 0: Number of CSRs dedicated to shadowing the <code>data</code> registers. If <code>dataaccess</code> is 1: Number of 32-bit words in the memory map dedicated to shadowing the <code>data</code> registers. Since there are at most 12 <code>data</code> registers, the value in this register must be 12 or smaller.	R	Preset
dataaddr	If <code>dataaccess</code> is 0: The number of the first CSR dedicated to shadowing the <code>data</code> registers. If <code>dataaccess</code> is 1: Address of RAM where the <code>data</code> registers are shadowed. This address is sign extended giving a range of -2048 to 2047, easily addressed with a load or store using <code>x0</code> as the address register.	R	Preset

### 3.14.4 Hart Array Window Select (`hawindowse1`, at `0x14`)

This register selects which of the 32-bit portion of the hart array mask register (see Section 3.3.2) is accessible in `hawindow`.



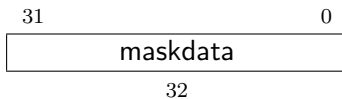


Field	Description	Access	Reset
hawindowssel	The high bits of this field may be tied to 0, depending on how large the array mask register is. E.g. on a hardware platform with 48 harts only bit 0 of this field may actually be writable.	WARL	0

### 3.14.5 Hart Array Window (hawindow, at 0x15)

This register provides R/W access to a 32-bit portion of the hart array mask register (see Section 3.3.2). The position of the window is determined by `hawindowssel`. I.e. bit 0 refers to hart `hawindowssel * 32`, while bit 31 refers to hart `hawindowssel * 32 + 31`.

Since some bits in the hart array mask register may be constant 0, some bits in this register may be constant 0, depending on the current value of `hawindowssel`.

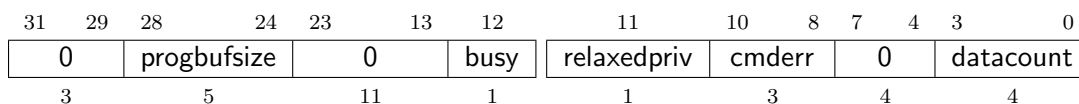


### 3.14.6 Abstract Control and Status (abstractcs, at 0x16)

Writing this register while an abstract command is executing causes `cmderr` to become 1 (busy) once the command completes (busy becomes 0).

---

`datacount` must be at least 1 to support RV32 harts, 2 to support RV64 harts, or 4 to support RV128 harts.



Field	Description	Access	Reset
progbuFSIZE	Size of the Program Buffer, in 32-bit words. Valid sizes are 0 - 16.	R	Preset
busy	0 (ready): There is no abstract command currently being executed. 1 (busy): An abstract command is currently being executed. This bit is set as soon as <code>command</code> is written, and is not cleared until that command has completed.	R	0

*Continued on next page*

Field	Description	Access	Reset
<code>relaxedpriv</code>	<p>This optional bit controls whether program buffer and abstract memory accesses are performed with the exact and full set of permission checks that apply based on the current architectural state of the hart performing the access, or with a relaxed set of permission checks (e.g. PMP restrictions are ignored). The details of the latter are implementation-specific.</p> <p>0 (full checks): Full permission checks apply.</p> <p>1 (relaxed checks): Relaxed permission checks apply.</p>	WARL	Preset
<code>cmderr</code>	<p>Gets set if an abstract command fails. The bits in this field remain set until they are cleared by writing 1 to them. No abstract command is started until the value is reset to 0.</p> <p>This field only contains a valid value if <code>busy</code> is 0.</p> <p>0 (none): No error.</p> <p>1 (busy): An abstract command was executing while <code>command</code>, <code>abstractcs</code>, or <code>abstractauto</code> was written, or when one of the <code>data</code> or <code>progbuf</code> registers was read or written. This status is only written if <code>cmderr</code> contains 0.</p> <p>2 (not supported): The command in <code>command</code> is not supported. It may be supported with different options set, but it will not be supported at a later time when the hart or system state are different.</p> <p>3 (exception): An exception occurred while executing the command (e.g. while executing the Program Buffer).</p> <p>4 (halt/resume): The abstract command couldn't execute because the hart wasn't in the required state (running/halted), or unavailable.</p> <p>5 (bus): The abstract command failed due to a bus error (e.g. alignment, access size, or timeout).</p> <p>6 (reserved): Reserved for future use.</p> <p>7 (other): The command failed for another reason.</p>	R/W1C	0
<code>datacount</code>	<p>Number of <code>data</code> registers that are implemented as part of the abstract command interface. Valid sizes are 1 – 12.</p>	R	Preset

### 3.14.7 Abstract Command (`command`, at `0x17`)

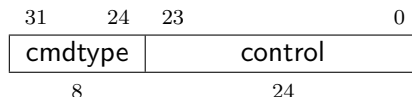
Writes to this register cause the corresponding abstract command to be executed.

Writing this register while an abstract command is executing causes `cmderr` to become 1 (busy) once the command completes (busy becomes 0).

If `cmderr` is non-zero, writes to this register are ignored.

---

`cmderr` inhibits starting a new command to accommodate debuggers that, for performance reasons, send several commands to be executed in a row without checking `cmderr` in between. They can safely do so and check `cmderr` at the end without worrying that one command failed but then a later command (which might have depended on the previous one succeeding) passed.



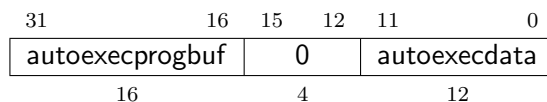
Field	Description	Access	Reset
cmdtype	The type determines the overall functionality of this abstract command.	WARZ	0
control	This field is interpreted in a command-specific manner, described for each abstract command.	WARZ	0

### 3.14.8 Abstract Command Autoexec (abstractauto, at 0x18)

This register is optional. Including it allows more efficient burst accesses. A debugger can detect whether it is supported by setting bits and reading them back.

If this register is implemented then bits corresponding to implemented `progbuf` and data registers must be writable. Other bits must be hard-wired to 0.

If this register is written while an abstract command is executing then the write is ignored and `cmderr` becomes 1 (busy) once the command completes (busy becomes 0).



Field	Description	Access	Reset
autoexecprogbuf	When a bit in this field is 1, read or write accesses to the corresponding <code>progbuf</code> word cause the DM to act as if the current value in <code>command</code> was written there again after the access to <code>progbuf</code> completes.	WARL	0

*Continued on next page*

Field	Description	Access	Reset
autoexecdata	When a bit in this field is 1, read or write accesses to the corresponding <code>data</code> word cause the DM to act as if the current value in <code>command</code> was written there again after the access to <code>data</code> completes.	WARL	0

### 3.14.9 Configuration Structure Pointer 0 (`confstrptr0`, at `0x19`)

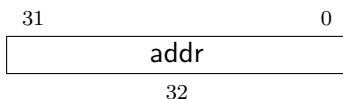
When `confstrptrvalid` is set, reading this register returns bits 31:0 of the configuration structure pointer. Reading the other `confstrptr` registers returns the upper bits of the address.

When system bus access is implemented, this must be an address that can be used with the System Bus Access module. Otherwise, this must be an address that can be used to access the configuration structure from the hart with ID 0.

If `confstrptrvalid` is 0, then the `confstrptr` registers hold identifier information which is not further specified in this document.

The configuration structure itself is a data structure of the same format as the data structure pointed to by `mconfigptr` as described in the Privileged Spec.

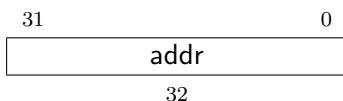
This entire register is read-only.



### 3.14.10 Configuration Structure Pointer 1 (`confstrptr1`, at `0x1a`)

When `confstrptrvalid` is set, reading this register returns bits 63:32 of the configuration structure pointer. See `confstrptr0` for more details.

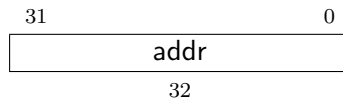
This entire register is read-only.



### 3.14.11 Configuration Structure Pointer 2 (`confstrptr2`, at `0x1b`)

When `confstrptrvalid` is set, reading this register returns bits 95:64 of the configuration structure pointer. See `confstrptr0` for more details.

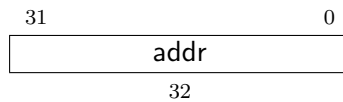
This entire register is read-only.



### 3.14.12 Configuration Structure Pointer 3 (confstrptr3, at 0x1c)

When [confstrptrvalid](#) is set, reading this register returns bits 127:96 of the configuration structure pointer. See [confstrptr0](#) for more details.

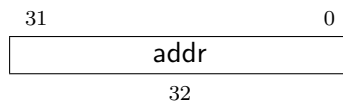
This entire register is read-only.



### 3.14.13 Next Debug Module (nextdm, at 0x1d)

If there is more than one DM accessible on this DMI, this register contains the base address of the next one in the chain, or 0 if this is the last one in the chain.

This entire register is read-only.



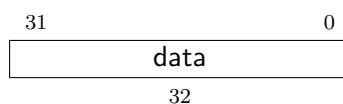
### 3.14.14 Abstract Data 0 (data0, at 0x04)

[data0](#) through [data11](#) are registers that may be read or changed by abstract commands. [datacount](#) indicates how many of them are implemented, starting at [data0](#), counting up. Table 3.1 shows how abstract commands use these registers.

Accessing these registers while an abstract command is executing causes [cmderr](#) to be set to 1 (busy) if it is 0.

Attempts to write them while [busy](#) is set does not change their value.

The values in these registers might not be preserved after an abstract command is executed. The only guarantees on their contents are the ones offered by the command in question. If the command fails, no assumptions can be made about the contents of these registers.



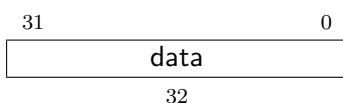
### 3.14.15 Program Buffer 0 (progbuf0, at 0x20)

`progbuf0` through `progbuf15` must provide write access to the optional program buffer. It may also be possible for the debugger to read from the program buffer through these registers. If reading is not supported, then all reads return 0.

`progbufsize` indicates how many `progbuf` registers are implemented starting at `progbuf0`, counting up.

Accessing these registers while an abstract command is executing causes `cmderr` to be set to 1 (busy) if it is 0.

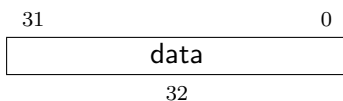
Attempts to write them while `busy` is set does not change their value.



### 3.14.16 Authentication Data (authdata, at 0x30)

This register serves as a 32-bit serial port to/from the authentication module.

When `authbusy` is clear, the debugger can communicate with the authentication module by reading or writing this register. There is no separate mechanism to signal overflow/underflow.



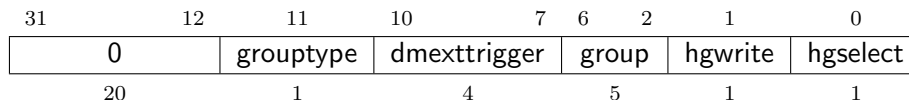
### 3.14.17 Debug Module Control and Status 2 (dmcs2, at 0x32)

This register contains DM control and status bits that didn't easily fit in `dmcontrol` and `dmstatus`. All are optional.

If halt groups are not implemented, then `group` will always be 0 when `grouptype` is 0.

If resume groups are not implemented, then `grouptype` will remain 0 even after 1 is written there.

The DM external triggers available to add to halt groups may be the same as or distinct from the DM external triggers available to add to resume groups.



Field	Description	Access	Reset
grouptype	0 (halt): The remaining fields in this register configure halt groups. 1 (resume): The remaining fields in this register configure resume groups.	WARL	0
dmexttrigger	This field contains the currently selected DM external trigger. If a non-existent trigger value is written here, the hardware will change it to a valid one or 0 if no DM external triggers exist.	WARL	0
group	When <a href="#">hgselect</a> is 0, contains the group of the hart specified by <a href="#">hartsel</a> . When <a href="#">hgselect</a> is 1, contains the group of the DM external trigger selected by <a href="#">dmexttrigger</a> . The value written to this field is ignored unless <a href="#">hgwrite</a> is also written 1. Group numbers are contiguous starting at 0, with the highest number being implementation-dependent, and possibly different between different group types. Debuggers should read back this field after writing to confirm they are using a hart group that is supported. If groups aren't implemented, then this entire field is 0.	WARL	preset
hgwrite	When 1 is written and <a href="#">hgselect</a> is 0, for every selected hart the DM will change its group to the value written to <a href="#">group</a> , if the hardware supports that group for that hart. Implementations may also change the group of a minimal set of unselected harts in the same way, if that is necessary due to a hardware limitation. When 1 is written and <a href="#">hgselect</a> is 1, the DM will change the group of the DM external trigger selected by <a href="#">dmexttrigger</a> to the value written to <a href="#">group</a> , if the hardware supports that group for that trigger. Writing 0 has no effect.	W1	-
hgselect	0 (harts): Operate on harts. 1 (triggers): Operate on DM external triggers. If there are no DM external triggers, this field must be tied to 0.	WARL	0

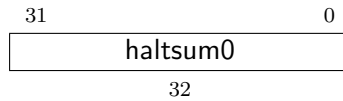
### 3.14.18 Halt Summary 0 (haltsum0, at 0x40)

Each bit in this read-only register indicates whether one specific hart is halted or not. Unavailable/nonexistent harts are not considered to be halted.

This register might not be present if fewer than 2 harts are connected to this DM.

The LSB reflects the halt status of hart  $\{\text{hartsel}[19:5], 5'h0\}$ , and the MSB reflects halt status of hart  $\{\text{hartsel}[19:5], 5'h1f\}$ .

This entire register is read-only.



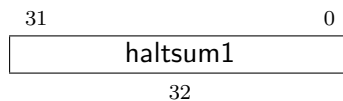
### 3.14.19 Halt Summary 1 (haltsum1, at 0x13)

Each bit in this read-only register indicates whether any of a group of harts is halted or not. Unavailable/nonexistent harts are not considered to be halted.

This register might not be present if fewer than 33 harts are connected to this DM.

The LSB reflects the halt status of harts  $\{\text{hartsel}[19:10], 10'h0\}$  through  $\{\text{hartsel}[19:10], 10'h1f\}$ . The MSB reflects the halt status of harts  $\{\text{hartsel}[19:10], 10'h3e0\}$  through  $\{\text{hartsel}[19:10], 10'h3ff\}$ .

This entire register is read-only.



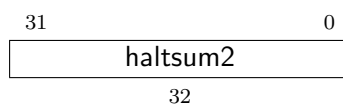
### 3.14.20 Halt Summary 2 (haltsum2, at 0x34)

Each bit in this read-only register indicates whether any of a group of harts is halted or not. Unavailable/nonexistent harts are not considered to be halted.

This register might not be present if fewer than 1025 harts are connected to this DM.

The LSB reflects the halt status of harts  $\{\text{hartsel}[19:15], 15'h0\}$  through  $\{\text{hartsel}[19:15], 15'h3ff\}$ . The MSB reflects the halt status of harts  $\{\text{hartsel}[19:15], 15'h7c00\}$  through  $\{\text{hartsel}[19:15], 15'h7fff\}$ .

This entire register is read-only.







Field	Description	Access	Reset
sbbusy	When 1, indicates the system bus manager is busy. (Whether the system bus itself is busy is related, but not the same thing.) This bit goes high immediately when a read or write is requested for any reason, and does not go low until the access is fully completed. Writes to <code>sbcs</code> while <code>sbbusy</code> is high result in undefined behavior. A debugger must not write to <code>sbcs</code> until it reads <code>sbbusy</code> as 0.	R	0
sbreadonaddr	When 1, every write to <code>sbaddress0</code> automatically triggers a system bus read at the new address.	R/W	0
sbaccess	Select the access size to use for system bus accesses. 0 (8bit): 8-bit 1 (16bit): 16-bit 2 (32bit): 32-bit 3 (64bit): 64-bit 4 (128bit): 128-bit If <code>sbaccess</code> has an unsupported value when the DM starts a bus access, the access is not performed and <code>sberror</code> is set to 4.	R/W	2
sbautoincrement	When 1, <code>sbaddress</code> is incremented by the access size (in bytes) selected in <code>sbaccess</code> after every system bus access.	R/W	0
sbreadondata	When 1, every read from <code>sbddata0</code> automatically triggers a system bus read at the (possibly auto-incremented) address.	R/W	0
sberror	When the Debug Module's system bus manager encounters an error, this field gets set. The bits in this field remain set until they are cleared by writing 1 to them. While this field is non-zero, no more system bus accesses can be initiated by the Debug Module. An implementation may report "Other" (7) for any error condition. 0 (none): There was no bus error. 1 (timeout): There was a timeout. 2 (address): A bad address was accessed. 3 (alignment): There was an alignment error. 4 (size): An access of unsupported size was requested. 7 (other): Other.	R/W1C	0
sbasize	Width of system bus addresses in bits. (0 indicates there is no bus access support.)	R	Preset
sbaccess128	1 when 128-bit system bus accesses are supported.	R	Preset

*Continued on next page*

Field	Description	Access	Reset
sbaccess64	1 when 64-bit system bus accesses are supported.	R	Preset
sbaccess32	1 when 32-bit system bus accesses are supported.	R	Preset
sbaccess16	1 when 16-bit system bus accesses are supported.	R	Preset
sbaccess8	1 when 8-bit system bus accesses are supported.	R	Preset

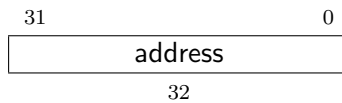
### 3.14.23 System Bus Address 31:0 (sbaddress0, at 0x39)

If `sbasize` is 0, then this register is not present.

When the system bus manager is busy, writes to this register will set `sdbusyerror` and don't do anything else.

If `sberror` is 0, `sdbusyerror` is 0, and `sbreadonaddr` is set then writes to this register start the following:

1. Set `sdbusy`.
2. Perform a bus read from the new value of `sbaddress`.
3. If the read succeeded and `sbautoincrement` is set, increment `sbaddress`.
4. Clear `sdbusy`.

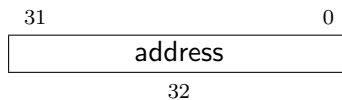


Field	Description	Access	Reset
address	Accesses bits 31:0 of the physical address in <code>sbaddress</code> .	R/W	0

### 3.14.24 System Bus Address 63:32 (sbaddress1, at 0x3a)

If `sbasize` is less than 33, then this register is not present.

When the system bus manager is busy, writes to this register will set `sdbusyerror` and don't do anything else.

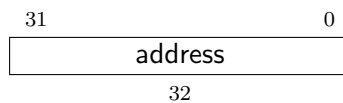


Field	Description	Access	Reset
address	Accesses bits 63:32 of the physical address in <b>sbaddress</b> (if the system address bus is that wide).	R/W	0

### 3.14.25 System Bus Address 95:64 (sbaddress2, at 0x3b)

If **sbasize** is less than 65, then this register is not present.

When the system bus manager is busy, writes to this register will set **sbbusyerror** and don't do anything else.

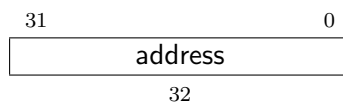


Field	Description	Access	Reset
address	Accesses bits 95:64 of the physical address in <b>sbaddress</b> (if the system address bus is that wide).	R/W	0

### 3.14.26 System Bus Address 127:96 (sbaddress3, at 0x37)

If **sbasize** is less than 97, then this register is not present.

When the system bus manager is busy, writes to this register will set **sbbusyerror** and don't do anything else.



Field	Description	Access	Reset
address	Accesses bits 127:96 of the physical address in <b>sbaddress</b> (if the system address bus is that wide).	R/W	0

### 3.14.27 System Bus Data 31:0 (sbddata0, at 0x3c)

If all of the **sbaccess** bits in **sbc** are 0, then this register is not present.

Any successful system bus read updates `sbddata`. If the width of the read access is less than the width of `sbddata`, the contents of the remaining high bits may take on any value.

If either `sberror` or `sbbusyerror` isn't 0 then accesses do nothing.

If the bus manager is busy then accesses set `sbbusyerror`, and don't do anything else.

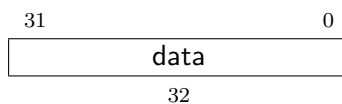
Writes to this register start the following:

1. Set `sbbusy`.
2. Perform a bus write of the new value of `sbddata` to `sbaddress`.
3. If the write succeeded and `sbautoincrement` is set, increment `sbaddress`.
4. Clear `sbbusy`.

Reads from this register start the following:

1. "Return" the data.
2. Set `sbbusy`.
3. If `sbreadondata` is set:
  - (a) Perform a system bus read from the address contained in `sbaddress`, placing the result in `sbddata`.
  - (b) If `sbautoincrement` is set and the read was successful, increment `sbaddress`.
4. Clear `sbbusy`.

Only `sbddata0` has this behavior. The other `sbddata` registers have no side effects. On systems that have buses wider than 32 bits, a debugger should access `sbddata0` after accessing the other `sbddata` registers.

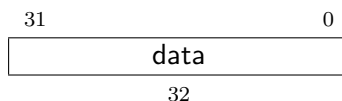


Field	Description	Access	Reset
<code>data</code>	Accesses bits 31:0 of <code>sbddata</code> .	R/W	0

### 3.14.28 System Bus Data 63:32 (`sbddata1`, at `0x3d`)

If `sbaccess64` and `sbaccess128` are 0, then this register is not present.

If the bus manager is busy then accesses set `sbbusyerror`, and don't do anything else.

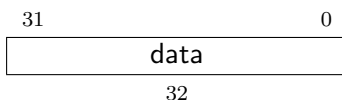


Field	Description	Access	Reset
<b>data</b>	Accesses bits 63:32 of <code>sbddata</code> (if the system bus is that wide).	R/W	0

### 3.14.29 System Bus Data 95:64 (`sbddata2`, at `0x3e`)

This register only exists if `sbaccess128` is 1.

If the bus manager is busy then accesses set `sbbusyerror`, and don't do anything else.

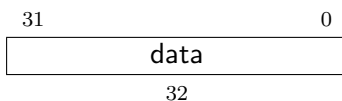


Field	Description	Access	Reset
<b>data</b>	Accesses bits 95:64 of <code>sbddata</code> (if the system bus is that wide).	R/W	0

### 3.14.30 System Bus Data 127:96 (`sbddata3`, at `0x3f`)

This register only exists if `sbaccess128` is 1.

If the bus manager is busy then accesses set `sbbusyerror`, and don't do anything else.



Field	Description	Access	Reset
<b>data</b>	Accesses bits 127:96 of <code>sbddata</code> (if the system bus is that wide).	R/W	0

### 3.14.31 Custom Features (`custom`, at `0x1f`)

This optional register may be used for non-standard features. Future version of the debug spec will not use this address.

### 3.14.32 Custom Features 0 (`custom0`, at `0x70`)

The optional `custom0` through `custom15` registers may be used for non-standard features. Future versions of the debug spec will not use these addresses.

## Chapter 4

# Sdext (ISA Extension)

This chapter describes the Sdext ISA extension. It must be implemented to make external debug work, and is only useful in conjunction with external debug.

Modifications to the RISC-V core to support debug are kept to a minimum. There is a special execution mode (Debug Mode) and a few extra CSRs. The DM takes care of the rest.

In order to be compatible with this specification an implementation must implement everything described in this chapter that is not explicitly listed as optional.

If Sdext is implemented and Sdtrig is not implemented, then accessing any of the Sdtrig CSRs must raise an illegal instruction exception.

### 4.1 Debug Mode

Debug Mode is a special processor mode used only when a hart is halted for external debugging. Because the hart is halted, there is no forward progress in the normal instruction stream. How Debug Mode is implemented is not specified here.

When executing code due to an abstract command, the hart stays in Debug Mode and the following apply:

1. All implemented instructions operate just as they do in M-mode, unless an exception is mentioned in this list.
2. All operations are executed with machine mode privilege, except that additional Debug Mode CSRs are accessible and MPRV in `mstatus` may be ignored according to `mprven`. Full permission checks, or a relaxed set of permission checks, will apply according to `relaxedpriv`.
3. All interrupts (including NMI) are masked.
4. Traps don't take place. Instead, they end execution of the program buffer and the hart remains in Debug Mode. Because they do not trap to M-mode, they do not update registers such as `mstatus`, `mepc`, `mcause`, `mtval`, `mtval2`, and `mtinst`. The same is true for the equivalent privileged registers that are updated when trapping to other modes. Registers that may be

- updated as part of execution before the exception are allowed to be updated. For example, vector load/store instructions which raise exceptions may partially update the destination register and set `vstart` appropriately.
5. Triggers don't match or fire.
  6. If `stopcount` is 0 then counters continue. If it is 1 then counters are stopped.
  7. If `stoptime` is 0 then `time` continues to update. If it is 1 then `time` will not update. It will resynchronize with `mtime` after leaving Debug Mode.
  8. Instructions that place the hart into a stalled state act as a `nop`. This includes `wfi`, `wrs.sto`, and `wrs.nto`.
  9. Almost all instructions that change the privilege mode have UNSPECIFIED behavior. This includes `ecall`, `mret`, `sret`, and `uret`. (To change the privilege mode, the debugger can write `prv` and `v` in `dcsr`). The only exception is `ebreak`, which ends execution of the Program Buffer when executed.
  10. All control transfer instructions may act as illegal instructions if their destination is in the Program Buffer. If one such instruction acts as an illegal instruction, all such instructions must act as illegal instructions.
  11. All control transfer instructions may act as illegal instructions if their destination is outside the Program Buffer. If one such instruction acts as an illegal instruction, all such instructions must act as illegal instructions.
  12. Instructions that depend on the value of the PC (e.g. `auipc`) may act as illegal instructions.
  13. Effective XLEN is DXLEN.
  14. Forward progress is guaranteed.

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*When `mprven = 1`, the external debugger can set MPRV and MPP appropriately to have hardware perform memory accesses with the appropriate endianness, address translation, permission checks, and PMP/PMA checks (subject to `relaxedpriv`). This is also the only way to access all of physical memory when 34-bit physical addresses are supported on a Sv32 hart. If hardware ties `mprven` to 0 then the external debugger is expected to simulate all the effects of MPRV, including any extensions that affect memory accesses. For these reasons it is recommended to tie `mprven` to 1.*

## 4.2 Load-Reserved/Store-Conditional Instructions

The reservation registered by an `lr` instruction on a memory address may be lost when entering Debug Mode or while in Debug Mode. This means that there may be no forward progress if Debug Mode is entered between `lr` and `sc` pairs.

---

*This is a behavior that debug users must be aware of. If they have a breakpoint set between a `lr` and `sc` pair, or are stepping through such code, the `sc` may never succeed. Fortunately in general use there will be very few instructions in such a sequence, and anybody debugging it will quickly notice that the reservation is not occurring. The solution in that case is to set a breakpoint on the first instruction after the `sc` and run to it. A higher level debugger may choose to automate this.*



### 4.3 Wait for Interrupt Instruction

If halt is requested while `wfi` is executing, then the hart must leave the stalled state, completing this instruction's execution, and then enter Debug Mode.

### 4.4 Wait-on-Reservation-Set Instructions

If halt is requested while `wrs.sto` or `wrs.nto` is executing, then the hart must leave the stalled state, completing this instruction's execution, and then enter Debug Mode.

### 4.5 Single Step

#### 4.5.1 Step Bit In Dcsr

This method is only available to external debuggers, and is the preferred way to single step.

An external debugger can cause a halted hart to execute a single instruction or trap and then re-enter Debug Mode by setting `step` before resuming. If `step` is set when a hart resumes then it will single step, regardless of the reason for resuming.

If control is transferred to a trap handler while executing the instruction, then Debug Mode is re-entered immediately after the PC is changed to the trap handler, and the appropriate `tval` and `cause` registers are updated. In this case none of the trap handler is executed, and if the cause was a pending interrupt no instructions might be executed at all.

If executing or fetching the instruction causes a trigger to fire with `action=1`, Debug Mode is re-entered immediately after that trigger has fired. In that case `cause` is set to 2 (trigger) instead of 4 (single step). Whether the instruction is executed or not depends on the specific configuration of the trigger.

If the instruction that is executed causes the PC to change to an address where an instruction fetch causes an exception, that exception does not occur until the next time the hart is resumed. Similarly, a trigger at the new address does not fire until the hart actually attempts to execute that instruction.

If the instruction being stepped over would normally stall the hart, then instead the instruction is treated as a `nop`. This includes `wfi`, `wrs.sto`, and `wrs.nto`.

#### 4.5.2 Icount Trigger

Native debuggers won't have access to `dcsr`, but can use the `icount` trigger by setting `count` to 1.

This approach does have some limitations:

1. Interrupts will fire as usual. Debuggers that want to disable interrupts while stepping must disable them by changing `mstatus`, and specially handle instructions that read `mstatus`.
2. `wfi` instructions are not treated specially and might take a very long time to complete.

This mechanism cleanly supports a system which supports multiple privilege levels, where the OS or a debug stub runs in M-Mode while the program being debugged runs in a less privileged mode. Systems that only support M-Mode can use `icount` as well, but count must be able to count several instructions (depending on the software implementation). See Section [B.3.1](#).

## 4.6 Reset

If the halt signal (driven by the hart's halt request bit in the Debug Module) or `resethaltreq` are asserted when a hart comes out of reset, the hart must enter Debug Mode before executing any instructions, but after performing any initialization that would usually happen before the first instruction is executed.

## 4.7 Halt

When a hart halts:

1. `cause` is updated.
2. `prv` and `v` are set to reflect current privilege mode.
3. `dpc` is set to the next instruction that should be executed.
4. If the current instruction can be partially executed and should be restarted to complete, then the relevant state for that is updated. E.g. if a halt occurs during a partially executed vector instruction, then `vstart` is updated, and `dpc` is updated to the address of the partially executed instruction. This is analogous to how vector instructions behave for exceptions.
5. The hart enters Debug Mode.

## 4.8 Resume

When a hart resumes:

1. `pc` changes to the value stored in `dpc`.
2. The current privilege mode and virtualization mode are changed to that specified by `prv` and `v`.
3. If the new privilege mode is less privileged than M-mode, `MPRV` in `mstatus` is cleared.
4. The hart is no longer in debug mode.

## 4.9 XLEN

While in Debug Mode, XLEN is DXLEN. It is up to the debugger to determine the XLEN during normal program execution (by looking at `misa`) and to clearly communicate this to the user.

## 4.10 Core Debug Registers

The supported Core Debug Registers must be implemented for each hart that can be debugged. They are CSRs, accessible using the RISC-V `csr` opcodes and optionally also using abstract debug commands.

Attempts to access an unimplemented Core Debug Register raise an illegal instruction exception.

These registers are only accessible from Debug Mode.

Table 4.1: Core Debug Registers

Address	Name	Page
0x7b0	Debug Control and Status ( <code>dcsr</code> )	<a href="#">53</a>
0x7b1	Debug PC ( <code>dpc</code> )	<a href="#">57</a>
0x7b2	Debug Scratch Register 0 ( <code>dscratch0</code> )	<a href="#">58</a>
0x7b3	Debug Scratch Register 1 ( <code>dscratch1</code> )	<a href="#">58</a>

### 4.10.1 Debug Control and Status (`dcsr`, at 0x7b0)

Upon entry into Debug Mode, `v` and `prv` are updated with the privilege level the hart was previously in, and `cause` is updated with the reason for Debug Mode entry. Other than these fields and `nmip`, the other fields of `dcsr` are only writable by the external debugger.

Table 4.2 shows the priorities of reasons for entering Debug Mode. Implementations should implement priorities as shown in the table. For compatibility with old versions of this spec, `resethaltreq` and `haltreq` are allowed to be at different positions than shown as long as:

1. `resethaltreq` is higher priority than `haltreq`
2. the relative order of the other four causes is maintained

Table 4.2: Priority of reasons for entering Debug Mode from highest to lowest.

<code>cause</code> encoding	Cause
5	<code>resethaltreq</code>
6	<code>halt group</code>
3	<code>haltreq</code>
2	trigger (See table 5.2 for detailed priority)
1	<code>ebreak</code>
4	<code>step</code>

Note that `mcontrol/mcontrol6` triggers which fire after the instruction which hit the trigger are considered to be high priority causes on the subsequent instruction. Therefore, an execute trigger with `timing=after` on an `ebreak` instruction is lower priority than the `ebreak` itself because the trigger will fire after the `ebreak` instruction. For the same reason, if a single instruction is stepped with both `icount` and `step` then the `step` has priority. See table 5.2 for the relative priorities of triggers with respect to the `ebreak` instruction.

Most multi-hart implementations will probably hardwire `stoptime` to 0, as the implementation can get complicated and the benefit is small.

This CSR is read/write.

31	28	27	18	17	16	15	14	13	12	11	
debugver	0		ebreakvs	ebreakvu	ebreakm	0	ebreaks	ebreaku	stepie		
4	10		1	1	1	1	1	1	1	1	
	10		9	8	6	5	4	3	2	1	0
	stopcount		stoptime	cause	v	mprven	nmip	step	prv		
	1		1	3	1	1	1	1	2		

Field	Description	Access	Reset
debugver	0 (none): There is no debug support. 4 (1.0): Debug support exists as it is described in this document. 15 (custom): There is debug support, but it does not conform to any available version of this spec.	R	Preset
ebreakvs	0 (exception): <code>ebreak</code> instructions in VS-mode behave as described in the Privileged Spec. 1 (debug mode): <code>ebreak</code> instructions in VS-mode enter Debug Mode. This bit is hardwired to 0 if the hart does not support virtualization mode.	WARL	0
ebreakvu	0 (exception): <code>ebreak</code> instructions in VU-mode behave as described in the Privileged Spec. 1 (debug mode): <code>ebreak</code> instructions in VU-mode enter Debug Mode. This bit is hardwired to 0 if the hart does not support virtualization mode.	WARL	0
ebreakm	0 (exception): <code>ebreak</code> instructions in M-mode behave as described in the Privileged Spec. 1 (debug mode): <code>ebreak</code> instructions in M-mode enter Debug Mode.	R/W	0
ebreaks	0 (exception): <code>ebreak</code> instructions in S-mode behave as described in the Privileged Spec. 1 (debug mode): <code>ebreak</code> instructions in S-mode enter Debug Mode. This bit is hardwired to 0 if the hart does not support S-mode.	WARL	0

*Continued on next page*

Field	Description	Access	Reset
<code>ebreaku</code>	<p>0 (exception): <code>ebreak</code> instructions in U-mode behave as described in the Privileged Spec.</p> <p>1 (debug mode): <code>ebreak</code> instructions in U-mode enter Debug Mode.</p> <p>This bit is hardwired to 0 if the hart does not support U-mode.</p>	WARL	0
<code>stepie</code>	<p>0 (interrupts disabled): Interrupts (including NMI) are disabled during single stepping with <code>step</code> set. This value should be supported.</p> <p>1 (interrupts enabled): Interrupts (including NMI) are enabled during single stepping with <code>step</code> set.</p> <p>Implementations may hard wire this bit to 0. In that case interrupt behavior can be emulated by the debugger.</p> <p>The debugger must not change the value of this bit while the hart is running.</p>	WARL	0
<code>stopcount</code>	<p>0 (normal): Increment counters as usual.</p> <p>1 (freeze): Don't increment any hart-local counters while in Debug Mode or on <code>ebreak</code> instructions that cause entry into Debug Mode. These counters include the <code>instret</code> CSR. On single-hart cores <code>cycle</code> should be stopped, but on multi-hart cores it must keep incrementing.</p> <p>An implementation may hardwire this bit to 0 or 1.</p>	WARL	Preset
<code>stoptime</code>	<p>0 (normal): <code>time</code> continues to reflect <code>mtime</code>.</p> <p>1 (freeze): <code>time</code> is frozen at the time that Debug Mode was entered. When leaving Debug Mode, <code>time</code> will reflect the latest value of <code>mtime</code> again.</p> <p>While all harts have <code>stoptime</code> =1 and are in Debug Mode, <code>mtime</code> is allowed to stop incrementing.</p> <p>An implementation may hardwire this bit to 0 or 1.</p>	WARL	Preset

*Continued on next page*

Field	Description	Access	Reset
cause	<p>Explains why Debug Mode was entered.</p> <p>When there are multiple reasons to enter Debug Mode in a single cycle, hardware should set <b>cause</b> to the cause with the highest priority. See table 4.2 for priorities.</p> <p>1 (ebreak): An <b>ebreak</b> instruction was executed.</p> <p>2 (trigger): A Trigger Module trigger fired with <b>action=1</b>.</p> <p>3 (haltreq): The debugger requested entry to Debug Mode using <b>haltreq</b>.</p> <p>4 (step): The hart single stepped because <b>step</b> was set.</p> <p>5 (resethaltreq): The hart halted directly out of reset due to <b>resethaltreq</b>. It is also acceptable to report 3 when this happens.</p> <p>6 (group): The hart halted because it's part of a halt group. Harts may report 3 for this cause instead.</p> <p>Other values are reserved for future use.</p>	R	0
v	<p>Extends the prv field with the virtualization mode the hart was operating in when Debug Mode was entered. The encoding is described in Table 4.6. A debugger can change this value to change the hart's virtualization mode when exiting Debug Mode. This bit is hardwired to 0 on harts that do not support virtualization mode.</p>	WARL	0
mprven	<p>0 (disabled): MPRV in <b>mstatus</b> is ignored in Debug Mode.</p> <p>1 (enabled): MPRV in <b>mstatus</b> takes effect in Debug Mode.</p> <p>Implementing this bit is optional. It may be tied to either 0 or 1.</p>	WARL	Preset
nmip	<p>When set, there is a Non-Maskable-Interrupt (NMI) pending for the hart.</p> <p>Since an NMI can indicate a hardware error condition, reliable debugging may no longer be possible once this bit becomes set. This is implementation-dependent.</p>	R	0
step	<p>When set and not in Debug Mode, the hart will only execute a single instruction and then enter Debug Mode. See Section 4.5.1 for details.</p> <p>The debugger must not change the value of this bit while the hart is running.</p>	R/W	0

*Continued on next page*

Field	Description	Access	Reset
prv	Contains the privilege mode the hart was operating in when Debug Mode was entered. The encoding is described in Table 4.6. A debugger can change this value to change the hart's privilege mode when exiting Debug Mode. Not all privilege modes are supported on all harts. If the encoding written is not supported or the debugger is not allowed to change to it, the hart may change to any supported privilege mode.	WARL	3

#### 4.10.2 Debug PC (dpc, at 0x7b1)

Upon entry to debug mode, `dpc` is updated with the virtual address of the next instruction to be executed. The behavior is described in more detail in Table 4.4.

Table 4.4: Virtual address in DPC upon Debug Mode Entry

Cause	Virtual Address in DPC
<code>ebreak</code>	Address of the <code>ebreak</code> instruction
single step	Address of the instruction that would be executed next if no debugging was going on. Ie. <code>pc + 4</code> for 32-bit instructions that don't change program flow, the destination PC on taken jumps/branches, etc.
trigger module	The address of the next instruction to be executed at the time that debug mode was entered. If the trigger is <code>mcontrol</code> and <code>timing</code> is 0 or if the trigger is <code>mcontrol6</code> and <code>hit1</code> is 0, this corresponds to the address of the instruction which caused the trigger to fire.
halt request	Address of the next instruction to be executed at the time that debug mode was entered

Executing the Program Buffer may cause the value of `dpc` to become UNSPECIFIED. If that is the case, it must be possible to read/write `dpc` using an abstract command with `postexec` not set. The debugger must attempt to save `dpc` between halting and executing a Program Buffer, and then restore `dpc` before leaving Debug Mode.

---

*Allowing `dpc` to become UNSPECIFIED upon Program Buffer execution allows for direct implementations that don't have a separate PC register, and do need to use the PC when executing the Program Buffer.*

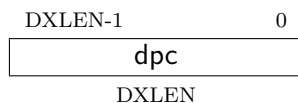
If the Access Register abstract command supports reading `dpc` while the hart is running, then the value read should be the address of a recently executed instruction.

If the Access Register abstract command supports writing `dpc` while the hart is running, then the executing program should jump to the written address shortly after the write occurs.

The writability of `dpc` follows the same rules as `mepc` as defined in the Privileged Spec. In particular, `dpc` must be able to hold all valid virtual addresses and the writability of the low bits depends on `IALIGN`.

When resuming, the hart's PC is updated to the virtual address stored in `dpc`. A debugger may write `dpc` to change where the hart resumes.

This CSR is read/write.



### 4.10.3 Debug Scratch Register 0 (`dscratch0`, at `0x7b2`)

Optional scratch register that can be used by implementations that need it. A debugger must not write to this register unless `hartinfo` explicitly mentions it (the Debug Module may use this register internally).

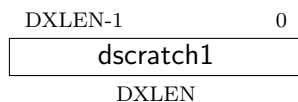
This CSR is read/write.



### 4.10.4 Debug Scratch Register 1 (`dscratch1`, at `0x7b3`)

Optional scratch register that can be used by implementations that need it. A debugger must not write to this register unless `hartinfo` explicitly mentions it (the Debug Module may use this register internally).

This CSR is read/write.



## 4.11 Virtual Debug Registers

A virtual register is one that doesn't exist directly in the hardware, but that the debugger exposes as if it does. Debug software should implement them, but hardware can skip this section. Virtual registers exist to give users access to functionality that's not part of standard debuggers without requiring them to carefully modify debug registers while the debugger is also accessing those same registers.



Table 4.6: Privilege Mode and Virtualization Mode Encoding

H extension supported	v	prv	Abbreviation	Name
No	0	0	U-mode	User mode
No	0	1	S-mode	Supervisor mode
No	0	3	M-mode	Machine mode
Yes	0	0	U-mode	User mode
Yes	0	1	HS-mode	Hypervisor-enabled supervisor mode
Yes	0	3	M-mode	Machine mode
Yes	1	0	VU-mode	Virtual user mode
Yes	1	1	VS-mode	Virtual supervisor mode

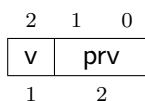
Table 4.5: Virtual Core Debug Registers

Address	Name	Page
virtual	Privilege Mode ( <code>priv</code> )	<a href="#">59</a>

#### 4.11.1 Privilege Mode (`priv`, at `virtual`)

Users can read this register to inspect the privilege mode that the hart was running in when the hart halted. Users can write this register to change the privilege mode that the hart will run in when it resumes.

This register contains `priv` and `v` from `dcsr`, but in a place that the user is expected to access. The user should not access `dcsr` directly, because doing so might interfere with the debugger.



Field	Description	Access	Reset
v	Contains the virtualization mode the hart was operating in when Debug Mode was entered. The encoding is described in Table 4.6, and matches the virtualization mode encoding from the Privileged Spec. A user can write this value to change the hart's virtualization mode when exiting Debug Mode.	WARL	0

*Continued on next page*

Field	Description	Access	Reset
<b>prv</b>	Contains the privilege mode the hart was operating in when Debug Mode was entered. The encoding is described in Table 4.6, and matches the privilege mode encoding from the Privileged Spec. A user can write this value to change the hart's privilege mode when exiting Debug Mode.	R/W	0

## Chapter 5

# Sdtrig (ISA Extension)

This chapter describes the Sdtrig ISA extension, which can be implemented independently of functionality described in the other chapters. It consists exclusively of the Trigger Module (TM).

Triggers can cause a breakpoint exception, entry into Debug Mode, or a trace action without having to execute a special instruction. This makes them invaluable when debugging code from ROM. They can trigger on execution of instructions at a given memory address, or on the address/data in loads/stores.

If Sdtrig is implemented, the Trigger Module must support at least one trigger. Accessing trigger CSRs that are not used by any of the implemented triggers must result in an illegal instruction exception. M-Mode and Debug Mode accesses to trigger CSRs that are used by any of the implemented triggers must succeed, regardless of the current type of the currently selected trigger.

A trigger matches when the conditions that it specifies (e.g. a load from a specific address) are met. A trigger fires when a trigger that matches performs the action configured for that trigger.

Triggers do not fire while in Debug Mode.

### 5.1 Enumeration

Each trigger may support a variety of features. A debugger can build a list of all triggers and their features as follows:

1. Write 0 to `tselect`. If this results in an illegal instruction exception, then there are no triggers implemented.
2. Read back `tselect` and check that it contains the written value. If not, exit the loop.
3. Read `tinfo`.
4. If that caused an exception, the debugger must read `tdata1` to discover the type. (If `type` is 0, this trigger doesn't exist. Exit the loop.)
5. If `info` is 1, this trigger doesn't exist. Exit the loop.
6. Otherwise, the selected trigger supports the types discovered in `info`.

7. Repeat, incrementing the value in `tselect`.

---

The above algorithm reads back `tselect` so that implementations which have  $2^n$  triggers only need to implement  $n$  bits of `tselect`.

The algorithm checks `tinfo` and `type` in case the implementation has  $m$  bits of `tselect` but fewer than  $2^m$  triggers.

## 5.2 Actions

Triggers can be configured to take one of several actions when they fire. Table 5.1 lists all options.

Table 5.1: `action` encoding

Value	Description
0	Raise a breakpoint exception. (Used when software wants to use the trigger module without an external debugger attached.) <code>xepc</code> must contain the virtual address of the next instruction that must be executed to preserve the program flow.
1	Enter Debug Mode. <code>dpc</code> must contain the virtual address of the next instruction that must be executed to preserve the program flow. This action is only legal when the trigger's <code>dmode</code> is 1. Since the <code>tdata</code> registers are WARL, hardware should clear the action field whenever the action field is 1, the new value of <code>dmode</code> would be 0, and the new value of the action field would be 1. This action can only be supported if <code>Sdext</code> is implemented on the hart.
2	Trace on, described in the trace specification.
3	Trace off, described in the trace specification.
4	Trace notify, described in the trace specification.
5	Reserved for use by the trace specification.
8 – 9	Send a signal to TM external trigger output 0 or 1 (respectively).
other	Reserved for future use.

---

Actions 8 and 9 are intended to increment custom event counters, but these signals could also be brought to outputs for use by external logic.

## 5.3 Priority

Table 5.2 lists the synchronous exceptions from the Privileged Spec, and where the various types of triggers fit in. The first 3 columns come from the Privileged Spec, and the final column shows

where triggers fit in. Priorities in the table are separated by horizontal lines, so e.g. etrigger and itrigger have the same priority. If this table contradicts the table in the Privileged Spec, then the latter takes precedence.

This table only applies if triggers are precise. Otherwise triggers will fire some indeterminate time after the event, and the priority is irrelevant. When triggers are chained, the priority is the lowest priority of the triggers in the chain.

Priority	Exception Code	Description	Trigger
<i>Highest</i>	3 3 3 3		etrigger icount itrigger mcontrol/mcontrol6 after (on previous instruction)
	3	Instruction address breakpoint	mcontrol/mcontrol6 execute address before
	12, 20, 1	During instruction address translation: First encountered page fault, guest-page fault, or access fault	
	1	With physical address for instruction: Instruction access fault	
	3		mcontrol/mcontrol6 execute data before
	2 22 0 8, 9, 10, 11 3 3	Illegal instruction Virtual instruction Instruction address misaligned Environment call Environment break Load/Store/AMO address breakpoint	mcontrol/mcontrol6 load/store address/data before
	4, 6	Optionally: Load/Store/AMO address misaligned	
	13, 15, 21, 23, 5, 7	During address translation for an explicit memory access: First encountered page fault, guest-page fault, or access fault	
	5, 7	With physical address for an explicit memory access: Load/store/AMO access fault	
	4, 6	If not higher priority: Load/store/AMO address misaligned	
<i>Lowest</i>	3		mcontrol/mcontrol6 load data before

Table 5.2: Synchronous exception priority in decreasing priority order.

When multiple triggers in the same priority fire at once, [hit](#) (if implemented) is set for all of them.

If more than one of these triggers has `action = 0` then `tval` is updated in accordance with one of them, but which one is UNSPECIFIED. If one of these triggers has the “enter Debug Mode” action (1) and another trigger has the “raise a breakpoint exception” action (0), the preferred behavior is to have both actions take place. It is implementation-dependent which of the two happens first. This ensures both that the presence of an external debugger doesn’t affect execution and that a trigger set by user code doesn’t affect the external debugger. If this is not implemented, then the hart must enter Debug Mode and ignore the breakpoint exception. In the latter case, `hit` of the trigger whose action is 0 must still be set, giving a debugger an opportunity to handle this case. What happens with trace actions when triggers with different actions are also firing is left to the trace specification.

## 5.4 Native Triggers

Triggers can be used for native debugging when `action = 0`. If supported by the hart and desired by the debugger, triggers will often be programmed to have `m = 0` so that when they fire they cause a breakpoint exception to trap to a more privileged mode. That breakpoint exception can either be taken in M-mode or it can be delegated to a less privileged mode. However, it is possible for triggers to fire in the same mode that the resulting exception will be handled in.

In these cases such a trigger may cause a breakpoint exception while already in a trap handler. This might leave the hart unable to resume normal execution because state such as `mcause` and `mepc` would be overwritten.

---

*In particular, when `action = 0`:*

1. *`mcontrol` and `mcontrol6` triggers with `m = 1` can cause a breakpoint exception that is taken from M-mode to M-mode (regardless of delegation).*
2. *`mcontrol` and `mcontrol6` triggers with `s = 1` can cause a breakpoint exception that is taken from S-mode to S-mode if `medeleg [3]=1`.*
3. *`mcontrol6` triggers with `vs = 1` can cause a breakpoint exception that is taken from VS-mode to VS-mode if `medeleg [3]=1` and `hedeleg [3]=1`.*
4. *`icount` triggers with `m = 1` can cause a breakpoint exception that is taken from M-mode to M-mode (regardless of delegation).*
5. *`icount` triggers with `s = 1` can cause a breakpoint exception that is taken from S-mode to S-mode if `medeleg [3]=1`.*
6. *`icount` triggers with `vs = 1` can cause a breakpoint exception that is taken from VS-mode to VS-mode if `medeleg [3]=1` and `hedeleg [3]=1`.*
7. *`etrigger` and `itrigger` triggers will always be taken from a trap handler before the first instruction of the handler. If `etrigger/itrigger` is set to trigger on exception/interrupt `X` and if `X` is delegated to mode `Y` then the trigger will cause a breakpoint exception that is taken from mode `Y` to mode `Y` unless breakpoint exceptions are delegated to a more privileged mode than `Y`.*
8. *`tmexttrigger` triggers are asynchronous and may occur in any mode and at any time.*

Harts that support triggers with `action = 0` should implement one of the following two solutions to solve the problem of reentrancy:

1. The hardware prevents triggers with `action = 0` from matching or firing while in M-mode and while MIE in `mstatus` is 0. If `medeleg [3]=1` then it prevents triggers with `action = 0`

from matching or firing while in S-mode and while SIE in `sstatus` is 0. If `medeleg` [3]=1 and `hedeleg` [3]=1 then it prevents triggers with `action` =0 from matching or firing while in VS-mode and while SIE in `vsstatus` is 0.

2. `mte` and `mpte` in `tcontrol` is implemented. `medeleg` [3] is hard-wired to 0.

---

*The first option has the limitation that interrupts might be disabled at times when a user still might want triggers to fire. It has the benefit that breakpoints are not required to be handled in M-mode.*

*The second option has the benefit that it only disables triggers during the trap handler, though it requires specific software support for this debug feature in the M-mode trap handlers. It can only work if breakpoints are not delegated to less privileged modes and therefore targets primarily implementations without S-mode.*

*Because `tcontrol` is not accessible to S-mode, the second option can not be extended to accommodate delegation without adding additional S-mode and VS-mode CSRs.*

*Both options prevent `etrigger` and `itrigger` from having any effect on exceptions and interrupts that are handled in M-mode. They also prevent triggering during some initial portion of each handler. Debuggers should use other mechanisms to debug these cases, such as patching the handler or setting a breakpoint on the instruction after MIE is cleared.*

## 5.5 Memory Access Triggers

`mcontrol` and `mcontrol6` both enable triggers on memory accesses. This section describes for both of them how certain corner cases are treated.

### 5.5.1 A Extension

If the A extension is supported, then triggers on loads/stores treat them as follows:

1. `lr` instructions are loads.
2. Successful `sc` instructions are stores.
3. It is UNSPECIFIED whether failing `sc` instructions are stores or not.
4. Each AMO instruction is a load for the read portion of the operation. The address is always available to trigger on, although the value loaded might not be, depending on the hardware implementation.
5. Each AMO instruction is a store for the write portion of the operation. The address is always available to trigger on, although the value stored might not be, depending on the hardware implementation.

If the destination register of any load or AMO is `zero` then it is UNSPECIFIED whether a data load trigger will match. Whether data store triggers match on AMOs is UNSPECIFIED.

### 5.5.2 Combined Accesses

Some instructions lead a hart to perform multiple memory accesses. This includes vector loads and stores, as well as `cm.push` and `cm.pop` instructions. The Trigger Module should match such accesses

as if they all happened individually. E.g. a vector load should be treated as if it performed multiple loads of size SEW (selected element width), and `cm.push` should be treated as if it performed multiple stores of size XLEN.

### 5.5.3 Cache Operations

Cache operations are infrequently performed, and code that uses them can have hard-to-find bugs. For the purposes of debug triggers, two classes of cache operations must match as stores:

1. Cache operations that enable software to maintain coherence between otherwise non-coherent implicit and explicit memory accesses.
2. Cache operations that perform block writes of constant data.

Only triggers with `size = 0` and `select = 0` will match. Since cache operations affect multiple addresses, there are multiple possible values to compare against. Implementations must implement one of the following options. From most desirable to least desirable, they are:

1. Every address from the effective address rounded down to the nearest cache block boundary (inclusive) to the effective address rounded up to the nearest cache block boundary (exclusive) is a compare value.
2. The effective address rounded down to the nearest cache block boundary is a compare value.
3. The effective address of the instruction is a compare value.

Cache operations encoded as HINTs do not match debug triggers.

---

*The above language intends to capture the trigger behavior with respect to the cache operations to be introduced in a forthcoming I/D consistency extension.*

*For RISC-V Base Cache Management Operation ISA Extensions 1.0.1, this means the following:*

1. `cbo.clean`, `cbo.flush`, and `cbo.inval` match as if they are stores because they affect consistency.
2. `cbo.zero` matches as if it is a store because it performs a block write of constant data.
3. The `prefetch` instructions don't match at all.

## 5.6 Multiple State Change Instructions

An instruction that performs multiple architectural state changes (e.g., register updates and/or memory accesses) might cause a trigger to fire at an intermediate point in its execution. As a result, architectural state changes up to that point might have been performed, while subsequent state changes, starting from the event that activated the trigger, might not have been. The definition of such an instruction will specify the order in which architectural state changes take place. Alternatively, it may state that partial execution is not allowed, implying that a mid-execution trigger must prevent any architectural state changes from occurring.

Debuggers won't be aware if an instruction has been partially executed. When they resume execution, they will execute the same instruction once more. Therefore, it's crucial that partially



executing the instruction and then executing it again leaves the hart in a state closely resembling the state it would have been in if the instruction had only been executed once.

## 5.7 Trigger Registers

These registers are CSRs, accessible using the RISC-V `csr` opcodes and optionally also using abstract debug commands.

Almost all trigger functionality is optional. All `tdata` registers follow write-any-read-legal semantics. If a debugger writes an unsupported configuration, the register will read back a value that is supported (which may simply be a disabled trigger). This means that a debugger must always read back values it writes to `tdata` registers, unless it already knows already what is supported. Writes to one `tdata` register must not modify the contents of other `tdata` registers, nor the configuration of any trigger besides the one that is currently selected.

The combination of these rules means that a debugger cannot simply set a trigger by writing `tdata1`, then `tdata2`, etc. The current value of `tdata2` might not be legal with the new value of `tdata1`. To help with this situation, it is guaranteed that writing 0 to `tdata1` disables the trigger, and leaves it in a state where `tdata2` and `tdata3` can be written with any value that makes sense for any trigger type supported by this trigger.

As a result, a debugger can write any supported trigger as follows:

1. Write 0 to `tdata1`. (This will result in `tdata1` containing a non-zero value, since the register is **WARL**.)
2. Write desired values to `tdata2` and `tdata3`.
3. Write desired value to `tdata1`.

Code that restores CSR context of triggers that might be configured to fire in the current privilege mode must use this same sequence to restore the triggers. This avoids the problem of a partially written trigger firing at a different time than is expected.

Attempts to access an unimplemented Trigger Register raise an illegal instruction exception.

The trigger registers, except `mscontext`, `scontext`, and `hcontext`, are only accessible in machine and Debug Mode to prevent untrusted user code from causing entry into Debug Mode without the OS's permission.

In this section `XLEN` means `MXLEN` when in M-mode, and `DXLEN` when in Debug Mode. On systems where those values of `XLEN` can differ, this is handled as follows. Fields retain their values regardless of `XLEN`, which only affects where in the register these fields appear (e.g. `type`). Some fields are wider when `XLEN` is 64 than when it is 32 (e.g. `svalue`). The high bits in such fields retain their value but are not readable when `XLEN` is 32. A modification of a register when `XLEN` is 32 clears any inaccessible bits in that register.

Table 5.3: Trigger Registers

Address	Name	Page
0x5a8	Supervisor Context ( <code>scontext</code> )	72
0x6a8	Hypervisor Context ( <code>hcontext</code> )	72
0x7a0	Trigger Select ( <code>tselect</code> )	68
0x7a1	Trigger Data 1 ( <code>tdata1</code> )	68
0x7a1	Match Control ( <code>mcontrol</code> )	74
0x7a1	Match Control Type 6 ( <code>mcontrol6</code> )	80
0x7a1	Instruction Count ( <code>icount</code> )	87
0x7a1	Interrupt Trigger ( <code>itrigger</code> )	89
0x7a1	Exception Trigger ( <code>etrigger</code> )	90
0x7a1	External Trigger ( <code>tmexttrigger</code> )	91
0x7a2	Trigger Data 2 ( <code>tdata2</code> )	70
0x7a3	Trigger Data 3 ( <code>tdata3</code> )	70
0x7a3	Trigger Extra (RV32) ( <code>textra32</code> )	92
0x7a3	Trigger Extra (RV64) ( <code>textra64</code> )	94
0x7a4	Trigger Info ( <code>tinfo</code> )	71
0x7a5	Trigger Control ( <code>tcontrol</code> )	71
0x7a8	Machine Context ( <code>mcontext</code> )	73
0x7aa	Machine Supervisor Context ( <code>mscontext</code> )	73

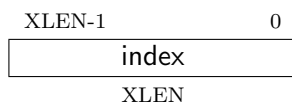
### 5.7.1 Trigger Select (`tselect`, at 0x7a0)

This register determines which trigger is accessible through the other trigger registers. It is optional if no triggers are implemented. The set of accessible triggers must start at 0, and be contiguous.

This register is **WARL**. Writes of values greater than or equal to the number of supported triggers may result in a different value in this register than what was written or may point to a trigger where `type = 0`. To verify that what they wrote is a valid index, debuggers can read back the value and check that `tselect` holds what they wrote and read `tdata1` to see that `type` is non-zero.

Since triggers can be used both by Debug Mode and M-mode, the external debugger must restore this register if it modifies it.

This CSR is read/write.



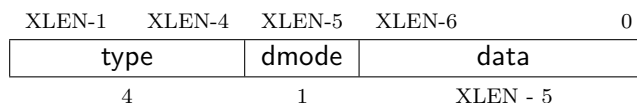
### 5.7.2 Trigger Data 1 (`tdata1`, at 0x7a1)

This register is optional if no triggers are implemented.

Writing 0 to this register must result in a trigger that is disabled. If this trigger supports multiple

types, then the hardware should disable it by changing `type` to 15.

This CSR is read/write.



Field	Description	Access	Reset
<code>type</code>	<p>0 (none): There is no trigger at this <code>tselect</code>.</p> <p>1 (legacy): The trigger is a legacy SiFive address match trigger. These should not be implemented and aren't further documented here.</p> <p>2 (mcontrol): The trigger is an address/data match trigger. The remaining bits in this register act as described in <code>mcontrol</code>.</p> <p>3 (icount): The trigger is an instruction count trigger. The remaining bits in this register act as described in <code>icount</code>.</p> <p>4 (itrigger): The trigger is an interrupt trigger. The remaining bits in this register act as described in <code>itrigger</code>.</p> <p>5 (etrigger): The trigger is an exception trigger. The remaining bits in this register act as described in <code>etrigger</code>.</p> <p>6 (mcontrol6): The trigger is an address/data match trigger. The remaining bits in this register act as described in <code>mcontrol6</code>. This is similar to a type 2 trigger, but provides additional functionality and should be used instead of type 2 in newer implementations.</p> <p>7 (tmexttrigger): The trigger is a trigger source external to the TM. The remaining bits in this register act as described in <code>tmexttrigger</code>.</p> <p>12–14 (custom): These trigger types are available for non-standard use.</p> <p>15 (disabled): This trigger is disabled. In this state, <code>tdata2</code> and <code>tdata3</code> can be written with any value that is supported for any of the types this trigger implements. The remaining bits in this register, except for <code>dmode</code>, are ignored. Other values are reserved for future use.</p>	WARL	Preset

*Continued on next page*

Field	Description	Access	Reset
<b>dmode</b>	If <b>type</b> is 0, then this bit is hard-wired to 0. 0 (both): Both Debug and M-mode can write the <b>tdata</b> registers at the selected <b>tselect</b> . 1 (dmode): Only Debug Mode can write the <b>tdata</b> registers at the selected <b>tselect</b> . Writes from other modes are ignored. This bit is only writable from Debug Mode. In ordinary use, external debuggers will always set this bit when configuring a trigger. When clearing this bit, debuggers should also set the action field (whose location depends on <b>type</b> ) to something other than 1.	WARL	0
<b>data</b>	If <b>type</b> is 0, then this field is hard-wired to 0. Trigger-specific data.	WARL	Preset

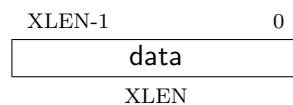
### 5.7.3 Trigger Data 2 (tdata2, at 0x7a2)

Trigger-specific data. It is optional if no implemented triggers use it.

If the trigger is disabled, then this register can be written with any value supported by any of the trigger types supported by this trigger.

If XLEN is less than DXLEN, writes to this register are sign-extended.

This CSR is read/write.



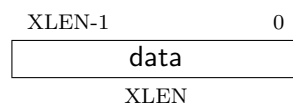
### 5.7.4 Trigger Data 3 (tdata3, at 0x7a3)

Trigger-specific data. It is optional if no implemented triggers use it.

If the trigger is disabled, then this register can be written with any value supported by any of the trigger types supported by this trigger.

If XLEN is less than DXLEN, writes to this register are sign-extended.

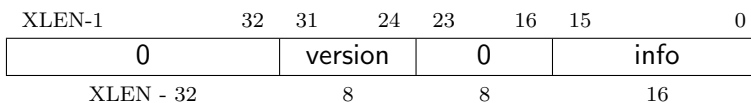
This CSR is read/write.



### 5.7.5 Trigger Info (tinfo, at 0x7a4)

This register is optional if no triggers are implemented, or if `type` is not writable and `version` would be 0. In this case the debugger can read the only supported type from `tdata1`.

Writing this read/write CSR has no effect.

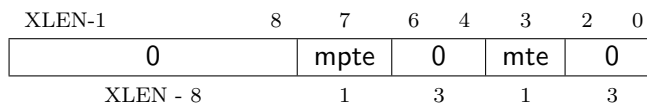


Field	Description	Access	Reset
version	Contains the version of the Sdtrig extension implemented. 0 (0): Supports triggers as described in this spec at commit 5a5c078, made on February 2, 2023. In these older versions: <ol style="list-style-type: none"> <li>1. <code>mcontrol6</code> has a timing bit identical to <code>timing</code></li> <li>2. <code>hit0</code> behaves just as <code>hit</code>.</li> <li>3. <code>hit1</code> is read-only 0.</li> <li>4. Encodings for <code>size</code> for access sizes larger than 64 bits are different.</li> </ol> 1 (1): Supports triggers as described in the ratified version 1.0 of this document.	R	Preset
info	One bit for each possible <code>type</code> enumerated in <code>tdata1</code> . Bit N corresponds to type N. If the bit is set, then that type is supported by the currently selected trigger. If the currently selected trigger doesn't exist, this field contains 1.	R	Preset

### 5.7.6 Trigger Control (tcontrol, at 0x7a5)

This optional register is only accessible in M-mode and Debug Mode and provides various control bits related to triggers.

This CSR is read/write.





Field	Description	Access	Reset
data	Supervisor mode software can write a context number to this register, which can be used to set triggers that only fire in that specific context. An implementation may tie any number of high bits in this field to 0. It's recommended to implement no more than 16 bits on RV32, and 34 on RV64.	WARL	0

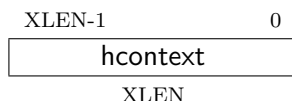
### 5.7.9 Machine Context (mcontext, at 0x7a8)

This register must be implemented if `hcontext` is implemented, and is optional otherwise. It is only accessible in M-mode and Debug mode.

---

`hcontext` is primarily useful to set triggers on hypervisor systems that only fire when a given VM is executing. It is also useful in systems where M-Mode implements something like a hypervisor directly.

This CSR is read/write.



Field	Description	Access	Reset
hcontext	M-Mode or HS-Mode (using <code>hcontext</code> ) software can write a context number to this register, which can be used to set triggers that only fire in that specific context. An implementation may tie any number of upper bits in this field to 0. If the H extension is not implemented, it's recommended to implement no more than 6 bits on RV32 and 13 on RV64 (as visible through the <code>mcontext</code> register). If the H extension is implemented, it's recommended to implement no more than 7 bits on RV32 and 14 on RV64.	WARL	0

### 5.7.10 Machine Supervisor Context (mscontext, at 0x7aa)

This optional register is an alias for `scontext`. It is only accessible in S/HS-mode, M-mode and Debug Mode. It is included for backward compatibility with version 0.13.

---

*The encoding of this CSR does not conform to the CSR Address Mapping Convention in the Privileged Spec. It is expected that new implementations will not support this encoding and that new debuggers will not use this CSR if `scontext` is available.*

### 5.7.11 Match Control (mcontrol, at 0x7a1)

This register is accessible as `tdata1` when `type` is 2. This trigger type is deprecated. It is included for backward compatibility with version 0.13.

---

*This trigger type only supports a subset of features of the newer `mcontrol6`. It is expected that new implementations will not support this trigger type and that new debuggers will not use it if `mcontrol6` is available.*

Address and data trigger implementation are heavily dependent on how the processor core is implemented. To accommodate various implementations, execute, load, and store address/data triggers may fire at whatever point in time is most convenient for the implementation. The debugger may request specific timings as described in [timing](#). Table 5.10 suggests timings for the best user experience.

A chain of triggers that don't all have the same [timing](#) value will never fire. That means to implement the suggestions in Table 5.10, both timings should be supported on load address triggers.

This trigger type may be limited to address comparisons (`select` is always 0) only. If that is the case and masking is not supported (match values 4, 5, 12, 13), then `tdata2` must be able to hold all valid virtual addresses but it need not be capable of holding other values.

The Privileged Spec says that breakpoint exceptions that occur on instruction fetches, loads, or stores update the `tval` CSR with either zero or the faulting virtual address. The faulting virtual address for an mcontrol trigger with `action` =0 is the address being accessed and which caused that trigger to fire. If multiple mcontrol triggers are chained then the faulting virtual address is the address which caused any of the chained triggers to fire.

If `textra32` or `textra64` are implemented for this trigger, it only matches when the conditions set there are satisfied.

This CSR is read/write.

XLEN-1				XLEN-4				XLEN-5				XLEN-6				XLEN-11				XLEN-12				23		22		21		20		19	
type				dmode				maskmax				0				sizehi		hit		select													
4				1				6				XLEN - 34				2		1		1													
18		17		16		15		12		11		10		7		6		5		4		3		2		1		0					
timing		sizelo		action		chain		match		m		0		s		u		execute		store		load											
1		2		4		1		4		1		1		1		1		1		1		1											



Field	Description	Access	Reset
maskmax	Specifies the largest naturally aligned powers-of-two (NAPOT) range supported by the hardware when <code>match</code> is 1. The value is the logarithm base 2 of the number of bytes in that range. A value of 0 indicates <code>match</code> 1 is not supported. A value of 63 corresponds to the maximum NAPOT range, which is $2^{63}$ bytes in size.	R	Preset
sizehi	This field only exists when XLEN is at least 64. It contains the 2 high bits of the access size. The low bits come from <code>sizelo</code> . See <code>sizelo</code> for how this is used.	WARL	0
hit	If this bit is implemented then it must become set when this trigger fires and may become set when this trigger matches. The trigger's user can set or clear it at any time. It is used to determine which trigger(s) matched. If the bit is not implemented, it is always 0 and writing it has no effect.	WARL	0
select	This bit determines the contents of the XLEN-bit compare values. 0 (address): There is at least one compare value and it contains the lowest virtual address of the access. It is recommended that there are additional compare values for the other accessed virtual addresses. (E.g. on a 32-bit read from 0x4000, the lowest address is 0x4000 and the other addresses are 0x4001, 0x4002, and 0x4003.) 1 (data): There is exactly one compare value and it contains the data value loaded or stored, or the instruction executed. Any bits beyond the size of the data access will contain 0.	WARL	0

*Continued on next page*

Field	Description	Access	Reset
timing	<p>0 (before): The action for this trigger will be taken just before the instruction that triggered it is retired, but after all preceding instructions are retired. <code>xepc</code> or <code>dpc</code> (depending on <code>action</code>) must be set to the virtual address of the instruction that matched.</p> <p>If this is combined with <code>load</code> and <code>select</code> =1 then a memory access will be performed (including any side effects of performing such an access) even though the load will not update its destination register. Debuggers should consider this when setting such breakpoints on, for example, memory-mapped I/O addresses.</p> <p>If an instruction matches this trigger and the instruction performs multiple memory accesses, it is UNSPECIFIED which memory accesses have completed before the trigger fires.</p> <p>1 (after): The action for this trigger will be taken after the instruction that triggered it is retired. It should be taken before the next instruction is retired, but it is better to implement triggers imprecisely than to not implement them at all. <code>xepc</code> or <code>dpc</code> (depending on <code>action</code>) must be set to the virtual address of the next instruction that must be executed to preserve the program flow.</p> <p>Most hardware will only implement one timing or the other, possibly dependent on <code>select</code>, <code>execute</code>, <code>load</code>, and <code>store</code>. This bit primarily exists for the hardware to communicate to the debugger what will happen. Hardware may implement the bit fully writable, in which case the debugger has a little more control.</p> <p>Data load triggers with <code>timing</code> of 0 will result in the same load happening again when the debugger lets the hart run. For data load triggers, debuggers must first attempt to set the breakpoint with <code>timing</code> of 1.</p> <p>If a trigger with <code>timing</code> of 0 matches, it is implementation-dependent whether that prevents a trigger with <code>timing</code> of 1 matching as well.</p>	WARL	0

*Continued on next page*

Field	Description	Access	Reset
size0	<p>This field contains the 2 low bits of the access size. The high bits come from <a href="#">sizehi</a>. The combined value is interpreted as follows:</p> <p>0 (any): The trigger will attempt to match against an access of any size. The behavior is only well-defined if <code>select = 0</code>, or if the access size is XLEN.</p> <p>1 (8bit): The trigger will only match against 8-bit memory accesses.</p> <p>2 (16bit): The trigger will only match against 16-bit memory accesses or execution of 16-bit instructions.</p> <p>3 (32bit): The trigger will only match against 32-bit memory accesses or execution of 32-bit instructions.</p> <p>4 (48bit): The trigger will only match against execution of 48-bit instructions.</p> <p>5 (64bit): The trigger will only match against 64-bit memory accesses or execution of 64-bit instructions.</p> <p>6 (80bit): The trigger will only match against execution of 80-bit instructions.</p> <p>7 (96bit): The trigger will only match against execution of 96-bit instructions.</p> <p>8 (112bit): The trigger will only match against execution of 112-bit instructions.</p> <p>9 (128bit): The trigger will only match against 128-bit memory accesses or execution of 128-bit instructions.</p> <p>An implementation must support the value of 0, but all other values are optional. When an implementation supports address triggers (<code>select = 0</code>), it is recommended that those triggers support every access size that the hart supports, as well as for every instruction size that the hart supports. Implementations such as RV32D or RV64V are able to perform loads and stores that are wider than XLEN. Custom extensions may also support instructions that are wider than XLEN. Because <code>tdata2</code> is of size XLEN, there is a known limitation that data value triggers (<code>select = 1</code>) can only be supported for access sizes up to XLEN bits. When an implementation supports data value triggers (<code>select = 1</code>), it is recommended that those triggers support every access size up to XLEN that the hart supports, as well as for every instruction length up to XLEN that the hart supports.</p>	WARL	0

*Continued on next page*

Field	Description	Access	Reset
action	The action to take when the trigger fires. The values are explained in Table 5.1.	WARL	0
chain	<p>0 (disabled): When this trigger matches, the configured action is taken.</p> <p>1 (enabled): While this trigger does not match, it prevents the trigger with the next index from matching.</p> <p>A trigger chain starts on the first trigger with <code>chain = 1</code> after a trigger with <code>chain = 0</code>, or simply on the first trigger if that has <code>chain = 1</code>. It ends on the first trigger after that which has <code>chain = 0</code>. This final trigger is part of the chain. The action on all but the final trigger is ignored. The action on that final trigger will be taken if and only if all the triggers in the chain match at the same time. Debuggers should not terminate a chain with a trigger with a different type. It is undefined when exactly such a chain fires.</p> <p>Because <code>chain</code> affects the next trigger, hardware must zero it in writes to <code>mcontrol</code> that set <code>dmode</code> to 0 if the next trigger has <code>dmode</code> of 1. In addition hardware should ignore writes to <code>mcontrol</code> that set <code>dmode</code> to 1 if the previous trigger has both <code>dmode</code> of 0 and <code>chain</code> of 1. Debuggers must avoid the latter case by checking <code>chain</code> on the previous trigger if they're writing <code>mcontrol</code>.</p> <p>Implementations that wish to limit the maximum length of a trigger chain (eg. to meet timing requirements) may do so by zeroing <code>chain</code> in writes to <code>mcontrol</code> that would make the chain too long.</p>	WARL	0

*Continued on next page*

Field	Description	Access	Reset
match	<p>0 (equal): Matches when any compare value equals <code>tdata2</code>.</p> <p>1 (napot): Matches when the top <math>M</math> bits of any compare value match the top <math>M</math> bits of <code>tdata2</code>. <math>M</math> is <math>XLEN - 1</math> minus the index of the least-significant bit containing 0 in <code>tdata2</code>. Debuggers should only write values to <code>tdata2</code> such that <math>M + \text{maskmax} \geq XLEN</math> and <math>M &gt; 0</math>, otherwise it's undefined on what conditions the trigger will match.</p> <p>2 (ge): Matches when any compare value is greater than (unsigned) or equal to <code>tdata2</code>.</p> <p>3 (lt): Matches when any compare value is less than (unsigned) <code>tdata2</code>.</p> <p>4 (mask low): Matches when <math>\frac{XLEN}{2} - 1:0</math> of any compare value equals <math>\frac{XLEN}{2} - 1:0</math> of <code>tdata2</code> after <math>\frac{XLEN}{2} - 1:0</math> of the compare value is ANDed with <math>XLEN - 1:\frac{XLEN}{2}</math> of <code>tdata2</code>.</p> <p>5 (mask high): Matches when <math>XLEN - 1:\frac{XLEN}{2}</math> of any compare value equals <math>\frac{XLEN}{2} - 1:0</math> of <code>tdata2</code> after <math>XLEN - 1:\frac{XLEN}{2}</math> of the compare value is ANDed with <math>XLEN - 1:\frac{XLEN}{2}</math> of <code>tdata2</code>.</p> <p>8 (not equal): Matches when <code>match</code> = 0 would not match.</p> <p>9 (not napot): Matches when <code>match</code> = 1 would not match.</p> <p>12 (not mask low): Matches when <code>match</code> = 4 would not match.</p> <p>13 (not mask high): Matches when <code>match</code> = 5 would not match.</p> <p>Other values are reserved for future use.</p> <p>All comparisons only look at the lower <math>XLEN</math> (in the current mode) bits of the compare values and of <code>tdata2</code>. When <code>select</code> = 1 and access size is <math>N</math>, this is further reduced, and comparisons only look at the lower <math>N</math> bits of the compare values and of <code>tdata2</code>.</p>	WARL	0
m	When set, enable this trigger in M-mode.	WARL	0
s	When set, enable this trigger in S/HS-mode. This bit is hard-wired to 0 if the hart does not support S-mode.	WARL	0
u	When set, enable this trigger in U-mode. This bit is hard-wired to 0 if the hart does not support U-mode.	WARL	0

Continued on next page

Field	Description	Access	Reset
execute	When set, the trigger fires on the virtual address or opcode of an instruction that is executed.	WARL	0
store	When set, the trigger fires on the virtual address or data of any store.	WARL	0
load	When set, the trigger fires on the virtual address or data of any load.	WARL	0

### 5.7.12 Match Control Type 6 (mcontrol6, at 0x7a1)

This register is accessible as `tdata1` when `type` is 6.

Implementing this trigger as described here requires that `version` is 1 or higher, which in turn means `tinfo` must be implemented.

This replaces `mcontrol` in newer implementations and serves to provide additional functionality.

Address and data trigger implementation are heavily dependent on how the processor core is implemented. To accommodate various implementations, `execute`, `load`, and `store` address/data triggers may fire at whatever point in time is most convenient for the implementation.

Table 5.10 suggests timings for the best user experience. The underlying principle is that firing just before the instruction gives a user more insight, so is preferable. However, depending on the instruction and conditions, it might not be possible to evaluate the trigger until the instruction has partially executed. In that case it is better to let the instruction retire before the trigger fires, to avoid extra memory accesses which might affect the state of the system.

Table 5.10: Suggested Trigger Timings

Match Type	Suggested Trigger Timing
Execute Address	Before
Execute Instruction	Before
Execute Address+Instruction	Before
Load Address	Before
Load Data	After
Load Address+Data	After
Store Address	Before
Store Data	Before
Store Address+Data	Before

A chain of triggers must only fire if every trigger in the chain was matched by the same instruction.

This trigger type may be limited to address comparisons (`select` is always 0) only. If that is the case and masking is not supported (match values 4, 5, 12, 13), then `tdata2` must be able to hold all valid virtual addresses but it need not be capable of holding other values.

The Privileged Spec says that breakpoint exceptions that occur on instruction fetches, loads, or stores update the `tval` CSR with either zero or the faulting virtual address. The faulting virtual

address for an mcontrol6 trigger with `action = 0` is the address being accessed and which caused that trigger to fire. If multiple mcontrol6 triggers are chained then the faulting virtual address is the address which caused any of the chained triggers to fire.

In implementations that support `match` mode 1 (NAPOT), not all NAPOT ranges may be supported. All NAPOT ranges between  $2^1$  and  $2^{\text{maskmax6}}$  are supported where  $\text{maskmax6} \geq 1$ . The value of `maskmax6` can be determined by the debugger via the following sequence:

1. Set `match = 1`.
2. Read `match`. If it is not 1 then NAPOT matching is not supported.
3. Write all ones to `tdata2`.
4. Read `tdata2`. The value of `maskmax6` is the index of the most significant 0 bit plus 1.

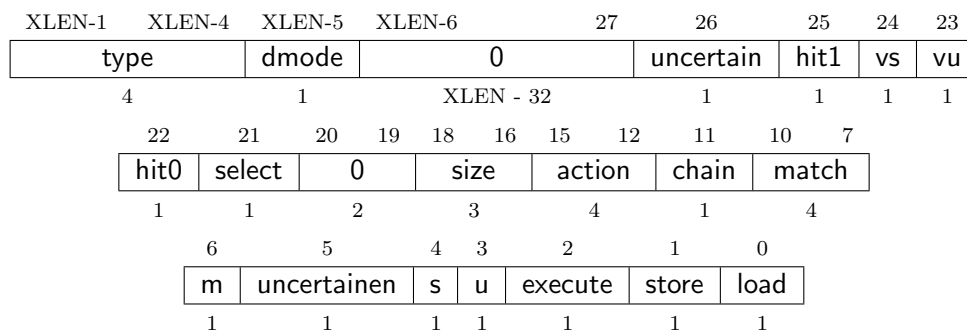
If `textra32` or `textra64` are implemented for this trigger, it only matches when the conditions set there are satisfied.

---

`uncertain` and `uncertainen` exist to accommodate systems where not every memory access is fully observed by the Trigger Module. Possible examples include data values in far AMOs, and the address/data/size of accesses by instructions that perform multiple memory accesses, such as vector, push, and pop instructions.

While the `uncertain` mechanism exists to deal with these situations, it can lead to an unusable number of false positives. Users will get a much better debug experience if the TM does have perfect visibility into the details of every memory access.

This CSR is read/write.



Field	Description	Access	Reset
<code>uncertain</code>	<p>This bit should be updated every time that <code>hit0</code> or <code>hit1</code> is updated.</p> <p>0 (certain): The trigger that fired satisfied the configured conditions, or this bit is not implemented.</p> <p>1 (uncertain): The trigger that fired might not have perfectly satisfied the configured conditions. Due to the implementation the hardware cannot be certain.</p>	WARL	0

*Continued on next page*

Field	Description	Access	Reset
vs	When set, enable this trigger in VS-mode. This bit is hard-wired to 0 if the hart does not support virtualization mode.	WARL	0
vu	When set, enable this trigger in VU-mode. This bit is hard-wired to 0 if the hart does not support virtualization mode.	WARL	0
hit0	<p>If they are implemented, hit1 (MSB) and hit0 (LSB) combine into a single 2-bit field. The TM updates this field when the trigger fires. After the debugger has seen the update, it will normally write 0 to this field to so it can see future changes. If either of the bits is not implemented, the unimplemented bits will be read-only 0.</p> <p>0 (false): The trigger did not fire.</p> <p>1 (before): The trigger fired before the instruction that matched it was retired, but after all preceding instructions are retired. This explicitly allows for instructions to be partially executed, as described in Section 5.6.</p> <p>xepc or dpc (depending on action) must be set to the virtual address of the instruction that matched.</p> <p>2 (after): The trigger fired after the instruction that triggered and at least one additional instruction were retired. xepc or dpc (depending on action) must be set to the virtual address of the next instruction that must be executed to preserve the program flow.</p> <p>3 (immediately after): The trigger fired just after the instruction that triggered it was retired, but before any subsequent instructions were executed. xepc or dpc (depending on action) must be set to the virtual address of the next instruction that must be executed to preserve the program flow.</p> <p>If the instruction performed multiple memory accesses, all of them have been completed.</p>	WARL	0

*Continued on next page*



Field	Description	Access	Reset
select	<p>This bit determines the contents of the XLEN-bit compare values.</p> <p>0 (address): There is at least one compare value and it contains the lowest virtual address of the access. In addition, it is recommended that there are additional compare values for the other accessed virtual addresses match. (E.g. on a 32-bit read from 0x4000, the lowest address is 0x4000 and the other addresses are 0x4001, 0x4002, and 0x4003.)</p> <p>1 (data): There is exactly one compare value and it contains the data value loaded or stored, or the instruction executed. Any bits beyond the size of the data access will contain 0.</p>	WARL	0

*Continued on next page*

Field	Description	Access	Reset
size	<p>0 (any): The trigger will attempt to match against an access of any size. The behavior is only well-defined if <code>select = 0</code>, or if the access size is <code>XLEN</code>.</p> <p>1 (8bit): The trigger will only match against 8-bit memory accesses.</p> <p>2 (16bit): The trigger will only match against 16-bit memory accesses or execution of 16-bit instructions.</p> <p>3 (32bit): The trigger will only match against 32-bit memory accesses or execution of 32-bit instructions.</p> <p>4 (48bit): The trigger will only match against execution of 48-bit instructions.</p> <p>5 (64bit): The trigger will only match against 64-bit memory accesses or execution of 64-bit instructions.</p> <p>6 (128bit): The trigger will only match against 128-bit memory accesses or execution of 128-bit instructions.</p> <p>An implementation must support the value of 0, but all other values are optional. When an implementation supports address triggers (<code>select = 0</code>), it is recommended that those triggers support every access size that the hart supports, as well as for every instruction size that the hart supports. Implementations such as RV32D or RV64V are able to perform loads and stores that are wider than <code>XLEN</code>. Custom extensions may also support instructions that are wider than <code>XLEN</code>. Because <code>tdata2</code> is of size <code>XLEN</code>, there is a known limitation that data value triggers (<code>select = 1</code>) can only be supported for access sizes up to <code>XLEN</code> bits. When an implementation supports data value triggers (<code>select = 1</code>), it is recommended that those triggers support every access size up to <code>XLEN</code> that the hart supports, as well as for every instruction length up to <code>XLEN</code> that the hart supports.</p>	WARL	0
action	The action to take when the trigger fires. The values are explained in Table 5.1.	WARL	0

*Continued on next page*

Field	Description	Access	Reset
chain	<p>0 (disabled): When this trigger matches, the configured action is taken.</p> <p>1 (enabled): While this trigger does not match, it prevents the trigger with the next index from matching.</p> <p>A trigger chain starts on the first trigger with <code>chain = 1</code> after a trigger with <code>chain = 0</code>, or simply on the first trigger if that has <code>chain = 1</code>. It ends on the first trigger after that which has <code>chain = 0</code>. This final trigger is part of the chain. The action on all but the final trigger is ignored. The action on that final trigger will be taken if and only if all the triggers in the chain match at the same time. Debuggers should not terminate a chain with a trigger with a different type. It is undefined when exactly such a chain fires.</p> <p>Because <code>chain</code> affects the next trigger, hardware must zero it in writes to <code>mcontrol6</code> that set <code>dmode</code> to 0 if the next trigger has <code>dmode</code> of 1. In addition hardware should ignore writes to <code>mcontrol6</code> that set <code>dmode</code> to 1 if the previous trigger has both <code>dmode</code> of 0 and <code>chain</code> of 1. Debuggers must avoid the latter case by checking <code>chain</code> on the previous trigger if they're writing <code>mcontrol6</code>.</p> <p>Implementations that wish to limit the maximum length of a trigger chain (eg. to meet timing requirements) may do so by zeroing <code>chain</code> in writes to <code>mcontrol6</code> that would make the chain too long.</p>	WARL	0

*Continued on next page*

Field	Description	Access	Reset
match	<p>0 (equal): Matches when any compare value equals <code>tdata2</code>.</p> <p>1 (napot): Matches when the top <math>M</math> bits of any compare value match the top <math>M</math> bits of <code>tdata2</code>. <math>M</math> is <math>XLEN - 1</math> minus the index of the least-significant bit containing 0 in <code>tdata2</code>. <code>tdata2</code> is WARL and if bits <code>maskmax6 - 1:0</code> are written with all ones then bit <code>maskmax6 - 1</code> will be set to 0 while the values of bits <code>maskmax6 - 2:0</code> are UNSPECIFIED. Legal values for <code>tdata2</code> require <math>M + \text{maskmax6} \geq XLEN</math> and <math>M &gt; 0</math>. See above for how to determine <code>maskmax6</code>.</p> <p>2 (ge): Matches when any compare value is greater than (unsigned) or equal to <code>tdata2</code>.</p> <p>3 (lt): Matches when any compare value is less than (unsigned) <code>tdata2</code>.</p> <p>4 (mask low): Matches when <math>\frac{XLEN}{2} - 1:0</math> of any compare value equals <math>\frac{XLEN}{2} - 1:0</math> of <code>tdata2</code> after <math>\frac{XLEN}{2} - 1:0</math> of the compare value is ANDed with <math>XLEN - 1:\frac{XLEN}{2}</math> of <code>tdata2</code>.</p> <p>5 (mask high): Matches when <math>XLEN - 1:\frac{XLEN}{2}</math> of any compare value equals <math>\frac{XLEN}{2} - 1:0</math> of <code>tdata2</code> after <math>XLEN - 1:\frac{XLEN}{2}</math> of the compare value is ANDed with <math>XLEN - 1:\frac{XLEN}{2}</math> of <code>tdata2</code>.</p> <p>8 (not equal): Matches when <code>match</code> = 0 would not match.</p> <p>9 (not napot): Matches when <code>match</code> = 1 would not match.</p> <p>12 (not mask low): Matches when <code>match</code> = 4 would not match.</p> <p>13 (not mask high): Matches when <code>match</code> = 5 would not match.</p> <p>Other values are reserved for future use.</p> <p>All comparisons only look at the lower <math>XLEN</math> (in the current mode) bits of the compare values and of <code>tdata2</code>. When <code>select</code> =1 and access size is <math>N</math>, this is further reduced, and comparisons only look at the lower <math>N</math> bits of the compare values and of <code>tdata2</code>.</p>	WARL	0
m	When set, enable this trigger in M-mode.	WARL	0

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Field	Description	Access	Reset
uncertainen	0 (disabled): This trigger will only match if the hardware can perfectly evaluate it. 1 (enabled): This trigger will match if it's possible that it would match if the Trigger Module had perfect information about the operations being performed.	WARL	0
s	When set, enable this trigger in S/HS-mode. This bit is hard-wired to 0 if the hart does not support S-mode.	WARL	0
u	When set, enable this trigger in U-mode. This bit is hard-wired to 0 if the hart does not support U-mode.	WARL	0
execute	When set, the trigger fires on the virtual address or opcode of an instruction that is executed.	WARL	0
store	When set, the trigger fires on the virtual address or data of any store.	WARL	0
load	When set, the trigger fires on the virtual address or data of any load.	WARL	0

### 5.7.13 Instruction Count (icount, at 0x7a1)

This register is accessible as `tdata1` when `type` is 3.

This trigger matches when:

1. An instruction retires after having been fetched in a privilege mode where the trigger is enabled. This explicitly includes `xRET` instructions.
2. A trap is taken from a privilege mode where the trigger is enabled. This explicitly includes traps taken due to interrupts.

If more than one of the above events occur during a single instruction execution, the trigger still only matches once for that instruction.

---

*For use in single step, icount must match for traps where the instruction will not be reexecuted after the handler, such as illegal instructions that are emulated by privileged software and the instruction being emulated never retires. Ideally, icount would not match for traps where the instruction will later be retried by the handler, such as page faults where privileged software modifies the page tables and returns to the faulting instruction which ultimately retires. Trying to distinguish the two cases leads to complex rules, so instead the rule is simply that all traps match. See also Section 4.5.2.*

When `count` is greater than 1 and the trigger matches, then `count` is decremented by 1.

When `count` is 1 and the trigger matches, then `pending` becomes set. In addition `count` will become 0 unless it is hard-wired to 1.

The only exception to the above is when the instruction the trigger matched on is a write to the

icount trigger. In that case `pending` might or might not become set if `count` was 1. Afterwards `count` contains the newly written value.

When `count` is 0 it stays at 0 until explicitly written.

When `pending` is set, the trigger fires just before any further instructions are executed in a mode where the trigger is enabled. As the trigger fires, `pending` is cleared. In addition, if `count` is hard-wired to 1 then `m`, `s`, `u`, `vs`, and `vu` are all cleared.

If the trigger fires with `action` =0 then zero is written to the `tval` CSR on the breakpoint trap.

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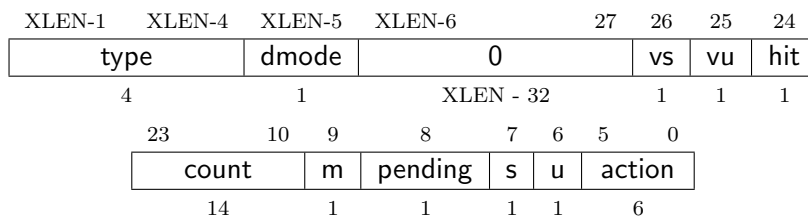
*The intent of `pending` is to cleanly handle the case where `action` is 0, `m` is 0, `u` is 1, `count` is 1, and the U-mode instruction being executed causes a trap into M-mode. In that case we want the entire M-mode handler to be executed, and the debug trap to be taken before the next U-mode instruction.*

---

*This trigger type is intended to be used as a single step for software monitor programs or native debug. Systems that support multiple privilege modes that want to debug software running in lower privilege modes don't need to support `count` greater than 1.*

If `textra32` or `textra64` are implemented for this trigger, it only matches when the conditions set there are satisfied.

This CSR is read/write.



Field	Description	Access	Reset
<code>vs</code>	When set, enable this trigger in VS-mode. This bit is hard-wired to 0 if the hart does not support virtualization mode.	WARL	0
<code>vu</code>	When set, enable this trigger in VU-mode. This bit is hard-wired to 0 if the hart does not support virtualization mode.	WARL	0
<code>hit</code>	If this bit is implemented, the hardware sets it when this trigger fires. The trigger's user can set or clear it at any time. It is used to determine which trigger(s) fires. If the bit is not implemented, it is always 0 and writing it has no effect.	WARL	0

*Continued on next page*

Field	Description	Access	Reset
count	The trigger will generally fire after <code>count</code> instructions in enabled modes have been executed. See above for the precise behavior.	WARL	1
m	When set, enable this trigger in M-mode.	WARL	0
pending	This bit becomes set when <code>count</code> is decremented from 1 to 0. It is cleared when the trigger fires, which will happen just before executing the next instruction in one of the enabled modes.	R/W	0
s	When set, enable this trigger in S/HS-mode. This bit is hard-wired to 0 if the hart does not support S-mode.	WARL	0
u	When set, enable this trigger in U-mode. This bit is hard-wired to 0 if the hart does not support U-mode.	WARL	0
action	The action to take when the trigger fires. The values are explained in Table 5.1.	WARL	0

#### 5.7.14 Interrupt Trigger (`ittrigger`, at `0x7a1`)

This register is accessible as `tdata1` when `type` is 4.

This trigger can fire when an interrupt trap is taken.

It can be enabled for individual interrupt numbers by setting the bit corresponding to the interrupt number in `tdata2`. The interrupt number is interpreted in the mode that the trap handler executes in. (E.g. virtualized interrupt numbers are not the same in every mode.) In addition the trigger can be enabled for non-maskable interrupts using `nmi`.

Hardware may only support a subset of interrupts for this trigger. A debugger must read back `tdata2` after writing it to confirm the requested functionality is actually supported.

When the trigger matches, it fires after the trap occurs, just before the first instruction of the trap handler is executed. If `action` = 0, the standard CSRs are updated for taking the breakpoint trap, and zero is written to the relevant `tval` CSR. If the breakpoint trap does not go to a higher privilege mode, this will lose CSR information for the original trap. See Section 5.4 for more information about this case.

If `textra32` or `textra64` are implemented for this trigger, it only matches when the conditions set there are satisfied.

This CSR is read/write.

XLEN-1	XLEN-4	XLEN-5	XLEN-6	XLEN-7	13	12	11
type	dmode	hit	0			vs	vu
4	1	1	XLEN - 19			1	1

10	9	8	7	6	5	0
nmi	m	0	s	u	action	
1	1	1	1	1	6	

Field	Description	Access	Reset
hit	If this bit is implemented, the hardware sets it when this trigger matches. The trigger's user can set or clear it at any time. It is used to determine which trigger(s) matched. If the bit is not implemented, it is always 0 and writing it has no effect.	WARL	0
vs	When set, enable this trigger for interrupts that are taken from VS mode. This bit is hard-wired to 0 if the hart does not support virtualization mode.	WARL	0
vu	When set, enable this trigger for interrupts that are taken from VU mode. This bit is hard-wired to 0 if the hart does not support virtualization mode.	WARL	0
nmi	When set, non-maskable interrupts cause this trigger to fire if the trigger is enabled for the current mode.	WARL	0
m	When set, enable this trigger for interrupts that are taken from M mode.	WARL	0
s	When set, enable this trigger for interrupts that are taken from S/HS mode. This bit is hard-wired to 0 if the hart does not support S-mode.	WARL	0
u	When set, enable this trigger for interrupts that are taken from U mode. This bit is hard-wired to 0 if the hart does not support U-mode.	WARL	0
action	The action to take when the trigger fires. The values are explained in Table 5.1.	WARL	0

### 5.7.15 Exception Trigger (etrigger, at 0x7a1)

This register is accessible as `tdata1` when `type` is 5.

This trigger may fire on up to XLEN of the Exception Codes defined in `mcause` (described in the Privileged Spec, with Interrupt=0). Those causes are configured by writing the corresponding bit in `tdata2`. (E.g. to trap on an illegal instruction, the debugger sets bit 2 in `tdata2`.)

Hardware may support only a subset of exceptions. A debugger must read back `tdata2` after writing it to confirm the requested functionality is actually supported.

When the trigger matches, it fires after the trap occurs, just before the first instruction of the trap handler is executed. If `action` =0, the standard CSRs are updated for taking the breakpoint trap,



and zero is written to the relevant `tval` CSR. If the breakpoint trap does not go to a higher privilege mode, this will lose CSR information for the original trap. See Section 5.4 for more information about this case.

If `textra32` or `textra64` are implemented for this trigger, it only matches when the conditions set there are satisfied.

This CSR is read/write.

XLEN-1	XLEN-4	XLEN-5	XLEN-6	XLEN-7	13	12	11	10	9	8	7	6	5	0
type	dmode	hit	0	vs	vu	0	m	0	s	u	action			
4	1	1	XLEN - 19	1	1	1	1	1	1	1	1	1	6	

Field	Description	Access	Reset
hit	If this bit is implemented, the hardware sets it when this trigger matches. The trigger's user can set or clear it at any time. It is used to determine which trigger(s) matched. If the bit is not implemented, it is always 0 and writing it has no effect.	WARL	0
vs	When set, enable this trigger for exceptions that are taken from VS mode. This bit is hard-wired to 0 if the hart does not support virtualization mode.	WARL	0
vu	When set, enable this trigger for exceptions that are taken from VU mode. This bit is hard-wired to 0 if the hart does not support virtualization mode.	WARL	0
m	When set, enable this trigger for exceptions that are taken from M mode.	WARL	0
s	When set, enable this trigger for exceptions that are taken from S/HS mode. This bit is hard-wired to 0 if the hart does not support S-mode.	WARL	0
u	When set, enable this trigger for exceptions that are taken from U mode. This bit is hard-wired to 0 if the hart does not support U-mode.	WARL	0
action	The action to take when the trigger fires. The values are explained in Table 5.1.	WARL	0

### 5.7.16 External Trigger (`tmexttrigger`, at `0x7a1`)

This register is accessible as `tdata1` when `type` is 7.

This trigger fires when any selected TM external trigger input signals. Up to 16 TM external trigger inputs coming from other blocks outside the TM, (e.g. signaling an hpmcounter overflow) can be selected. Hardware may support none or just a few TM external trigger inputs (starting

with TM external trigger input 0 and continuing sequentially). Unsupported inputs are hardwired to be inactive.

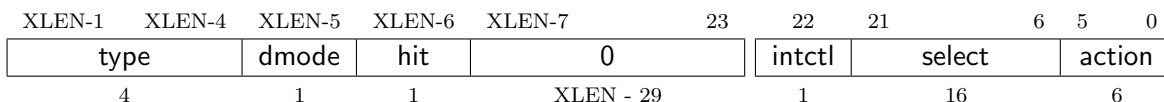
If the trigger fires with `action` =0 then zero is written to the `tval` CSR on the breakpoint trap. This trigger fires asynchronously but it is subject to delegation by `medeleg`[3] like the other triggers.

The TM external trigger input can signal when the trigger is prevented from firing due to one of the mechanisms in section 5.4. An implementation may either ignore the signal altogether when it cannot fire (dropping the trigger event) or it may hold the action as pending and fire the trigger once it is legal to do so.

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*`intctl` is intended to be used by the `clicintrig` mechanism from the Core-Local Interrupt Controller (CLIC) RISC-V Privileged Architecture Extensions.*

This CSR is read/write.



Field	Description	Access	Reset
hit	If this bit is implemented, the hardware sets it when this trigger matches. The trigger's user can set or clear it at any time. It is used to determine which trigger(s) matched. If the bit is not implemented, it is always 0 and writing it has no effect.	WARL	0
intctl	This optional bit, when set, causes this trigger to fire whenever an attached interrupt controller signals a trigger.	WARL	0
select	Selects any combination of up to 16 TM external trigger inputs that cause this trigger to fire.	WARL	0
action	The action to take when the trigger fires. The values are explained in Table 5.1.	WARL	0

### 5.7.17 Trigger Extra (RV32) (`textra32`, at `0x7a3`)

This register is accessible as `tdata3` when `type` is 2, 3, 4, 5, or 6 and `XLEN`=32.

All functionality in this register is optional. The `value` bits may tie any number of upper bits to 0. The `select` bits may only support 0 (ignore).

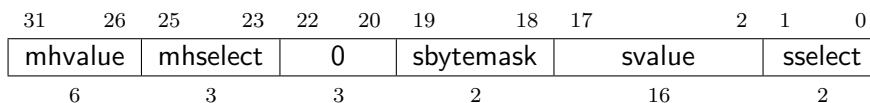
Byte-granular comparison of `scontext` to `svalue` allows `scontext` to be defined to include more than one element of comparison. For example, software instrumentation can program the `scontext` value to be the concatenation of different ID contexts such as process ID and thread ID. The user can then program byte compares based on `sbytemask` to include one or more of the contexts in the compare.

Byte masking only applies to `scontext` comparison; i.e when `sselect` is 1.

---

*Note that `sselect` and `mhselect` filtering apply in all modes, including M-mode and S-mode. If desired, debuggers can use a trigger's mode filtering bits to restrict the matching to modes where it considers ASID/VMID/`scontext`/`hcontext` to be active.*

This CSR is read/write.



Field	Description	Access	Reset
mhvalue	Data used together with <code>mhselect</code> .	WARL	0
mhselect	0 (ignore): Ignore <code>mhvalue</code> . 4 (mcontext): This trigger will only match or fire if the low bits of <code>mcontext</code> / <code>hcontext</code> equal <code>mhvalue</code> . 1, 5 (mcontext_select): This trigger will only match or fire if the low bits of <code>mcontext</code> / <code>hcontext</code> equal { <code>mhvalue</code> , <code>mhselect</code> [2]}. 2, 6 (vmid_select): This trigger will only match or fire if VMID in <code>hcatp</code> equals the lower VMID-MAX (defined in the Privileged Spec) bits of { <code>mhvalue</code> , <code>mhselect</code> [2]}. 3, 7 (reserved): Reserved. If the H extension is not supported, the only legal values are 0 and 4.	WARL	0
sbytemask	When the least significant bit of this field is 1, it causes bits 7:0 in the comparison to be ignored, when <code>sselect</code> =1. When the next most significant bit of this field is 1, it causes bits 15:8 to be ignored in the comparison, when <code>sselect</code> =1.	WARL	0
svalue	Data used together with <code>sselect</code> . This field should be tied to 0 when S-mode is not supported.	WARL	0

*Continued on next page*

Field	Description	Access	Reset
sselect	0 (ignore): Ignore <a href="#">svalue</a> . 1 (scontext): This trigger will only match or fire if the low bits of <a href="#">scontext</a> equal <a href="#">svalue</a> . 2 (asid): This trigger will only match or fire if: <ul style="list-style-type: none"> <li>the mode is VS-mode or VU-mode and ASID in <a href="#">vsatp</a> equals the lower ASID-MAX (defined in the Privileged Spec) bits of <a href="#">svalue</a>.</li> <li>in all other modes, ASID in <a href="#">satp</a> equals the lower ASIDMAX (defined in the Privileged Spec) bits of <a href="#">svalue</a>.</li> </ul> This field should be tied to 0 when S-mode is not supported.	WARL	0

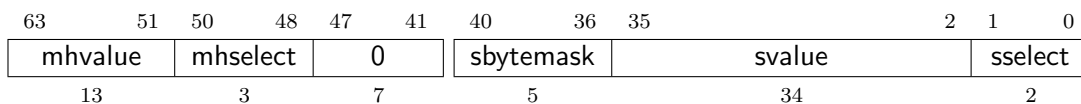
### 5.7.18 Trigger Extra (RV64) ([textra64](#), at [0x7a3](#))

This register is accessible as [tdata3](#) when [type](#) is 2, 3, 4, 5, or 6 and XLEN=64. The fields are defined above, in [textra32](#).

Byte-granular comparison of [scontext](#) to [svalue](#) in [textra64](#) allows [scontext](#) to be defined to include more than one element of comparison. For example, software instrumentation can program the [scontext](#) value to be the concatenation of different ID contexts such as process ID and thread ID. The user can then program byte compares based on [sbytemask](#) to include one or more of the contexts in the compare.

Byte masking only applies to [scontext](#) comparison; i.e when [sselect](#) is 1.

This CSR is read/write.



Field	Description	Access	Reset
sbytemask	When the least significant bit of this field is 1, it causes bits 7:0 in the comparison to be ignored, when <a href="#">sselect</a> =1. Likewise, the second bit controls the comparison of bits 15:8, third bit controls the comparison of bits 23:16, fourth bit controls the comparison of bits 31:24, and fifth bit controls the comparison of bits 33:32.	WARL	0

## Chapter 6

# Debug Transport Module (DTM) (non-ISA extension)

Debug Transport Modules provide access to the DM over one or more transports (e.g. JTAG or USB).

There may be multiple DTMs in a single hardware platform. Ideally every component that communicates with the outside world includes a DTM, allowing a hardware platform to be debugged through every transport it supports. For instance a USB component could include a DTM. This would trivially allow any hardware platform to be debugged over USB. All that is required is that the USB module already in use also has access to the Debug Module Interface.

Using multiple DTMs at the same time is not supported. It is left to the user to ensure this does not happen.

This specification defines a JTAG DTM in Section 6.1. Additional DTMs may be added in future versions of this specification.

An implementation can be compatible with this specification without implementing any of this section. In that case it must be advertised as conforming to “RISC-V Debug Specification 1.0-STABLE, with custom DTM.” If the JTAG DTM described here is implemented, it must be advertised as conforming to the “RISC-V Debug Specification 1.0-STABLE, with JTAG DTM.”

### 6.1 JTAG Debug Transport Module

This Debug Transport Module is based around a normal JTAG Test Access Port (TAP). The JTAG TAP allows access to arbitrary JTAG registers by first selecting one using the JTAG instruction register (IR), and then accessing it through the JTAG data register (DR).

### 6.1.1 JTAG Background

JTAG refers to IEEE Std 1149.1-2013. It is a standard that defines test logic that can be included in an integrated circuit to test the interconnections between integrated circuits, test the integrated circuit itself, and observe or modify circuit activity during the component's normal operation. This specification uses the latter functionality. The JTAG standard defines a Test Access Port (TAP) that can be used to read and write a few custom registers, which can be used to communicate with debug hardware in a component.

### 6.1.2 JTAG DTM Registers

JTAG TAPs used as a DTM must have an IR of at least 5 bits. When the TAP is reset, IR must default to 00001, selecting the IDCODE instruction. A full list of JTAG registers along with their encoding is in Table 6.1. If the IR actually has more than 5 bits, then the encodings in Table 6.1 should be extended with 0's in their most significant bits, except for the 0x1f encoding of BYPASS, which must be extended with 1's in the most significant bits. The only regular JTAG registers a debugger might use are BYPASS and IDCODE, but this specification leaves IR space for many other standard JTAG instructions. Unimplemented instructions must select the BYPASS register.

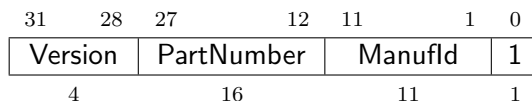
Table 6.1: JTAG DTM TAP Registers

Address	Name	Description	Page
0x00	BYPASS	JTAG recommends this encoding	
0x01	IDCODE	To identify a specific silicon version	
0x10	DTM Control and Status ( <i>dtmcs</i> )	For Debugging	97
0x11	Debug Module Interface Access ( <i>dmi</i> )	For Debugging	98
0x12	Reserved (BYPASS)	Reserved for future RISC-V debugging	
0x13	Reserved (BYPASS)	Reserved for future RISC-V debugging	
0x14	Reserved (BYPASS)	Reserved for future RISC-V debugging	
0x15	Reserved (BYPASS)	Reserved for future RISC-V standards	
0x16	Reserved (BYPASS)	Reserved for future RISC-V standards	
0x17	Reserved (BYPASS)	Reserved for future RISC-V standards	
0x1f	BYPASS	JTAG requires this encoding	

### 6.1.3 IDCODE (at 0x01)

This register is selected (in IR) when the TAP state machine is reset. Its definition is exactly as defined in IEEE Std 1149.1-2013.

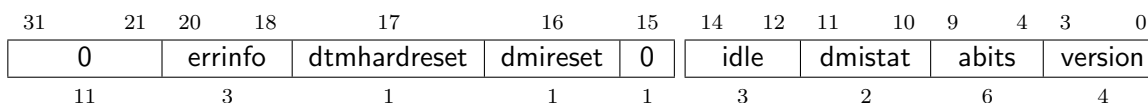
This entire register is read-only.



Field	Description	Access	Reset
Version	Identifies the release version of this part.	R	Preset
PartNumber	Identifies the designer's part number of this part.	R	Preset
Manufld	Identifies the designer/manufacturer of this part. Bits 6:0 must be bits 6:0 of the designer/manufacturer's Identification Code as assigned by JEDEC Standard JEP106. Bits 10:7 contain the modulo-16 count of the number of continuation characters (0x7f) in that same Identification Code.	R	Preset

#### 6.1.4 DTM Control and Status (dtmcs, at 0x10)

The size of this register will remain constant in future versions so that a debugger can always determine the version of the DTM.



Field	Description	Access	Reset
errinfo	<p>This optional field may provide additional detail about an error that occurred when communicating with a DM. It is updated whenever <code>op</code> is updated by the hardware or when 1 is written to <code>dmireset</code>.</p> <p>0 (not implemented): This field is not implemented.</p> <p>1 (dmi error): There was an error between the DTM and DMI.</p> <p>2 (communication error): There was an error between the DMI and a DMI subordinate.</p> <p>3 (device error): The DMI subordinate reported an error.</p> <p>4 (unknown): There is no error to report, or no further information available about the error. This is the reset value if the field is implemented. Other values are reserved for future use by this specification.</p>	R	4

*Continued on next page*

Field	Description	Access	Reset
dtmhardreset	Writing 1 to this bit does a hard reset of the DTM, causing the DTM to forget about any outstanding DMI transactions, and returning all registers and internal state to their reset value. In general this should only be used when the Debugger has reason to expect that the outstanding DMI transaction will never complete (e.g. a reset condition caused an inflight DMI transaction to be cancelled).	W1	-
dmireset	Writing 1 to this bit clears the sticky error state and resets <code>errinfo</code> , but does not affect outstanding DMI transactions.	W1	-
idle	This is a hint to the debugger of the minimum number of cycles a debugger should spend in Run-Test/Idle after every DMI scan to avoid a ‘busy’ return code ( <code>dmistat</code> of 3). A debugger must still check <code>dmistat</code> when necessary. 0: It is not necessary to enter Run-Test/Idle at all. 1: Enter Run-Test/Idle and leave it immediately. 2: Enter Run-Test/Idle and stay there for 1 cycle before leaving. And so on.	R	Preset
dmistat	Read-only alias of <code>op</code> .	R	0
abits	The size of <code>address</code> in <code>dmi</code> .	R	Preset
version	0 (0.11): Version described in spec version 0.11. 1 (1.0): Version described in spec versions 0.13 and 1.0. 15 (custom): Version not described in any available version of this spec.	R	1

### 6.1.5 Debug Module Interface Access (`dmi`, at `0x11`)

This register allows access to the Debug Module Interface (DMI).

In Update-DR, the DTM starts the operation specified in `op` unless the current status reported in `op` is sticky.

In Capture-DR, the DTM updates `data` with the result from that operation, updating `op` if the current `op` isn’t sticky.

See Section [B.2.1](#) for examples of how this is used.

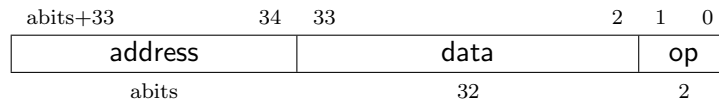
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*The still-in-progress status is sticky to accommodate debuggers that batch together a number of scans, which must all be executed or stop as soon as there’s a problem.*

*For instance a series of scans may write a Debug Program and execute it. If one of the writes*



fails but the execution continues, then the Debug Program may hang or have other unexpected side effects.



Field	Description	Access	Reset
<b>address</b>	Address used for DMI access. In Update-DR this value is used to access the DM over the DMI. <b>op</b> defines what this register contains after every possible operation.	R/W	0
<b>data</b>	The data to send to the DM over the DMI during Update-DR, and the data returned from the DM as a result of the previous operation.	R/W	0

*Continued on next page*

Field	Description	Access	Reset
op	<p>When the debugger writes this field, it has the following meaning:</p> <p>0 (nop): Ignore <code>data</code> and <code>address</code>. Don't send anything over the DMI during Update-DR. This operation should never result in a busy or error response. The address and data reported in the following Capture-DR are undefined. This operation leaves the values in <code>address</code> and <code>data</code> UNSPECIFIED.</p> <p>1 (read): Read from <code>address</code>. When this operation succeeds, <code>address</code> contains the address that was read from, and <code>data</code> contains the value that was read.</p> <p>2 (write): Write <code>data</code> to <code>address</code>. This operation leaves the values in <code>address</code> and <code>data</code> UNSPECIFIED.</p> <p>3 (reserved): Reserved.</p> <p>When the debugger reads this field, it means the following:</p> <p>0 (success): The previous operation completed successfully.</p> <p>1 (reserved): Reserved.</p> <p>2 (failed): A previous operation failed. The data scanned into <code>dmi</code> in this access will be ignored. This status is sticky and can be cleared by writing <code>dmireset</code> in <code>dtmcs</code>. This indicates that the DM itself or the DMI responded with an error. There are no specified cases in which the DM would respond with an error, and DMI is not required to support returning errors. If a debugger sees this status, there might be additional information in <code>errinfo</code>.</p> <p>3 (busy): An operation was attempted while a DMI request is still in progress. The data scanned into <code>dmi</code> in this access will be ignored. This status is sticky and can be cleared by writing <code>dmireset</code> in <code>dtmcs</code>. If a debugger sees this status, it needs to give the target more TCK edges between Update-DR and Capture-DR. The simplest way to do that is to add extra transitions in Run-Test/Idle.</p>	R/W	0

### 6.1.6 BYPASS (at 0x1f)

1-bit register that has no effect. It is used when a debugger does not want to communicate with this TAP.

This entire register is read-only.

0
0
1

### 6.1.7 JTAG Connector

#### 6.1.7.1 Recommended JTAG Connector

To make it easy to acquire debug hardware, this spec recommends a connector that is compatible with the MIPI-10 .05 inch connector specification, as described in MIPI Debug & Trace Connector Recommendations, Version 1.20, 2 July 2021.

The connector has .05 inch spacing, gold-plated male header with .016 inch thick hardened copper or beryllium bronze square posts (SAMTEC FTSH or equivalent). Female connectors are compatible 20 $\mu$ m gold connectors.

Viewing the male header from above (the pins pointing at your eye), a target's connector looks as it does in Table 6.5. The function of each pin is described in Table 6.6.

Table 6.5: MIPI 10-pin JTAG + nRESET Connector Diagram

VREF DEBUG	1	2	TMS
GND	3	4	TCK
GND	5	6	TDO
GND or KEY	7	8	TDI
GND	9	10	nRESET

If a hardware platform requires nTRST then it is permissible to reuse the nRESET pin as the nTRST signal, resulting in a MIPI 10-pin JTAG + nTRST connector.

#### 6.1.7.2 Alternate JTAG Connector

The MIPI-10 connector should provide plenty of signals for all modern hardware. If a design does need legacy JTAG signals, then the MIPI-20 connector should be used. Pins whose functionality isn't needed may be left unconnected.

Its physical connector is virtually identical to MIPI-10, except that it's twice as long, supporting twice as many pins. Its pinout is shown in Table 6.7. The function of each pin is described in Table 6.6.

Table 6.6: JTAG Connector Pin Functions

Essential	GND	Connected to ground.
	TCK	JTAG TCK signal, driven by the debug adapter.
	TDI	JTAG TDI signal, driven by the debug adapter.
	TDO	JTAG TDO signal, driven by the target.
	TMS	JTAG TMS signal, driven by the debug adapter.
	VREF DEBUG	Reference voltage for logic high.
Recommended	nRESET	Open drain active low reset signal, usually driven by the debug adapter. The signal may be used bi-directional to drive or sense the target reset signal. Asserting reset should reset any RISC-V cores as well as any other peripherals on the PCB. It should not reset the debug logic. This pin is optional but strongly encouraged.  nRESET should never be connected to the TAP reset, otherwise the debugger might not be able to debug through a reset to discover the cause of a crash or to maintain execution control after the reset.
	KEY	This pin may be cut on the male and plugged on the female header to ensure the header is always plugged in correctly. It is, however, recommended to use this pin as an additional ground, to allow for fastest TCK speeds. A shrouded connector should be used to prevent the cable from being plugged in incorrectly.
Advanced	EXT	Reserved for custom use. Could be an input or an output.
	TRIGIN	Not used by this specification, to be driven by debug adapter. (Can be used for extended functions like UART or boot mode selection by some debug adapters).
	TRIGOUT	Not used by this specification, driven by the target.
Specialized	nTRST	Test reset, driven by the debug adapter. Asserting nTRST initializes the JTAG DTM asynchronously. It is used in systems where the JTAG DTM is not ready to be used after a normal power up. This signal is sometimes called TRST*.
Legacy	RTCK	Return test clock, driven by the target. A target may relay the TCK signal here once it has processed it, allowing a debugger to adjust its TCK frequency in response.  This signal should only be used to support legacy components that rely on this functionality.
	nTRST_PD	Test reset pull-down, driven by the debug adapter. Same function as nTRST, but with pull-down resistor on target.  This signal should only be used to support legacy components that rely on this functionality.

Table 6.7: MIPI 20-pin JTAG Connector Diagram

VREF DEBUG	1	2	TMS
GND	3	4	TCK
GND	5	6	TDO
GND or KEY	7	8	TDI
GND	9	10	nRESET
GND	11	12	GND or RTCK
GND	13	14	NC or nTRST_PD
GND	15	16	nTRST or NC
GND	17	18	TRIGIN or NC
GND	19	20	TRIGOUT or GND

### 6.1.8 cJTAG

This spec does not have specific recommendations on how to use the cJTAG protocol.

When implementing cJTAG access to a JTAG DTM, the MIPI 10-pin Narrow JTAG connector should be used. Pins whose functionality isn't needed may be left unconnected.

Viewing the male header from above (the pins pointing at your eye), a target's connector looks as it does in Table 6.8.

Table 6.8: MIPI 10-pin Narrow JTAG Connector Diagram

VREF DEBUG	1	2	TMSC
GND	3	4	TCKC
GND	5	6	EXT or NC
GND or KEY	7	8	NC or nTRST_PD
GND	9	10	nRESET

# Appendix A

## Hardware Implementations

Below are two possible implementations. A designer could choose one, mix and match, or come up with their own design.

### A.1 Abstract Command Based

Halting happens by stalling the hart execution pipeline.

Muxes on the register file(s) allow for accessing GPRs and CSRs using the Access Register abstract command.

Memory is accessed using the Abstract Access Memory command or through System Bus Access.

This implementation could allow a debugger to collect information from the hart even when that hart is unable to execute instructions.

### A.2 Execution Based

This implementation only implements the Access Register abstract command for GPRs on a halted hart, and relies on the Program Buffer for all other operations. It uses the hart's existing pipeline and ability to execute from arbitrary memory locations to avoid modifications to a hart's datapath.

When the halt request bit is set, the Debug Module raises a special interrupt to the selected harts. This interrupt causes each hart to enter Debug Mode and jump to a defined memory region that is serviced by the DM and is only accessible to the harts in Debug Mode. Accesses to this memory should be uncached to avoid side effects from debugging operations. When taking this jump, `pc` is saved to `dpc` and `cause` is updated in `dcscr`. This jump is similar to a trap but it is not architecturally considered a trap, so for instance doesn't count as a trap for trigger behavior.

The code in the Debug Module causes the hart to execute a "park loop." In the park loop the hart writes its `mhartid` to a memory location within the Debug Module to indicate that it is halted.

To allow the DM to individually control one out of several halted harts, each hart polls for flags in a DM-controlled memory location to determine whether the debugger wants it to execute the Program Buffer or perform a resume.

To execute an abstract command, the DM first populates some internal words of program buffer according to `command`. When `transfer` is set, the DM populates these words with `lw <gpr>, 0x400(zero)` or `sw 0x400(zero), <gpr>`. 64- and 128-bit accesses use `ld/sd` and `lq/sq` respectively. If `transfer` is not set, the DM populates these instructions as `nops`. If `postexec` is set, execution continues to the debugger-controlled Program Buffer, otherwise the DM causes a `ebreak` to execute immediately.

When `ebreak` is executed (indicating the end of the Program Buffer code) the hart returns to its park loop. If an exception is encountered, the hart jumps to an address within the Debug Module. The code there causes the hart to write to the Debug Module indicating an exception. Then the hart jumps back to the park loop. The DM infers from the write that there was an exception, and sets `cmderr` appropriately. Typically the hart will execute a `fence` instruction before entering the park loop, to ensure that any effects from the abstract command, such as a write to `data0`, take effect before the DM returns `busy` to 0.

To resume execution, the debug module sets a flag which causes the hart to execute a `dret`. `dret` is an instruction that only has meaning while in Debug Mode and not executing from the Program Buffer. Its recommended encoding is `0x7b200073`. When `dret` is executed, `pc` is restored from `dpc` and normal execution resumes at the privilege set by `prv` and `v`.

`data0` etc. are mapped into regular memory at an address relative to `zero` with only a 12-bit `imm`. The exact address is an implementation detail that a debugger must not rely on. For example, the `data` registers might be mapped to `0x400`.

For additional flexibility, `progbuf0`, etc. are mapped into regular memory immediately preceding `data0`, in order to form a contiguous region of memory which can be used for either program execution or data transfer.

The PMP must not disallow fetches, loads, or stores in the address range associated with the Debug Module when the hart is in Debug Mode, regardless of how the PMP is configured. The same is true of PMA. Without this guarantee, the park loop would enter an infinite loop of traps and debug would not be possible.

### A.3 Debug Module Interface Signals

As stated in section 3.1 the details of the DMI are left to the system designer. It is quite often the case that only one DTM and one DM is implemented. In this case it might be useful to comply with the signals suggested in table A.1, which is the implementation used in the open-source `rocket-chip` RISC-V core.

The DTM can start a request when the DM sets `REQ_READY` to 1. When this is the case `REQ_OP` can be set to 1 for a read or 2 for a write request. The desired address is driven with the `REQ_ADDRESS` signal. Finally `REQ_VALID` is set high, indicating to the DM that a valid request is pending.

The DM must respond to a request from the DTM when RSP\_READY is high. The status of the response is indicated by the RSP\_OP signal (see [op](#)). The data of the response is driven to RSP\_DATA. A pending response is signalled by setting RSP\_VALID.

Signal	Width	Source	Description
REQ_VALID	1	DTM	Indicates that a valid request is pending
REQ_READY	1	DM	Indicates that the DM is able to process a request
REQ_ADDRESS	<a href="#">abits</a>	DTM	Requested address
REQ_DATA	32	DTM	Requested data
REQ_OP	2	DTM	Same meaning as the <a href="#">op</a> field
RSP_VALID	1	DM	Indicates that a valid respond is pending
RSP_READY	1	DTM	Indicates that the DTM is able to process a respond
RSP_DATA	32	DM	Response data
RSP_OP	2	DM	Same meaning as the <a href="#">op</a> field

Table A.1: Signals for the suggested DMI between one DTM and one DM



# Appendix B

## Debugger Implementation

### B.1 C Header File

<https://github.com/riscv/riscv-debug-spec> contains instructions for generating a C header file that defines macros for every field in every register/abstract command mentioned in this document.

### B.2 External Debugger Implementation

This section details how an external debugger might use the described debug interface to perform some common operations on RISC-V cores using the JTAG DTM described in Section 6.1. All these examples assume a 32-bit core but it should be easy to adapt the examples to 64- or 128-bit cores.

To keep the examples readable, they all assume that everything succeeds, and that they complete faster than the debugger can perform the next access. This will be the case in a typical JTAG setup. However, the debugger must always check the sticky error status bits after performing a sequence of actions. If it sees any that are set, then it should attempt the same actions again, possibly while adding in some delay, or explicit checks for status bits.

#### B.2.1 Debug Module Interface Access

To read an arbitrary Debug Module register, select `dmi`, and scan in a value with `op` set to 1, and `address` set to the desired register address. In Update-DR the operation will start, and in Capture-DR its results will be captured into `data`. If the operation didn't complete in time, `op` will be 3 and the value in `data` must be ignored. The busy condition must be cleared by writing `dmireset` in `dtmcs`, and then the second scan must be performed again. This process must be repeated until `op` returns 0. In later operations the debugger should allow for more time between Update-DR and Capture-DR.

To write an arbitrary Debug Bus register, select `dmi`, and scan in a value with `op` set to 2, and

[address](#) and [data](#) set to the desired register address and data respectively. From then on everything happens exactly as with a read, except that a write is performed instead of the read.

It should almost never be necessary to scan IR, avoiding a big part of the inefficiency in typical JTAG use.

## B.2.2 Checking for Halted Harts

A user will want to know as quickly as possible when a hart is halted (e.g. due to a breakpoint). To efficiently determine which harts are halted when there are many harts, the debugger uses the [haltsum](#) registers. Assuming the maximum number of harts exist, first it checks [haltsum3](#). For each bit set there, it writes [hartsel](#), and checks [haltsum2](#). This process repeats through [haltsum1](#) and [haltsum0](#). Depending on how many harts exist, the process should start at one of the lower [haltsum](#) registers.

## B.2.3 Halting

To halt one or more harts, the debugger selects them, sets [haltreq](#), and then waits for [allhalted](#) to indicate the harts are halted. Then it can clear [haltreq](#) to 0, or leave it high to catch a hart that resets while halted.

## B.2.4 Running

First, the debugger should restore any registers that it has overwritten. Then it can let the selected harts run by setting [resumereq](#). Once [allresumeack](#) is set, the debugger knows the selected harts have resumed. Harts might halt very quickly after resuming (e.g. by hitting a software breakpoint) so the debugger cannot use [allhalted](#)/[anyhalted](#) to check whether the hart resumed.

## B.2.5 Single Step

Using the hardware single step feature is almost the same as regular running. The debugger just sets [step](#) in [dcsr](#) before letting the hart run. The hart behaves exactly as in the running case, except that interrupts may be disabled (depending on [stepie](#)) and it only fetches and executes a single instruction before re-entering Debug Mode.

## B.2.6 Accessing Registers

### B.2.6.1 Using Abstract Command

Read [s0](#) using abstract command:

Op	Address	Value	Comment
Write	<code>command</code>	<code>aarsize = 2, transfer, regno = 0x1008</code>	Read <code>s0</code>
Read	<code>data0</code>	-	Returns value that was in <code>s0</code>

Write `mstatus` using abstract command:

Op	Address	Value	Comment
Write	<code>data0</code>	new value	
Write	<code>command</code>	<code>aarsize = 2, transfer, write, regno = 0x300</code>	Write <code>mstatus</code>

### B.2.6.2 Using Program Buffer

Abstract commands are used to exchange data with GPRs. Using this mechanism, other registers can be accessed by moving their value into/out of GPRs.

Write `mstatus` using program buffer:

Op	Address	Value	Comment
Write	<code>progbuf0</code>	<code>csrw s0, MSTATUS</code>	
Write	<code>progbuf1</code>	<code>ebreak</code>	
Write	<code>data0</code>	new value	
Write	<code>command</code>	<code>aarsize = 2, postexec, transfer, write, regno = 0x1008</code>	Write <code>s0</code> , then execute program buffer

Read `f1` using program buffer:

Op	Address	Value	Comment
Write	<code>progbuf0</code>	<code>fmv.x.s s0, f1</code>	
Write	<code>progbuf1</code>	<code>ebreak</code>	
Write	<code>command</code>	<code>postexec</code>	Execute program buffer
Write	<code>command</code>	<code>transfer, regno = 0x1008</code>	read <code>s0</code>
Read	<code>data0</code>	-	Returns the value that was in <code>f1</code>

## B.2.7 Reading Memory

### B.2.7.1 Using System Bus Access

With system bus access, addresses are physical system bus addresses.

Read a word from memory using system bus access:

Op	Address	Value	Comment
Write	<code>sbc</code>	<code>sbaccess = 2, sbreadonaddr</code>	Setup
Write	<code>sbaddress0</code>	address	
Read	<code>sbddata0</code>	-	Value read from memory

Read block of memory using system bus access:

Op	Address	Value	Comment
Write	<code>sbc</code>	<code>sbaccess = 2, sbreadonaddr, sbreadondata, sbautoincrement</code>	Turn on autoread and autoincrement
Write	<code>sbaddress0</code>	address	Writing address triggers read and increment
Read	<code>sbddata0</code>	-	Value read from memory
Read	<code>sbddata0</code>	-	Next value read from memory
...	...	...	...
Write	<code>sbc</code>	0	Disable autoread
Read	<code>sbddata0</code>	-	Get last value read from memory.

### B.2.7.2 Using Program Buffer

Through the Program Buffer, the hart performs the memory accesses. Addresses are physical or virtual (depending on `mprven` and other system configuration).

Read a word from memory using program buffer:

Op	Address	Value	Comment
Write	<code>progbuf0</code>	<code>lw s0, 0(s0)</code>	
Write	<code>progbuf1</code>	<code>ebreak</code>	
Write	<code>data0</code>	address	
Write	<code>command</code>	<code>transfer, write, postexec, regno = 0x1008</code>	Write <code>s0</code> , then execute program buffer
Write	<code>command</code>	<code>regno = 0x1008</code>	Read <code>s0</code>
Read	<code>data0</code>	-	Value read from memory

Read block of memory using program buffer:

Op	Address	Value	Comment
Write	<a href="#">progbuf0</a>	lw s1, 0(s0)	
Write	<a href="#">progbuf1</a>	addi s0, s0, 4	
Write	<a href="#">progbuf2</a>	ebreak	
Write	<a href="#">data0</a>	address	
Write	<a href="#">command</a>	transfer, write, postexec, regno = 0x1008	Write s0, then execute program buffer
Write	<a href="#">command</a>	postexec, regno = 0x1009	Read s1, then execute program buffer
Write	<a href="#">abstractauto</a>	autoexecdata [0]	Set <a href="#">autoexecdata</a> [0]
Read	<a href="#">data0</a>	-	Get value read from memory, then execute program buffer
Read	<a href="#">data0</a>	-	Get next value read from memory, then execute program buffer
...	...	...	...
Write	<a href="#">abstractauto</a>	0	Clear <a href="#">autoexecdata</a> [0]
Read	<a href="#">data0</a>	-	Get last value read from memory.

### B.2.7.3 Using Abstract Memory Access

Abstract memory accesses act as if they are performed by the hart, although the actual implementation may differ.

Read a word from memory using abstract memory access:

Op	Address	Value	Comment
Write	<a href="#">data1</a>	address	
Write	<a href="#">command</a>	cmdtype=2, <a href="#">aamsize</a> =2	
Read	<a href="#">data0</a>	-	Value read from memory

Read block of memory using abstract memory access:

Op	Address	Value	Comment
Write	<a href="#">abstractauto</a>	1	Re-execute the command when <a href="#">data0</a> is accessed
Write	<a href="#">data1</a>	address	
Write	<a href="#">command</a>	cmdtype=2, <a href="#">aamsize</a> =2, <a href="#">aampostincrement</a> =1	
Read	<a href="#">data0</a>	-	Read value, and trigger reading of next address
...	...	...	...
Write	<a href="#">abstractauto</a>	0	Disable auto-exec
Read	<a href="#">data0</a>	-	Get last value read from memory.

## B.2.8 Writing Memory

### B.2.8.1 Using System Bus Access

With system bus access, addresses are physical system bus addresses.

Write a word to memory using system bus access:

Op	Address	Value	Comment
Write	<code>sbc</code>	<code>sbaccess = 2</code>	Configure access size
Write	<code>sbaddress0</code>	address	
Write	<code>sbd</code>	value	

Write a block of memory using system bus access:

Op	Address	Value	Comment
Write	<code>sbc</code>	<code>sbaccess = 2, sbautoincrement</code>	Turn on autoincrement
Write	<code>sbaddress0</code>	address	
Write	<code>sbd</code>	value0	
Write	<code>sbd</code>	value1	
...	...	...	...
Write	<code>sbd</code>	valueN	

### B.2.8.2 Using Program Buffer

Through the Program Buffer, the hart performs the memory accesses. Addresses are physical or virtual (depending on `mprven` and other system configuration).

Write a word to memory using program buffer:

Op	Address	Value	Comment
Write	<code>progbuf0</code>	<code>sw s1, 0(s0)</code>	
Write	<code>progbuf1</code>	<code>ebreak</code>	
Write	<code>data0</code>	address	
Write	<code>command</code>	<code>transfer, write, regno = 0x1008</code>	Write <code>s0</code>
Write	<code>data0</code>	value	
Write	<code>command</code>	<code>transfer, write, postexec, regno = 0x1009</code>	Write <code>s1</code> , then execute program buffer

Write block of memory using program buffer:

Op	Address	Value	Comment
Write	<code>progbuf0</code>	<code>sw s1, 0(s0)</code>	
Write	<code>progbuf1</code>	<code>addi s0, s0, 4</code>	
Write	<code>progbuf2</code>	<code>ebreak</code>	
Write	<code>data0</code>	address	
Write	<code>command</code>	<code>transfer, write, regno = 0x1008</code>	Write <code>s0</code>
Write	<code>data0</code>	value0	
Write	<code>command</code>	<code>transfer, write, postexec, regno = 0x1009</code>	Write <code>s1</code> , then execute program buffer
Write	<code>abstractauto</code>	<code>autoexecdata [0]</code>	Set <code>autoexecdata [0]</code>
Write	<code>data0</code>	value1	
...	...	...	...
Write	<code>data0</code>	valueN	
Write	<code>abstractauto</code>	0	Clear <code>autoexecdata [0]</code>

### B.2.8.3 Using Abstract Memory Access

Abstract memory accesses act as if they are performed by the hart, although the actual implementation may differ.

Write a word to memory using abstract memory access:

Op	Address	Value	Comment
Write	<code>data1</code>	address	
Write	<code>data0</code>	value	
Write	<code>command</code>	<code>cmdtype=2, aamsize =2, write=1</code>	

Write a block of memory using abstract memory access:

Op	Address	Value	Comment
Write	<code>data1</code>	address	
Write	<code>data0</code>	value0	
Write	<code>command</code>	<code>cmdtype=2, aamsize =2, write=1, aampostincrement =1</code>	
Write	<code>abstractauto</code>	1	Re-execute the command when <code>data0</code> is accessed
Write	<code>data0</code>	value1	
Write	<code>data0</code>	value2	
...	...	...	...
Write	<code>data0</code>	valueN	
Write	<code>abstractauto</code>	0	Disable auto-exec

### B.2.9 Triggers

A debugger can use hardware triggers to halt a hart when a certain event occurs. Below are some examples, but as there is no requirement on the number of features of the triggers implemented by a hart, these examples might not be applicable to all implementations. When a debugger wants to set a trigger, it writes the desired configuration, and then reads back to see if that configuration is supported. All examples assume XLEN=32.

Enter Debug Mode when the instruction at 0x80001234 is executed, to be used as an instruction breakpoint in ROM:

<b>tdata1</b>	0x6980105c	type=6, dmode=1, action=1, select=0, match=0, m=1, s=1, u=1, vs=1, vu=1, execute=1
<b>tdata2</b>	0x80001234	address

Enter Debug Mode when performing a load at address 0x80007f80 in M-mode or S-mode or U-mode:

<b>tdata1</b>	0x68001059	type=6, dmode=1, action=1, select=0, match=0, m=1, s=1, u=1, load=1
<b>tdata2</b>	0x80007f80	address

Enter Debug Mode when storing to an address between 0x80007c80 and 0x80007cef (inclusive) in VS-mode or VU-mode when hgatp.VMID=1:

<b>tdata1 0</b>	0x69801902	type=6, dmode=1, action=1, chain=1, select=0, match=2, vs=1, vu=1, store=1
<b>tdata2 0</b>	0x80007c80	start address (inclusive)
<b>textra32 0</b>	0x03000000	mhselect=6, mhvalue=0
<b>tdata1 1</b>	0x69801182	type=6, dmode=1, action=1, select=0, match=3, vs=1, vu=1, store=1
<b>tdata2 1</b>	0x80007cf0	end address (exclusive)
<b>textra32 1</b>	0x03000000	mhselect=6, mhvalue=0

Enter Debug Mode when storing to an address between 0x81230000 and 0x8123fff (inclusive):

<b>tdata1</b>	0x698010da	type=6, dmode=1, action=1, select=0, match=1, m=1, s=1, u=1, vs=1, vu=1, store=1
<b>tdata2</b>	0x81237fff	16 upper bits to match exactly, then 0, then all ones.

Enter Debug Mode when loading from an address between 0x86753090 and 0x8675309f or between 0x96753090 and 0x9675309f (inclusive):

<b>tdata1 0</b>	0x69801a59	type=6, dmode=1, action=1, chain=1, match=4, m=1, s=1, u=1, vs=1, vu=1, load=1
<b>tdata2 0</b>	0xfff03090	Mask for low half, then match for low half
<b>tdata1 1</b>	0x698012d9	type=6, dmode=1, action=1, match=5, m=1, s=1, u=1, vs=1, vu=1, load=1
<b>tdata2 1</b>	0xefff8675	Mask for high half, then match for high half



### B.2.10 Handling Exceptions

Generally the debugger can avoid exceptions by being careful with the programs it writes. Sometimes they are unavoidable though, e.g. if the user asks to access memory or a CSR that is not implemented. A typical debugger will not know enough about the hardware platform to know what's going to happen, and must attempt the access to determine the outcome.

When an exception occurs while executing the Program Buffer, `cmderr` becomes set. The debugger can check this field to see whether a program encountered an exception. If there was an exception, it's left to the debugger to know what must have caused it.

### B.2.11 Quick Access

There are a variety of instructions to transfer data between GPRs and the `data` registers. They are either loads/stores or CSR reads/writes. The specific addresses also vary. This is all specified in `hartinfo`. The examples here use the pseudo-op `transfer dest, src` to represent all these options.

Halt the hart for a minimum amount of time to perform a single memory write:

Op	Address	Value	Comment
Write	<code>progbuf0</code>	<code>transfer arg2, s0</code>	Save <code>s0</code>
Write	<code>progbuf1</code>	<code>transfer s0, arg0</code>	Read first argument (address)
Write	<code>progbuf2</code>	<code>transfer arg0, s1</code>	Save <code>s1</code>
Write	<code>progbuf3</code>	<code>transfer s1, arg1</code>	Read second argument (data)
Write	<code>progbuf4</code>	<code>sw s1, 0(s0)</code>	
Write	<code>progbuf5</code>	<code>transfer s1, arg0</code>	Restore <code>s1</code>
Write	<code>progbuf6</code>	<code>transfer s0, arg2</code>	Restore <code>s0</code>
Write	<code>progbuf7</code>	<code>ebreak</code>	
Write	<code>data0</code>	address	
Write	<code>data1</code>	data	
Write	<code>command</code>	<code>0x10000000</code>	Perform quick access

This shows an example of setting the `m` bit in `mcontrol` to enable a hardware breakpoint in M-mode. Similar quick access instructions could have been used previously to configure the trigger that is being enabled here:

Op	Address	Value	Comment
Write	<code>progbuf0</code>	<code>transfer arg0, s0</code>	Save <code>s0</code>
Write	<code>progbuf1</code>	<code>li s0, (1 &lt;&lt; 6)</code>	Form the mask for <code>m</code> bit
Write	<code>progbuf2</code>	<code>csrwr x0, tdata1, s0</code>	Apply the mask to <code>mcontrol</code>
Write	<code>progbuf3</code>	<code>transfer s0, arg2</code>	Restore <code>s0</code>
Write	<code>progbuf4</code>	<code>ebreak</code>	
Write	<code>command</code>	<code>0x10000000</code>	Perform quick access

## B.3 Native Debugger Implementation

The spec contains a few features to aid in writing a native debugger. This section describes how some common tasks might be achieved.

### B.3.1 Single Step

Single step is straightforward if the OS or a debug stub runs in M-Mode while the program being debugged runs in a less privileged mode. When a step is required, the OS or debug stub writes `count =1`, `action =0`, `m =0` before returning control to the lower user program with an `mret` instruction.

Stepping code running in the same privilege mode as the debugger is more complicated, depending on what other debug features are implemented.

If hardware implements `mpte` and `mte`, then stepping through non-trap code which doesn't allow for nested interrupts is also straightforward.

If hardware automatically prevents `action =0` triggers from matching when entering a trap handler as described in Section 5.4, then a carefully written trap handler can ensure that interrupts are disabled whenever the `icount` trigger must not match.

If neither of these features exist, then single step is doable, but tricky to get right. To single step, the debug stub would execute something like:

```
li    t0, \FcsrIcountCount=4, \FcsrIcountAction=0, \FcsrIcountM=1
cswr tdata1, t0    /* Write the trigger. */
lw    t0, 8(sp)    /* Restore t0, count decrements to 3 */
lw    sp, 0(sp)    /* Restore sp, count decrements to 2 */
mret                                /* Return to program being debugged. count decrements to 1 */
```

There is an additional problem with using `icount` to single step. An instruction may cause an exception into a more privileged mode where the trigger is not enabled. The exception handler might address the cause of the exception, and then restart the instruction. Examples of this include page faults, FPU instructions when the FPU is not yet enabled, and interrupts. When a user is single stepping through such code, they will have to step twice to get past the restarted instruction. The first time the exception handler runs, and the second time the instruction actually executes. That is confusing and usually undesirable.

To help users out, debuggers should detect when a single step restarted an instruction, and then step again. This way the users see the expected behavior of stepping over the instruction. Ideally the debugger would notify the user that an exception handler executed the first time.

The debugger should perform this extra step when the PC doesn't change during a regular step.

---

*It is safe to perform an extra step when the PC changes, because every RISC-V instruction either changes the PC or has side effects when repeated, but never both.*

To avoid an infinite loop if the exception handler does not address the cause of the exception, the debugger must execute no more than a single extra step.

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# Appendix C

## Change Log

Revision	Date	Author(s)	Description
92835a8	2023-10-11	Tim Newsome	AR: Clarify itrigger and trigger number translation
b699db6	2023-10-09	Tim Newsome	AR: Clarify itrigger behavior.
9f3fd6f	2023-10-05	Tim Newsome	Add people with git PRs to credits.
aaab8fb	2023-10-03	Tim Newsome	AR: Update priority table from latest privspec
1b81755	2023-10-04	Tim Newsome	AR: Clarify triggers and multiple state changes
4aad021	2023-10-04	Tim Newsome	Revert "AR: Clarify triggers and multiple state changes"
79e33fb	2023-10-04	Tim Newsome	AR: Clarify triggers and multiple state changes
9ddcbcb	2023-10-04	Tim Newsome	AR: Define mcontrol* triggers and multiple accesses (#883)
9590e02	2023-10-04	Tim Newsome	AR: Clarify what tmexttrigger.intctl is for (#898)
a83e118	2023-10-02	Tim Newsome	AR: Clarify why icount matches on every trap.
6a5800c	2023-10-02	Tim Newsome	AR: Clarify section/chapter in Sdext.
2e78518	2023-09-28	Tim Newsome	Tighten up 'make clean'
d404fd9	2023-09-28	Tim Newsome	Don't build serial.tex.
1aab674	2023-09-18	Tim Newsome	AR: Remove Message Registers.
276b60a	2023-09-26	Tim Newsome	Fix debugger examples: use transfer with write
182beaa	2023-09-26	Tim Newsome	dscratch[01] are DXLEN bits wide.
3ef1852	2023-09-20	Tim Newsome	AR: List which extensions were considered.
0563d2f	2023-09-20	Tim Newsome	AR: Clarify mcontext/hcontext.
51fc095	2023-09-19	Tim Newsome	AR: tcontrol.mie applies to all traps, not just breakpoints
b4a51a3	2023-09-18	Tim Newsome	AR: Clarify unimplemented Sdtrig.
2cd44e5	2023-09-16	Evgeniy Naydanov	Address review comments on 'debug_defines.h'
2ace90c	2023-09-14	Tim Newsome	AR: Clarify what debuggers can assume about MRs
3f0a9e1	2023-09-14	Tim Newsome	AR: Remove comment about privspec/CSR behavior
2c8f38a	2023-09-11	Tim Newsome	Rebuild PDF.
b972694	2023-09-11	Tim Newsome	ARC: Comment on unimplemented CSR behavior.
d67cca8	2023-07-17	Evgeniy Naydanov	Add decoding facilities into debug_defines.h

b92c427	2023-08-30	Tim Newsome	Sdtrig must implement at least one trigger.
470ca90	2023-08-24	Tim Newsome	Clarify that each DMI address points at a 32-bit register.
925e7ce	2023-08-22	Jan Matyas	Program Buffer: Clarify requirements for case progbufsize==1 (#868)
648ed94	2023-08-15	Tim Newsome	Make ‘make clean‘ more exhaustive.
4472612	2023-08-15	Tim Newsome	AR: Define ”native debug” in Background section.
83483b1	2023-08-21	Tim Newsome	AR: Don’t use ”non-existent” in reference to register.
dcb37f2	2023-08-15	Tim Newsome	AR: Clarify message register uses.
e3671ac	2023-08-15	Tim Newsome	Recommend debug ROM is not cached.
2c19654	2023-08-11	Paul Donahue	Update chapter names per ARC review feedback
f44beaf	2023-07-31	Tim Newsome	AR: Use manager/subordinate for master/slave.
314b09f	2023-07-26	Tim Newsome	Fix type in execution appendix: execute -i postexec
ebde108	2023-07-10	Tim Newsome	Rebuild PDF.
d13d1e2	2023-06-23	Tim Newsome	In Debug Mode, triggers don’t match or fire.
9001957	2023-06-16	ved-rivos	Debug interaction with wrs.sto/wrs.nto (#838)
f698285	2023-06-14	Tim Newsome	Rebuild PDF for AR.
6a56dea	2023-06-12	Tim Newsome	External trigger signals inputs don’t cause trigger output.
5967998	2023-06-08	Tim Newsome	Rebuild PDF.
ed97c71	2023-05-10	Tim Newsome	Add dtmcs.errinfo.
0bc6cde	2023-05-16	Tim Newsome	Clarify trap in progbuf behavior.
a5fadee	2023-05-10	Tim Newsome	D-Mode exceptions may modify DPC
909fa7e	2023-05-09	Tim Newsome	Don’t say master/slave
f67fbb0	2023-05-04	Tim Newsome	Clarify that icount decrements on interrupt traps.
a06983a	2023-04-28	Tim Newsome	Clarify when hit0/hit1 are updated by TM.
b5652d8	2023-04-18	Tim Newsome	Reentrant trigger prevention inhibits match *or fire*
6c642e9	2023-04-14	Tim Newsome	Rebuild PDF for substantive freeze.
f4381fe	2023-04-14	Tim Newsome	Describe how CMO instructions match triggers. (#821)
92bf5c6	2023-04-14	Tim Newsome	Rebuild PDF.
fb70252	2023-04-14	Tim Newsome	Rebuild PDF.
8efe2d6	2023-04-13	Paul Donahue	Clarify that stepie applies to dcsr.step but not icount.
4458d62	2023-04-07	Tim Newsome	Clarify action 8/9 signal TM external triggers.
da5fd2a	2023-04-07	Tim Newsome	Rename dtmc.dmihardreset field to dtmhardreset
fb32087	2023-04-06	Paul Donahue	Remove references to the timing bit. (#831)
f22a2fb	2023-04-05	Tim Newsome	Document native single step options.
1862d2f	2023-03-30	Tim Newsome	icount: Remove confusing statement about inhibited trigger.
c17e508	2023-03-23	Tim Newsome	extra inhibits firing in addition to matching.
58bc953	2023-03-23	Tim Newsome	Document behavior for abstract dpc access while running.
852dd18	2023-03-23	Tim Newsome	Move dpc info from Program Buffer into dpc register.
328be72	2023-03-15	Tim Newsome	Describe what happens on halt.
754d677	2023-03-13	Tim Newsome	Clarify that dmi.op describe dmi.address contents.
3ca18ed	2023-03-13	Tim Newsome	Rebuild PDF.
2fdf271	2023-02-21	Tim Newsome	Document trigger behavior with vector/push/pop accesses.

efb16bc	2022-12-28	Tim Newsome	Clarify: non-existent CSR access raises an exception.
f297263	2023-03-06	Paul Donahue	Fix #817
c604228	2023-02-06	Tim Newsome	Add mcontrol6.uncertain,uncertainen bits.
5f77aa1	2023-02-21	Tim Newsome	Remove mcontrol6.size encodings for 80, 96, 112 bit
0374fd6	2023-02-17	Tim Newsome	Remove timing bit again. Introduce tinfo.version.
ef54e62	2023-02-21	Tim Newsome	Refactor AMO discussion.
16a3df8	2023-02-20	Tim Newsome	Clarify trigger action=1 requires Sdext.
58fb64c	2023-02-20	Tim Newsome	stepie=0 should be supported.
57e7c6d	2023-02-16	Tim Newsome	Rebuild PDF.
89800fd	2023-02-16	Tim Newsome	Revert "Remove mcontrol6.timing."
69066a6	2023-02-15	Tim Newsome	Harts/triggers belong to exactly one halt/resume group.
7e3bf7c	2023-01-27	Tim Newsome	Suggest debugger behavior when icount stepping
12750d4	2023-02-10	Tim Newsome	Clarify JTAG connector section. (#782)
1776835	2023-02-09	Tim Newsome	Rebuild PDF.
9244505	2023-02-07	Tim Newsome	Say "retired" instead of "committed."
91b7ae6	2023-01-23	Tim Newsome	Remove mcontrol6.timing.
7d7031e	2023-01-23	Tim Newsome	Expand mcontrol6.hit into hit0 and hit1.
5a5c078	2023-01-27	Tim Newsome	Clarify icount decrements on instruction execution
318639d	2023-01-27	Tim Newsome	Fix typo.
d1c7970	2023-01-17	Tim Newsome	Make capitalization more consistent.
c93823e	2023-01-09	Tim Newsome	Rebuild PDF.
36ae10c	2022-12-28	Tim Newsome	Fix git mess.
069b0a3	2022-12-28	Tim Newsome	Revert "Explicitly allow debug mode to access debug mode CSRs for #772. (#773)"
e0c3ebf	2022-12-28	Tim Newsome	Clarify: non-existent CSR access raises an exception.
98bebe3	2022-12-28	Paul Donahue	Explicitly allow debug mode to access debug mode CSRs for #772. (#773)
47637a1	2022-12-27	Tim Newsome	Clarify what fire and match mean for a trigger. (#776)
05252da	2022-12-12	Tim Newsome	Resumereq is W1 so cannot be cleared. (#788)
4c72024	2022-12-08	Tim Newsome	Rebuild PDF.
246028c	2022-12-08	Paul Donahue	Clarify the term "trap" in execution-based implementations. (#767)
8601134	2022-12-06	Tim Newsome	Remove longtable.sty. (#781)
dbc2388	2022-12-05	Paul Donahue	Low bit of dpc is read-only. (#779)
e481f91	2022-11-28	Tim Newsome	Improve Native Triggers formatting. (#777)
6e6cac1	2022-11-21	Paul Donahue	Address #771 (#774)
bd39106	2022-11-15	Tim Newsome	halt/resume groups don't apply to unavailable harts. (#764)
d1c1f0e	2022-10-31	Tim Newsome	Clarify stoptime behavior. (#769)
68aa0dc	2022-10-26	Tim Newsome	Dmode is still active for trigger type 15. (#768)
0e4d91e	2022-10-20	Jan Matyas	abstractcs.relaxedpriv: Clarify the description (#765)
1a39eb7	2022-10-19	Tim Newsome	dmi.data,address are unspecified after write/nop (#762)
8f396dc	2022-10-19	Paul Donahue	More prominently say that mcontrol is deprecated. (#759)



0250a06	2022-10-13	Tim Newsome	Rebuild PDF.
86e748a	2022-09-29	Tim Newsome	Clarify that PMP shouldn't prevent debug. (#758)
aed6d27	2022-09-19	Tim Newsome	Clarify dcsr.cause when halt groups are used. (#752)
14a8d62	2022-09-12	Tim Newsome	Rebuild PDF.
9dc1ee4	2022-09-03	YenHaoChen	Clarify 'dmode is cleared' means 'the new value of dmode would be 0' (#748)
d657b8a	2022-08-31	Paul Donahue	Simplify definition of Minimal Debug Spec.
7882ff4	2022-08-31	Paul Donahue	Minor wording change
ed697a2	2022-08-30	Paul Donahue	Change document title.
21706b1	2022-08-29	Paul Donahue	Additional architecture review feedback. Fixed oddly-behaving link to DM chapter.
fe2a7f4	2022-08-29	Tim Newsome	Powers- <sub>i</sub> Power in NAPOT definition.
c2a57e1	2022-08-29	Tim Newsome	Move contributors to preface.
fd97d9e	2022-08-29	Tim Newsome	Change preface warning for stable builds.
b71e15d	2022-08-24	Tim Newsome	Elaborate on version detection. (#741)
d749752	2022-07-18	Tim Newsome	Remove license from debug_defines.h. (#743)
198ab7c	2022-06-27	Tim Newsome	Remove clarifying information that is false. (#740)
e5ce2b2	2022-06-16	Paul Donahue	Fix order of states to be consistent with a few sentences earlier. (#738)
1744041	2022-06-17	Yanqi Yang	Clarify the current version in readme. (#739)
d401f57	2022-06-09	Tim Newsome	Rebuild PDF.
c26287e	2022-06-09	Paul Donahue	Change dcsr.cause priorities (#693)
baae96	2022-05-26	Jiuyang Liu	Add GitHub Action to auto release. (#733)
182b9c4	2022-05-20	Tim Newsome	List trigger actions described by the trace spec (#734)
c40847a	2022-05-20	Tim Newsome	Add trigger action values, for debug_defines.h. (#732)
efd25df	2022-05-19	Tim Newsome	Add field values to XML register description. (#727)
3ae76ea	2022-05-18	Tim Newsome	The program buffer may not be readable. (#731)
42224a6	2022-05-16	Tim Newsome	Make data registers Message Registers. (#728)
df943f2	2022-05-12	Tim Newsome	Describe debug_defines.h. (#729)
1c095fd	2022-05-12	Tim Newsome	Fix formatting error at the end of document. (#730)
9137a14	2022-05-11	Tim Newsome	Add arguments to C macros that need them. (#724)
4c0859b	2022-05-11	Tim Newsome	mte/mpte apply only to breakpoint traps (#723)
02a2793	2022-05-10	Tim Newsome	Tighten up language around icount matching. (#722)
127e223	2022-05-10	Tim Newsome	Select and issue a command can be simultaneous (#725)
8c22825	2022-05-09	Tim Newsome	Rebuild PDF.
17187e4	2022-04-26	Tim Newsome	Use trigger type 15 to indicate disabled. (#721)
8705fd1	2022-04-22	Paul Donahue	Clarifications/bug fixes: mscontext and tselect (#702)
12a3eb6	2022-04-07	Tim Newsome	Rebuild PDF.
abefdcf	2022-04-07	Paul Donahue	Clarify behavior of low bits of tdata2 when match=1. (#680)
191c8bf	2022-04-07	Tim Newsome	dtmcs.dmistat is a read-only alias of dmi.op. (#720)
70bdd3f	2022-04-06	Tim Newsome	Define what 0 in register fields means. (#719)

598d4eb	2022-04-01	Paul Donahue	constrptr points to the configuration structure, not string. (#714)
ab0341e	2022-04-01	Paul Donahue	hawindow/hawindowssel are WARL (#715)
eef4d5	2022-03-29	Paul Donahue	If masking is supported then tdata2 must be able to hold all XLEN bits. (#716)
356feb1	2022-03-24	Tim Newsome	Fix contradiction in icount description. (#665)
2676fa1	2022-03-24	Tim Newsome	Mention textra in the trigger types that it affects. (#711)
3b16253	2022-03-24	Tim Newsome	Clarify writing to dmcs2.dmexttrigger. (#705)
9c52c14	2022-03-22	Tim Newsome	Fix reset value of etrigger.type (#712)
e9994d7	2022-03-21	Tim Newsome	Move nmi from etrigger to itrigger. (#709)
98203bb	2022-03-16	Paul Donahue	Clarify tmexttrigger with action=0. (#707)
72ab3ae	2022-03-08	Paul Donahue	ackunavail doesn't clear unavail on harts that are still unavailable (#706)
b659d7d	2022-02-10	Tim Newsome	Rebuild PDF.
fe3d1e6	2022-02-09	Tim Newsome	Clarify partial abstract reg writes. (#704)
5b219ed	2022-02-08	Tim Newsome	Fix uses of \index macro. (#703)
6ea272f	2022-01-18	Paul Donahue	Stoptime clarification per architecture review (#699)
e83441e	2022-01-10	Tim Newsome	Rebuild PDF.
aa07a68	2022-01-10	Tim Newsome	Clarify that tdata1 is WARL (#701)
f3323bb	2022-01-07	Paul Donahue	Remove hcxe and scxe in tcontrol, using Smstateen instead. (#700)
e350c4b	2022-01-07	Paul Donahue	Define mtval behavior on action=0 triggers (#695)
106dd3c	2021-12-29	Paul Donahue	"privilege level" -i "privilege mode" per architecture review feedback (#698)
337b8a5	2021-12-29	Paul Donahue	Clarify behavior of writes to abstractauto while busy=1 (#696)
b187cff	2021-12-16	Paul Donahue	Specify accessibility of mscontext (#694)
92aef97	2021-12-16	Tim Newsome	Resume ack is set for any resume cause. (#692)
f7dcc8c	2021-12-13	Tim Newsome	Label Ztrigger and Zdebug chapters. (#669)
3f44538	2021-12-09	Tim Newsome	Rebuild PDF.
f3c0997	2021-11-09	Tim Newsome	autoexec* is the same as writing command (#685)
df0ea10	2021-11-08	Tim Newsome	Rebuild PDF.
e3bd339	2021-11-08	Paul Donahue	Fix obvious error in dcsr.cause=2 description. (#681)
017e943	2021-11-08	Paul Donahue	Debuggers may set action!=1 when clearing dmode (#679)
3dfe4f7	2021-10-26	Paul Donahue	Make mcontrol6.hit description consistent with mcontrol.hit (#687)
42b8be9	2021-10-20	A.M	fixed wrong ref in Terminology section and rebuilt PDF (#684)
1a72df3	2021-10-07	Tim Newsome	Rebuild PDF.
b992bc5	2021-10-07	Paul Donahue	Clarify that etrigger.nmi ignores vs and vu (not just m, s, u) (#676)
55c492d	2021-09-17	Paul Donahue	Clarify ASID filtering (#667)
031ff77	2021-09-15	Tim Newsome	SBA write memory example: configure size (#673)
80563e2	2021-09-14	Tim Newsome	Run 'apt update' (#672)
fc52dc5	2021-09-14	Tim Newsome	Don't call read/write CSRs read-only. (#668)

57c1233	2021-09-13	Tim Newsome	Rebuild PDF.
6e127fc	2021-08-16	Paul Donahue	Native triggers and reentrancy (#660)
6a9b355	2021-08-13	Tim Newsome	Clarify unused debug connector pins names/uses. (#664)
f77291e	2021-08-12	Tim Newsome	Rebuild PDF.
33cd3a6	2021-08-10	Paul Donahue	Clarify dcsr.cause=2 (#663)
8f7873d	2021-07-28	Paul Donahue	Clarify dpc for non-mcontrol triggers (#624)
70e3db2	2021-07-13	Daniel Mangum	Update link to mailing list in README.md (#658)
63c985f	2021-07-12	Tim Newsome	Rebuild PDF.
9ac6506	2021-07-08	Paul Donahue	xepc instead of exception PC (#654)
e6b7b6c	2021-07-08	Tim Newsome	Clarify postincrement/transfer behavior. (#655)
f393a8d	2021-07-08	Daniel Mangum	Fix small grammatical error in suggested DMI signals (#656)
d584f0f	2021-07-08	Tim Newsome	Attempt to build the document using github actions. (#657)
022d62f	2021-06-29	Paul Donahue	Clarify hstatus when the H extension isn't implemented (#646)
c391ffe	2021-06-16	Paul Donahue	MPRV is in mstatus, not mcontrol. (#651)
a1e05fb	2021-06-14	Paul Donahue	Clarify what happens to halted harts upon DM reset (#648)
0f4ea2f	2021-06-10	Tim Newsome	Rebuild PDF.
641cd87	2021-06-09	Tim Newsome	Unselected harts may change groups when hgwrite=1 (#642)
55333c4	2021-06-09	Tim Newsome	hasel etc. are only ignored by abstract commands. (#643)
3704486	2021-06-09	Tim Newsome	Mention license in debug_defines.h. (#641)
b68c265	2021-06-08	Paul Donahue	Single stepping an instruction that fires a trigger (#622)
fdecf04	2021-06-08	Tim Newsome	Labels must come after (or in) their caption. (#640)
5350603	2021-06-04	Paul Donahue	Traps caused by action=0 triggers can be delegated via medeleg (#637)
11bf0db	2021-05-28	Tim Newsome	Divide the spec into ISA and non-ISA. (#635)
3334910	2021-05-18	Tim Newsome	Further diagram update. (#632)
b344c8f	2021-05-14	Tim Newsome	External debuggers should set dmode. (#634)
d9434bc	2021-05-14	Paul Donahue	Fix #630 (#633)
39fab4b	2021-05-10	Tim Newsome	Rebuild PDF.
090eac8	2021-05-04	Tim Newsome	Update Run/Halt Debug State Machine. (#629)
31a5f61	2021-04-08	Tim Newsome	Rebuild PDF.
e3e408c	2021-03-25	Paul Donahue	Clarify triggers (#628)
e0e0b4d	2021-03-23	Tim Newsome	Remove latest draft link. (#627)
1b665f3	2021-03-22	Paul Donahue	clarify abstractauto (#625)
7242fe1	2021-03-17	Paul Donahue	Clarify DM behavior for non-existent registers (#621)
9ae7560	2021-03-10	Paul Donahue	Fix minor typos (#620)
bd786dc	2021-02-08	Tim Newsome	Rebuild PDF.
132ffb8	2021-02-02	Thomas Wicki	Clarify MCONTROL/MCONTROL6 'timing' description (#614)
075de11	2021-02-02	Paul Donahue	Clarify aarpostincrement wrapping (#613)

fb15173	2021-01-25	Paul Donahue	hartsel and hasel are WARL (#612)
fcc2a33	2021-01-11	Tim Newsome	Rebuild PDF.
1b9caa1	2021-01-11	Tim Newsome	Define some more terms/acronyms. (#610)
822f63b	2021-01-05	Tim Newsome	Move dret completely into the appendix. (#611)
57e271d	2020-12-30	Tim Newsome	Use A-mode instead of "A mode" consistently. (#609)
08e072a	2020-12-29	Ernie Edgar	update version of priv spec to one with hypervisor (#608)
bb578a4	2020-12-28	Tim Newsome	Clarify details around fence and progbuf. (#601)
7e47254	2020-12-28	Tim Newsome	Rebuild PDF.
90ba168	2020-12-21	Tim Newsome	Clarify how many bits mcontrol/mcontrol6 compare. (#604)
27d735c	2020-12-18	Tim Newsome	Clarify maskmax corner cases. (#607)
2930c1d	2020-12-18	Tim Newsome	Clarify that ASID might come from satp or vsatp. (#606)
c98d7e6	2020-12-18	Tim Newsome	Mark 1.0 as STABLE. (#605)
edd6482	2020-12-17	Tim Newsome	Breakpoint in trap handler *might* be unrecoverable. (#603)
c49e9c3	2020-12-16	Tim Newsome	Rebuild PDF.
4772b19	2020-12-03	Ernie Edgar	Add version 1.0 value to dtmcs.version field (#602)
393d965	2020-11-23	Tim Newsome	Document changes since 0.13. (#600)
f2ff7a6	2020-11-16	Tim Newsome	Rebuild PDF.
2d7190c	2020-11-16	Tim Newsome	Add clic bit to tmexttrigger. (#599)
0198481	2020-11-16	Tim Newsome	Explain how to simply write any trigger. (#598)
b35af12	2020-11-16	Tim Newsome	Clarify trigger CSR behavior when XLEN changes. (#597)
f116aea	2020-11-16	benscotstaveley	add dmstatus.ndmresetpending to allow a debugger to determine when ndmreset is complete (#594)
3dd952c	2020-11-13	Tim Newsome	Chains should all have the same type. (#596)
34f80c6	2020-11-13	Tim Newsome	AMO operations may be ignored by mcontrol triggers. (#595)
9fea4c5	2020-11-13	Tim Newsome	Add keepalive feature. (#592)
53191a4	2020-11-12	Bruce Able-dinger	Added sbytemask field to textra32 and textra64 (#588)
26040fd	2020-11-11	Tim Newsome	hit must be set on fire, may be set on match (#593)
a2cf8fc	2020-11-11	Tim Newsome	Create LaTeX macros for fields without descriptions. (#591)
2bde1e4	2020-11-10	Tim Newsome	Change version number to 1.0. (#590)
75a7607	2020-11-10	Paul Donahue	Fix #587 (#589)
cabb06d	2020-11-09	Tim Newsome	Rebuild PDF.
c89895c	2020-11-06	Tim Newsome	Change quick access exceptions to halt the target. (#585)
46804a6	2020-11-04	Paul Donahue	Recommend mprven=1 (#580)
219d105	2020-11-03	Tim Newsome	Add pending state/bit to icount. (#574)
b860d53	2020-11-03	Tim Newsome	RISC-V External Debug Support -i RISC-V Debug Support (#581)
fd654d4	2020-11-02	Paul Donahue	textra32/64 also affects the new mcontrol6 triggers (#578)

3d24926	2020-10-30	Tim Newsome	List all DM registers in Table 3.8. (#579)
971d0aa	2020-10-30	Paul Donahue	Add support for the A extension (#561)
d589bc3	2020-10-30	Paul Donahue	Remove all uses of the ambiguous term "may not" (#576)
db8e814	2020-10-29	Paul Donahue	Triggers affect harts, not the system. For instance, there may be (#575)
c877e9c	2020-10-29	Paul Donahue	Added exttrigger capability as type 7 (#543)
0a81ec3	2020-10-28	Tim Newsome	Debuggers should know when harts are unavailable. (#520)
65af35f	2020-10-26	Ernie Edgar	Update debug_module.tex (#577)
99dfc98	2020-10-23	Scott Johnson	Don't decrement icount.count when exception is blocked by tcontrol.mte (#557)
66481cf	2020-10-16	Paul Donahue	Fix broken reference (#567)
0b42843	2020-10-16	benscotstaveley	require polling of dmactive low as well as dmactive high transitions (#566)
1e81d58	2020-10-16	Paul Donahue	Use official RISC-V terminology (#564)
5c65dfe	2020-10-09	Tim Newsome	Rebuild PDF.
9c083f1	2020-10-05	Ernie Edgar	Clarify that PMP must allow access to DM for debug to be possible (#554)
7aa5978	2020-09-22	Paul Donahue	Hypervisor support (#549)
597281c	2020-09-16	Ernie Edgar	Update rocket-chip link to specific commit for permanence (#552)
072affe	2020-09-16	Paul Donahue	Follow suggestion in #544 (#548)
8345674	2020-09-15	Paul Donahue	Add dret to rule 8. (#547)
175090c	2020-09-11	Paul Donahue	Add mcontrol6. (#538)
de1ec1a	2020-09-10	Paul Donahue	Fix links to point to fields in the correct registers. (#546)
4f625ca	2020-09-03	Ernie Edgar	Add scontext2 alias for scontext (#535)
7c0a6d5	2020-08-28	Jan Matyas	aamvirtual: Clarification for systems without address translations (#542)
30b1a97	2020-08-24	Tim Newsome	Remove end-of-line whitespace in generated comments. (#540)
6e90a60	2020-08-21	Tim Newsome	Add header to debug_defines.h (#539)
0200b27	2020-08-21	Tim Newsome	Improve formatting of autogenerated C header files (#537)
97d51c2	2020-08-11	Tim Newsome	Rebuild PDF.
fcf4002	2020-08-11	Tim Newsome	authdata should only should be implemented if used (#521)
0570f14	2020-08-05	Paul Donahue	Add abstractcs.relaxedpriv (#536)
2210002	2020-07-07	Tim Newsome	Rebuild PDF.
b9959e5	2020-06-30	Tim Newsome	Make explicit that aampostincrement is optional. (#532)
67fed8f	2020-06-19	Tim Newsome	Explicitly allow uni-directional external triggers. (#526)
9c69bf3	2020-06-09	Tim Newsome	Rebuild PDF.
85bf4df	2020-05-21	Tim Newsome	Add Kai Meinhard to contributors list.
2f1c133	2020-05-21	Kai Meinhard	Appendix B suggests signals for a DMI with one DTM connected to one DM (#524)

708b1e0	2020-04-10	Tim Newsome	Add Larry Madar.
e02a8b6	2020-04-07	Tim Newsome	Rebuild PDF.
372b27f	2020-03-23	Tim Newsome	All tdata functionality is optional... (#444)
50f5c8f	2020-03-11	Tim Newsome	Explicitly allow hard-coded halt/resume groups. (#517)
f4794bb	2020-03-10	Tim Newsome	Rebuild PDF.
e3ec24e	2020-02-13	bdwyatt	Adding version encoding for 0.14 spec. (#512)
cf9a884	2020-02-11	Tim Newsome	Rebuild PDF.
fdd5ad6	2020-02-11	Philipp Wagner	dcsr.prv should be WARL, not R/W (#498)
38b2794	2020-02-11	Tim Newsome	sizehi only exists if Xlen <sub>i</sub> 64. (#514)
5a54283	2020-01-16	Tim Newsome	Use exception, trap, and interrupt as in ISA spec (#511)
a989a71	2020-01-13	Tim Newsome	Clarify dmireset/dmihardreset. (#508)
d10d8d0	2020-01-06	Tim Newsome	Rebuild PDF.
efc0143	2020-01-06	Tim Newsome	Clarify action=1 (enter Debug Mode) with dmode=0 (#501)
439fb93	2020-01-06	Tim Newsome	Fix conflict in sbdata0/sbautoincrement definition. (#507)
d35ce10	2019-12-10	Tim Newsome	Add resume groups. (#506)
2726f30	2019-12-06	Tim Newsome	Rebuild PDF.
a310a37	2019-12-04	Tim Newsome	Make haltsum0 optional if there is only one hart. (#505)
349c826	2019-11-26	Tim Newsome	Halt state may not be preserved across reset. (#504)
4ab79d7	2019-11-26	Tim Newsome	Clear MPRV when resuming into lower privilege mode. (#503)
c9c286b	2019-11-22	Tim Newsome	Time may pass before dmactive becomes high. (#500)
9d55a57	2019-11-21	Megan Wachs	Make the emitted registers chisel3
014505f	2019-10-08	Tim Newsome	Rebuild PDF.
62c63b8	2019-10-04	Tim Newsome	Document forward progress guarantees in Debug Mode. (#496)
d933bec	2019-10-02	Tim Newsome	Rewrite/clarify DM Reset Control (#494)
039bd5a	2019-09-23	Philipp Wagner	Fix wrong table reference (#484)
106b4f2	2019-09-16	Tim Newsome	DM reset must also reset all the DM's harts. (#493)
8bfcd17	2019-09-13	Tim Newsome	Explicitly list cmderr=6 (reserved). (#491)
448de85	2019-09-12	Philipp Wagner	dmcontrol.hartreset is WARL, not R/W (#490)
8637b3c	2019-09-10	Tim Newsome	Rebuild PDF.
f00f436	2019-09-10	Philipp Wagner	Tiny style fix for email "link" on title page (#486)
3646788	2019-09-10	Philipp Wagner	Fix page references in cmdtype table (#487)
99ae160	2019-09-09	Megan Wachs	Update implementations.tex (#482)
f9c9ed4	2019-09-04	Philipp Wagner	Update registers.py to use Python 3 (#483)
37d8ee1	2019-09-03	Philipp Wagner	Git ignore intermediate and output files (#485)
1e99ce7	2019-08-13	Tim Newsome	Tighten up trigger specification. (#478)
a121ee1	2019-08-13	Tim Newsome	Rebuild PDF.
7d126a9	2019-07-16	Tim Newsome	Mention the scontext reg number isn't conventional (#474)
b5df5bd	2019-07-16	Tim Newsome	Explicitly document confstrptr[1-3]. (#475)

e6311af	2019-07-12	Tim Newsome	Change R/W1C to reduce requirements on hardware. (#472)
178e749	2019-07-11	Tim Newsome	Define what we mean by virtual address. (#473)
340c302	2019-07-09	Tim Newsome	Rebuild PDF.
77d58e6	2019-07-08	Tim Newsome	Numerous tweaks, responding to Marc Gauthier (#463)
ab89a86	2019-07-04	Tim Newsome	Addressing more feedback from Marc Gauthier. (#465)
624a6b8	2019-06-26	Tim Newsome	Without S-mode, textra.svalue and .sselect should be 0 (#469)
1977166	2019-06-11	Tim Newsome	Rebuild PDF.
b06eb70	2019-06-06	Tim Newsome	Clarify mcontrol.size. (#460)
165f120	2019-05-29	Tim Newsome	Fully qualify register/field macro names. (#457)
c47f0a0	2019-05-29	Paul Donahue	Fix #452 (#459)
633ee13	2019-05-28	Paul Donahue	Fixed #453 (#458)
96ef519	2019-05-20	Tim Newsome	The *external* debugger must restore tselect. (#456)
e11f777	2019-05-08	Tim Newsome	Rebuild PDF.
034d0d6	2019-04-30	Tim Newsome	Clarify that debuggers should honor maskmax. (#440)
4369eb8	2019-04-30	pdonahue-ventana	Finesse ligatures to work with Adobe Acrobat Reader search and cut-and-paste (#442)
d125b9b	2019-04-30	pdonahue-ventana	serror and sbusyerror don't both have to be non-zero to prevent (#447)
859e167	2019-04-30	Tim Newsome	Tweak address matches. (#449)
96b2b28	2019-04-25	Tim Newsome	Clarify not supported cmderr. (#446)
658417f	2019-04-16	Tim Newsome	When extending IR, BYPASS still is all ones. (#437)
2e24bab	2019-04-16	Tim Newsome	JTAG does not suggest any specific IDCODE encoding (#439)
c50efcb	2019-04-09	Tim Newsome	Rebuild PDF.
281e4ad	2019-03-21	Tim Newsome	Don't run text off a page when longtable is used. (#434)
76874e9	2019-03-20	Tim Newsome	Explain how to detect the version. (#433)
a543b76	2019-03-12	Tim Newsome	Rebuild PDF.
a686747	2019-02-21	Tim Newsome	All trigger registers are optional (#431)
d6e4cd8	2019-02-19	Josh Scheid	Fix typo. (#426)
e773936	2019-02-19	Tim Newsome	Try to get travis to build the release branch. (#430)
3621456	2019-02-19	Tim Newsome	Abstract memory accesses use the low bits of arg0. (#429)
94a5f9c	2019-02-12	Tim Newsome	Clarify that harts halt out of reset if haltreq=1 (#419)
518e732	2019-02-12	Tim Newsome	Rebuild PDF.
62f36e1	2019-02-11	Tim Newsome	Errata go in 0.13.x, this is 0.14. (#424)
66c3117	2019-01-31	Tim Newsome	Address triggers may fire on any accessed address. (#421)
6102412	2019-01-31	Tim Newsome	\Faamsize does not affect Argument Width. (#420)
1ea1a9b	2019-01-09	Tim Newsome	Add nmi bit to etrigger. (#408)
d1c7a3f	2019-01-09	Tim Newsome	Reserve trigger types for non-standard use. (#417)

83b12fb	2019-01-08	Tim Newsome	Rebuild PDF.
b4b3b5c	2019-01-07	Tim Newsome	\Fversion may be invalid when \Factive=0 (#414)
800450f	2019-01-01	Tim Newsome	mte only applies when action=0 (#411)
67c7fe2	2018-12-13	Tim Newsome	Add pre-built PDF of the 0.13 release.
5e7cb72	2018-12-12	Tim Newsome	Stopcount only applies to hart-local counters. (#405)
e5902fc	2018-12-12	Tim Newsome	Reserve some DMI space for non-standard use. (#406)
3c0dc6a	2018-12-11	Tim Newsome	Rebuild PDFs.
aeee8f3	2018-12-04	Tim Newsome	Add halt groups and external triggers. (#404)
814406d	2018-11-13	Tim Newsome	Clarify what the 4 states are. (#403)
cb64db0	2018-11-06	Tim Newsome	Rebuild PDFs.
70da60c	2018-11-05	Tim Newsome	sselect applies to svalue. (#402)
66fe38e	2018-11-05	Tim Newsome	Fix trigger example value. (#401)
688ccaf	2018-11-05	Tim Newsome	Resume ack is set after resume. (#400)
553dda7	2018-11-05	Tim Newsome	Fix sbdata0 read order of operations. (#392)
b864f54	2018-10-31	Tim Newsome	Add Compatibility section to the introduction. (#399)
0b205b1	2018-10-31	Tim Newsome	Create errata document. (#398)
5390063	2018-10-26	Tim Newsome	Bump version to 0.13.1. (#391)
e46c2db	2018-10-08	bdwyatt	Fix link to PDF (#387)
ed66f39	2018-10-02	Tim Newsome	Rebuild PDF.
f2873e7	2018-10-02	Tim Newsome	Run/Halt figure applies only to single-hart systems. (#385)
a79945f	2018-10-02	Tim Newsome	Add ASID and context compare for triggers (#363)
9bb7da6	2018-10-02	Tim Newsome	Clean up language of #383. (#384)
fce4da5	2018-10-02	Tim Newsome	Make haltreq and resumereq proper write-only. (#383)
e5da11e	2018-10-02	Tim Newsome	Minimal implementations can't access all registers (#381)
e1be8f4	2018-10-02	Tim Newsome	Format quotes correctly. (#382)
e9103ba	2018-10-02	Tim Newsome	Change from AVR debug connector to MIPI-10,20. (#375)
8841a7a	2018-10-02	Tim Newsome	Abstract reg access is independent of run/halt. (#380)
71c54bb	2018-10-02	Tim Newsome	Explicitly state what's required for compliance. (#379)
4edb285	2018-10-01	Tim Newsome	Rebuild PDF.
b0420b3	2018-10-01	Tim Newsome	Final cleanups! Mostly table formatting. (#377)
d43f5a4	2018-10-01	Tim Newsome	Clarify W1. (#372)
72618f3	2018-10-01	Tim Newsome	Leave space for trace, but don't specify anything. (#376)
b7db4ce	2018-10-01	Tim Newsome	Add dcsr.cause for being halted out of reset. (#370)
42ab2a1	2018-09-28	Tim Newsome	Clean up language, formatting, consistency. (#371)
7801874	2018-09-28	Tim Newsome	Little language and formatting cleanups. (#366)
38ae12f	2018-09-27	Tim Newsome	Reset dmi.op to 0 instead of 2. (#369)
b50dc0d	2018-09-27	Tim Newsome	Formatting, language, consistency. (#373)
425e9b1	2018-09-27	Tim Newsome	Distinguish draft and release builds. (#364)
c7b4e1c	2018-09-26	Tim Newsome	Stepping over wfi does not enter wait state. (#368)



4725879	2018-09-25	Tim Newsome	Language, formatting, and abstract cmd arguments. (#367)
62bf89d	2018-09-25	Tim Newsome	Rebuild PDF.
10dfa65	2018-09-24	Tim Newsome	Allow global reset to reset the DM. (#350)
84ec8a5	2018-09-18	Tim Newsome	Harts can be in exactly 1 of 4 states. (#354)
308eaf6	2018-09-17	Tim Newsome	Mostly match "official" style for credits. (#362)
b6187ff	2018-09-17	Tim Newsome	Specify ackhavereset as W1. (#361)
41d9f06	2018-09-14	Tim Newsome	Abstract commands might work on a hung hart. (#360)
fa561bd	2018-09-14	Tim Newsome	Can't change harts during operations, and the current hart becoming unavailable may terminate the abstract command with error. (#322)
900cdbf	2018-09-11	Tim Newsome	Rebuild PDF.
514ef6f	2018-09-07	Tim Newsome	Clarify lack of notification for other reset harts (#349)
e0ff31e	2018-09-07	Tim Newsome	Clarify postexec when there is no Program Buffer (#352)
3dacc00	2018-09-07	Florian Zaruba	Move regno table to the actual access reg command (#345)
5d25cd5	2018-09-06	Tim Newsome	don't set most bits of DMCONTROL during abstract commands (#324)
12655e0	2018-09-06	Tim Newsome	Document breakpoint exception + enter debug mode (#299)
6894f4b	2018-09-05	Tim Newsome	Define DXLEN as the widest supported XLEN. (#298)
114a208	2018-09-04	Tim Newsome	Restrict how many bits may be set in dmcontrol. (#348)
4cd1563	2018-09-03	Tim Newsome	Don't change selected harts during hart reset. (#337)
1529c26	2018-09-03	Tim Newsome	On trigger chains, only the last action is taken. (#341)
18a3531	2018-08-31	Tim Newsome	Authdata is bidirectional. (#347)
7d14f95	2018-08-27	Tommy Thorn	m "LaTeX/english issues: eg. -i, e.g., etc" (#342)
0fb41b9	2018-08-27	Tim Newsome	Don't change step/stepie while running. (#340)
ff09418	2018-08-21	Tim Newsome	Rebuild PDF.
6bd15ac	2018-08-20	Tim Newsome	Be more clear about running signal. (#338)
e967b3b	2018-08-20	Tim Newsome	mprven may be tied high or low. (#339)
0f120c0	2018-08-20	Tim Newsome	Solution to native triggers in M mode only systems (#309)
13d5c08	2018-08-17	Tim Newsome	Thank John Hauser.
b52d9fe	2018-08-17	Tim Newsome	Allow control xfers in progbuf to act as illegal. (#331)
19058ef	2018-08-17	Tim Newsome	Clarify that resumereq is not level-sensitive. (#321)
497352c	2018-08-16	Tim Newsome	Side effects happen for abstract register accesses (#334)
fd5cf62	2018-08-15	Tim Newsome	Triggers do not fire in Debug Mode. (#335)
762d308	2018-08-15	Tim Newsome	Add aarpostincrement to abstract register access. (#333)
45b7636	2018-08-14	Tim Newsome	Clearing hasel does not clear the ha mask reg. (#327)

2ca20aa	2018-08-13	Tim Newsome	clrresethaltreq trumps setresethaltreq (#332)
57df3f3	2018-08-10	Tim Newsome	\Rcommand is not readable. (#328)
81df032	2018-08-10	Tim Newsome	Explain what we mean by Preset. (#323)
b51c6db	2018-08-10	Tim Newsome	Clarify ebreak behavior when ebreak* are 0. (#311)
a14d868	2018-08-10	Tim Newsome	Allow extra harts to be reset. (#330)
6d60ad9	2018-08-07	Tim Newsome	Rebuild PDF
f4bd15f	2018-08-02	Tim Newsome	Define cmderr for non-existent register access. (#325)
2d7d3d0	2018-07-20	Tim Newsome	Fix typo in data0 definition.
c8a64d1	2018-07-19	Tim Newsome	Rebuild PDF.
9d2944f	2018-07-18	Tim Newsome	Add size to mcontrol. (#310)
6bd1a4c	2018-07-16	Tim Newsome	Put the description of dmstatus first. (#303)
25e81e5	2018-07-12	Tim Newsome	Fix typo in trigger example. (#308)
8462c94	2018-07-09	Tim Newsome	Rebuild pdf.
38fde94	2018-07-09	Tim Newsome	datacount cannot be 0 (#286)
800ca8d	2018-07-06	Tim Newsome	Clarifications requested by Jeremy Bennett (#280)
b363afa	2018-07-06	Tim Newsome	Add missing .tex file to dependencies. (#302)
93340e4	2018-07-06	Tim Newsome	Clarify that trigger registers are WARL. (#306)
95af58a	2018-07-06	Tim Newsome	Force the register-address in place. (#304)
d83039d	2018-07-06	Tim Newsome	\Fcause priority numbers: higher means higher (#307)
921c6a3	2018-07-03	Tim Newsome	Completing progbuf exec is I/O for fence insts. (#305)
99e01fa	2018-06-27	Tim Newsome	Add target-specific bits to abstract access memory. (#295)
4a0152d	2018-06-19	Tim Newsome	Only write busy to \Fcmderr if \Fcmderr is 0. (#296)
b0dc615	2018-06-16	Tim Newsome	Rebuild the PDF.
90873eb	2018-06-16	Tim Newsome	Fix typo in abstract access memory examples. (#297)
5fe8e08	2018-06-16	Tim Newsome	dret is a section, not a subsection of reset (#294)
abfd8a0	2018-06-14	Tim Newsome	Revert "Only write busy to \Fcmderr if \Fcmderr is 0."
7c66968	2018-06-14	Tim Newsome	Only write busy to \Fcmderr if \Fcmderr is 0.
0f28f27	2018-06-08	Tim Newsome	Abstract memory (#283)
7c840dd	2018-06-08	Tim Newsome	Specify an Exception Trigger (#266)
9d0d8af	2018-06-06	Tim Newsome	Clarify what address space these registers are in (#281)
a7f293d	2018-06-03	Tim Newsome	Add missing dependency to Makefile (#285)
37893aa	2018-05-30	Tim Newsome	Make trigger types writable. (#279)
6730cc0	2018-05-29	Tim Newsome	Explain priority assignment rationale. (#277)
b6d5d66	2018-05-25	Tim Newsome	Prevent M mode triggers affecting D mode ones (#282)
08ee84f	2018-05-22	Tim Newsome	Reading tselect doesn't guarantee a valid trigger. (#271)
6dfe375	2018-04-18	Megan Wachs	Debug Module should be capitalized
dac2120	2018-04-11	Megan Wachs	resethaltreq: Proposal for forcing a hart into debug mode out of reset

3b6442f	2018-05-16	Tim Newsome	tdata2 need only hold valid addresses if select=0 (#278)
68501cb	2018-04-26	mwachs5	mprven: Add a bit to enable MPRV to take effect in debug mode
9fcabe0	2018-05-03	Megan Wachs	Appendix: correct and clarify what debugger vs DM does
30773fd	2018-05-03	Tim Newsome	Debuggers must not write sbcs while sbbusy is set (#270)
50d8cd8	2018-05-03	Megan Wachs	Remove merge commits from the changelog
3b7a296	2018-05-02	Tim Newsome	Fix typo.
b26072b	2018-05-02	Tim Newsome	Explain that 1 in hart array mask means selected
41f6026	2018-05-02	Megan Wachs	Examples: Give an example of CSR access with Quick Access (#268)
675bb14	2018-05-01	Tim Newsome	Replace XLEN with MXLEN. #257
848cca1	2018-04-30	Megan Wachs	Overview Diagram: increase number of Progbuf words (#267)
a719ee6	2018-04-25	Megan Wachs	fix misspelled name
097c701	2018-04-23	Tim Newsome	Fix typo.
01dabd5	2018-04-23	Tim Newsome	Incorporate review feedback.
ca7a9d0	2018-04-18	Tim Newsome	Add trigger examples for match types 1, 4, and 5
cd5a15c	2018-04-16	Tim Newsome	Give a few trigger examples.
4375927	2018-04-12	Tim Newsome	Clarify that maskmax applies only to NAPOT trigger
acadfe9	2018-04-13	Megan Wachs	NMI: debugging may not be possible if an NMI happens
8fb190c	2018-04-12	Tim Newsome	Another attempt at SBA errors.
714c5d1	2018-04-11	Megan Wachs	Core Debug: all interrupts are masked includes NMI
56fbd9d	2018-04-11	Megan Wachs	DCSR: add nmip bit to indicate NMI is pending
ffe3c2	2018-04-10	Tim Newsome	Clarify SBA unsupported access size error.
b4006ac	2018-04-10	Tim Newsome	Clarify high bits of sbdata in narrow reads.
4ca83dd	2018-03-28	Tim Newsome	Clarify progbuf=1 some more
3b62243	2018-03-26	Tim Newsome	Clarify debugger requirements when progbufsize=1
ffba4d0	2018-03-26	Tim Newsome	Explain why progbufsize=1 is special
6b88905	2018-03-19	Megan Wachs	haltsum1: correct its address to be BWC and not overlap with ABSTRACTAUTO
2382e2e	2018-03-06	Megan Wachs	Correct some inaccuracies in the chisel generated files
3e88e11	2018-03-06	Megan Wachs	travis: add 'make chisel' target to regression
32cbb9b	2018-03-19	Tim Newsome	Nonexistent/unavailable harts are not halted.
f8a7bb7	2018-03-19	Tim Newsome	More clarification.
e21ae4c	2018-03-16	Tim Newsome	Allow any bit in hart array mask to be tied to 0
efb7e45	2018-03-15	Tim Newsome	Change dcsr.prv reset value to 3
f19946b	2018-03-15	Tim Newsome	Clarify hart array mask register size.
ddec145	2018-03-14	Tim Newsome	Be more precise about core vs hart
4e5f4ad	2018-03-14	Tim Newsome	Review feedback.
8ac9273	2018-03-14	Tim Newsome	Be more precise about processor vs hart
83c9774	2018-03-14	Tim Newsome	Clarify abstract command errors.
4ebc177	2018-03-14	Tim Newsome	hawindowssel can be smaller, depends on # of harts
11e1b5c	2018-03-14	Tim Newsome	Split future ideas section into a notes doc
bafeeaa	2018-03-13	Tim Newsome	Rebuild PDF

6a85d53	2018-03-13	Tim Newsome	Incorporate review feedback.
f213315	2018-03-09	Tim Newsome	Clarify user responsibilities when debugging lr/sc
3641305	2018-03-09	Tim Newsome	Remove implemented features from Future Ideas.
1135bf3	2018-03-06	Tim Newsome	Incorporate feedback.
8f35e7e	2018-03-05	Megan Wachs	gt_1024: Clarify that some registers may not be present for small numbers of harts
683ae37	2018-02-14	Megan Wachs	hartsum- $\hat{}$ haltsum
ee51758	2018-02-14	Megan Wachs	Modification of $\hat{}$ 1024 hart proposal that maintains backwards compatibility
370d222	2018-03-05	Tim Newsome	Rephrase description of hit bit.
eee5e0c	2018-03-05	Tim Newsome	Clarify multiple DMs/harts
4d5acef	2018-02-28	Tim Newsome	Clarify what happens when $\backslash$ Fauthenticated is clear
6a0c9ec	2018-02-27	Tim Newsome	Move hit bit per review feedback.
097bd8e	2018-02-21	Tim Newsome	Fix link to pre-built pdf
d21774b	2018-02-21	Omer Faruk IR-MAK	Python interpreter to be used should default to Python2
a8c10cf	2018-02-20	Tim Newsome	Incorporate review feedback.
a0f947c	2018-02-20	Tim Newsome	Make trigger hit bit optional.
77e4634	2018-02-08	Tim Newsome	Add hit bit to hardware triggers.
140390a	2018-02-05	Tim Newsome	Better wording.
e35b1ff	2018-02-05	Tim Newsome	Move Reg Access Abbrev table after sample register
e887433	2018-02-05	Tim Newsome	Use longtable instead of xtabular.
5c84437	2018-01-31	Tim Newsome	Abstract Command data usage depends on the command
3d508ea	2018-01-25	Tim Newsome	HARTSELBITS- $\hat{}$ HARTSELLEN and other feedback
eb653f7	2018-01-24	Tim Newsome	Be explicit about the size of $\backslash$ Fhartsel.
822bd81	2018-01-24	Tim Newsome	Revert incrementing version number.
4c755af	2018-01-24	Tim Newsome	$\backslash$ Fsbbusyerror also inhibits new accesses.
457413d	2018-01-24	Tim Newsome	Update how to enumerate all harts.
2180801	2018-01-18	Tim Newsome	Fix ambiguity in busy error reporting.
3140efa	2018-01-09	Tim Newsome	Re-apply e698a5001aa4583d31dde484d78f4f10e4e3148f . No need to list out all the consecutive registers.
390daa7	2018-01-18	mwachs5	sbadress: Only writes to address will actually cause an error. Reads while busy are permitted.
5c820f3	2018-01-18	Megan Wachs	Remove reference to "caches"
4533648	2018-01-18	Megan Wachs	correct access spelling
d37c1ac	2018-01-16	Tim Newsome	Fix table column overruns by going full manual
e9100ea	2018-01-16	Tim Newsome	Correct when sbbusy error is set for being busy.
c029cc7	2018-01-16	Tim Newsome	Complete partial sentence.
494338a	2018-01-15	Tim Newsome	Add clarifications about error handling.
e14c34e	2018-01-15	Tim Newsome	Incorporate review feedback.
68720e5	2018-01-15	Tim Newsome	Remove H bits from triggers.
b8eb62a	2018-01-15	Tim Newsome	Clarify when sbaccess is checked for validity
8b50d29	2018-01-12	Tim Newsome	Add $\backslash$ Fsbbusy, to avoid race clearing $\backslash$ Fsberror
50b1b41	2018-01-12	Tim Newsome	Clarify: writes to $\backslash$ Rsbdata0 write the new data
7f26759	2018-01-12	Tim Newsome	Clarify exactly which bits are used for SB access.
47a019c	2018-01-11	Tim Newsome	Fix typo.

a49d6ad	2018-01-11	Tim Newsome	sbreadonaddr is R/W
42195c2	2018-01-11	Tim Newsome	Fix cut-and-paste error.
6c95235	2018-01-11	Tim Newsome	Add sbaddress3, for future proofing.
e3345ea	2018-01-11	Tim Newsome	Incorporate review feedback.
6da48f8	2018-01-11	Tim Newsome	Remove dmerr.
e99c092	2018-01-10	Tim Newsome	Add system bus version field.
a6aa531	2018-01-10	Tim Newsome	Talk about all data and progbuf regs in first reg
af272db	2018-01-09	Megan Wachs	Update dret font
3d579d8	2018-01-09	Tim Newsome	Explicitly list data[1-10] and progbuf[1-15]
c6481ae	2018-01-09	Tim Newsome	Revert "Explicitly list data[1-10] and progbuf[1-15]"
e698a50	2018-01-09	Tim Newsome	Explicitly list data[1-10] and progbuf[1-15]
e547ed5	2018-01-09	Tim Newsome	Clarify that we deal in physical addresses only.
b377b89	2018-01-09	Tim Newsome	Revert "Clarify that we deal in physical addresses only."
f7da066	2018-01-09	Tim Newsome	Clarify that we deal in physical addresses only.
99a1599	2018-01-09	Tim Newsome	Clarify that \Fdatasize contains at most 12.
ae6e88a	2018-01-09	mwachs5	dret: Legal only in Debug Mode
18f392d	2017-11-24	Tim Newsome	Get rid of sbsingleread in favor of sbreadonaddr
5754a3b	2018-01-05	Megan Wachs	Use a different word than "clobbered"
aca7e0b	2018-01-03	Megan Wachs	Add missing "to"s to abstractauto description
d59ddf3	2018-01-03	Megan Wachs	Correct plurality of halted harts in haltsum
57c53ed	2017-12-22	Tim Newsome	Put parens around all macros that need it.
7ded846	2017-12-18	Tim Newsome	Refer to existing hart instead of "valid"
68b8ac8	2017-12-15	Tim Newsome	Make \Fhaltsel WARL.
6a72f45	2017-12-18	Tim Newsome	Mark this as a draft, which it is.
dd8d871	2017-12-18	Tim Newsome	Properly deal with \ chars in the changelog.
42f920c	2017-12-18	Tim Newsome	Deal with \ chars in the changelog.
b13891c	2017-12-15	Tim Newsome	Revert "Make \Fhaltsel WARL."
26d76a0	2017-12-15	Tim Newsome	Make \Fhaltsel WARL.
afda8d7	2017-11-28	mwachs5	update PDF
134d310	2017-11-28	Megan Wachs	Correct compressed version of ebreak
caa1258	2017-11-27	Megan Wachs	badaddr -i tval (Priv Spec 1.9 -i 1.9.1)
32b0f08	2017-11-22	Tim Newsome	Incorporate feedback.
2f7aa54	2017-11-22	Tim Newsome	Simplify, and explain trigger behavior.
3e5887f	2017-11-21	Tim Newsome	Clarify some single step corner cases.
f4b9ae2	2017-11-21	Tim Newsome	Make ackhavereset write-only. (#178)
efe3dc8	2017-11-21	Tim Newsome	Make hartreset R/W (#177)
ce1b359	2017-11-17	Megan Wachs	Reset clarifications (#172)
852a70d	2017-11-16	Megan Wachs	icount: remove warning (#173)
363348f	2017-11-16	Tim Newsome	Explain cache coherency wrt to system bus access (#171)
26ea898	2017-11-15	Tim Newsome	Refer to ISA and priv docs.
ffc8c62	2017-11-03	Tim Newsome	Mention the index in "about this doc"
a4257ef	2017-11-02	Tim Newsome	Add an index to the document.
f5f45a5	2017-10-30	Megan Wachs	Add 'has reset' status and control (#168)
46f3f54	2017-10-25	Tim Newsome	Incorporate review feedback.
104247f	2017-10-24	Megan Wachs	Update README.md
6dd5c80	2017-10-24	Megan Wachs	Update README.md

cb1a847	2017-10-24	Megan Wachs	Add a note to the README about the built PDF
e00625f	2017-10-18	Tim Newsome	Include pdf.
c23e729	2017-10-18	Tim Newsome	Clarify more.
83f9faf	2017-10-11	Tim Newsome	Clarify what <code>\Fimpebreak</code> does.
78082b5	2017-10-11	Tim Newsome	Mention <code>\Fimpebreak</code> in Program Buffer description.
0378324	2017-10-11	mwachs5	Add legend and update some transitions on the Abstract Command State Machine diagram
fa2b600	2017-10-11	Megan Wachs	add missing period
0610630	2017-10-11	Megan Wachs	Just do simple <code>hmode -i dmode</code> replacement
16e11f3	2017-10-11	Tim Newsome	Remove <code>hmode</code> reference, to fix build.
84b9a6a	2017-10-11	Tim Newsome	Add <code>\Fimpebreak</code> , to support of implicit <code>ebreak</code> .
cc90b77	2017-10-11	mwachs5	Remove reference to 'H' mode from the figure
cc6a9de	2017-10-11	Megan Wachs	Change old reference to 'hmode' to 'dmode'
ea2877d	2017-10-10	Tim Newsome	Move how-to-debug into the relevant section.
486ecc6	2017-10-05	Tim Newsome	Refuse unsupported bus accesses.
6ca221d	2017-10-05	Tim Newsome	<code>haltreq</code> , <code>resumereq</code> , <code>hartreset</code> are per-hart bits
d4118ab	2017-09-30	Tim Newsome	<code>ndmreset</code> can't reset logic required to access DM.
c6bd8d1	2017-09-29	Tim Newsome	and <code>-i</code> or
58c2441	2017-09-29	Tim Newsome	Mention <code>\Fstepie</code> in Single Step
94c5f78	2017-09-29	Tim Newsome	Clarify <code>ndmreset</code> .
12810b4	2017-09-29	Tim Newsome	Clarify that <code>sbaddress</code> is physical.
5862fdf	2017-09-29	Tim Newsome	Unify M mode and <code>mprv</code> comment.
aea1bd5	2017-09-29	Tim Newsome	Define behavior when <code>haltreq</code> and <code>resumereq</code> are set
146b348	2017-09-28	Megan Wachs	remove superflous 'an'
a5d16c4	2017-09-28	Megan Wachs	remove superfluous 'a'
052a8ab	2017-09-28	Tim Newsome	Clarify that a debugger can lose hart control.
cc52cff	2017-09-28	Tim Newsome	Add <code>\Fdmerr</code> .
25685eb	2017-09-28	Tim Newsome	Explain that bus master or <code>progbuf</code> is required.
f75ee7d	2017-09-28	Tim Newsome	Clarify debugger can discover "almost" everything
71e6788	2017-09-27	Tim Newsome	Remove description of manual stepping.
9aea347	2017-09-27	Tim Newsome	Move Running/Single Step near Halting.
2090d9b	2017-09-27	Tim Newsome	<code>data0</code> should be <code>sbdata0</code> in this table.
5858cfe	2017-09-27	Tim Newsome	Clarify why <code>\Rpriv</code> exists.
bc3c2aa	2017-09-27	Tim Newsome	Mention where <code>priv</code> encoding comes from.
ef77cc4	2017-09-27	Tim Newsome	One more attempt to clarify DPC after single step.
80a288e	2017-09-27	Tim Newsome	Clarify <code>instret</code> not incrementing on <code>ebreak</code> .
c163d22	2017-09-20	Tim Newsome	Remove <code>ebreakh</code> .
9971075	2017-09-20	Tim Newsome	Clarify we're talking about privilege
3fbe495	2017-09-20	Tim Newsome	Clarify that we're talking about *implementation*
3684854	2017-09-20	Tim Newsome	Use steps environment in <code>sbdata0</code> .
d4eda18	2017-09-20	Tim Newsome	Explain that only <code>sbdata0</code> has side effects.
ae781c6	2017-09-20	Tim Newsome	Don't refer to internal system bus registers.
875922e	2017-09-20	Tim Newsome	Explain <code>sbdata0</code> being stale a bit more.
cd44fd5	2017-09-20	Tim Newsome	Clarify <code>autoread</code>
194484b	2017-09-20	Tim Newsome	Clarify <code>hawindow</code> .
02f1aac	2017-09-20	Tim Newsome	Clarify that <code>\Fdataaddr</code> is relative to <code>\Rzero</code> .
0e9b6ae	2017-09-20	Tim Newsome	Clarify nonexistent vs unavailable.
b55ff41	2017-09-20	Tim Newsome	Fix <code>devtreevalid</code> .

2eccb86	2017-09-20	Tim Newsome	Explicitly state which registers are read-only.
4af505c	2017-09-20	Tim Newsome	Show section numbers for registers.
cbd5573	2017-09-20	Tim Newsome	Thank Nikhil
19c206f	2017-09-20	Tim Newsome	Clarify how to determine whether progbuf is RAM
0651f7d	2017-09-20	Tim Newsome	Explain what happens if ebreak is missing.
e889dae	2017-09-20	Tim Newsome	Move figure of states into its own section.
cff7b80	2017-09-20	Tim Newsome	Explain when <code>\Ftransfer</code> might be used.
6b2ee61	2017-09-20	Tim Newsome	Explain where <code>\Fsize</code> encoding came from.
c9f3b73	2017-09-14	Tim Newsome	Fix typo.
4b25400	2017-09-13	Tim Newsome	Mention <code>dpc</code> in CSRs abstract register numbers.
c3ee426	2017-09-13	Tim Newsome	Move abstract regno table closer to its reference.
111b9a3	2017-09-13	Tim Newsome	cycle <code>-i</code> operation
994afdc	2017-09-13	Tim Newsome	Account for multiple selected harts.
aa4a297	2017-09-13	Tim Newsome	Halt Control <code>-i</code> Run Control
e97c821	2017-09-13	Tim Newsome	continuous <code>-i</code> contiguous
97f73ff	2017-09-13	Tim Newsome	Clarify <code>ndmreset</code> behavior.
6078220	2017-09-13	Tim Newsome	Explain <code>ndmreset</code>
a3d4f30	2017-09-13	Tim Newsome	Describe ‘halt region’
272b3d9	2017-09-13	Tim Newsome	Clarify accessing unimplemented DM DMI regs
3e91f1b	2017-09-13	Tim Newsome	Clarify either Prog Buf or Sys Bus Acc is required
e8a6145	2017-09-13	Tim Newsome	Clarify CSR access; remove serial port
ce20766	2017-09-13	Tim Newsome	Remove section referencing itself.
1195a61	2017-09-18	Tim Newsome	Generate constants to be unsigned for clang.
8967b0a	2017-08-16	Megan Wachs	Compressed instructions are <code>c.foo</code> , not <code>foo.c</code>
b5698a9	2017-08-16	Megan Wachs	clarify <code>progbufsize</code> description
d221bab	2017-08-16	Megan Wachs	Remove <code>progbufsize</code> enums from register description
0498102	2017-08-16	Megan Wachs	appendix: Use standard assembly format for sw
4456d99	2017-08-09	Tim Newsome	Rename <code>progsz</code> to <code>progbufsize</code> .
55d5b66	2017-08-09	Tim Newsome	Clarify that trigger comparisons are unsigned.
21e35ef	2017-08-09	Tim Newsome	Configuration String <code>-i</code> Device Tree
f044f45	2017-08-02	Tim Newsome	Don’t require a target to provide 25mA on VCC.
c883943	2017-08-02	Tim Newsome	Add table of Abstract Command Types
985a3df	2017-08-02	Tim Newsome	Fix and speed up build.
95b9108	2017-08-02	mwachs5	DTM: Clarify that there are no cases when DMI would actually return an error.
9c9e0c0	2017-08-02	mwachs5	SystemBus: No longer returns error. So DMI has no ‘error’ return code.
5ba18f9	2017-07-27	Tim Newsome	Fix more typos.
dbc65bf	2017-07-26	Tim Newsome	Fix typos.
bba0ad9	2017-07-26	Tim Newsome	Tighten up introduction lists.
e22d5eb	2017-07-26	Tim Newsome	Add version constants for ”not compatible”.
c79038e	2017-07-26	Tim Newsome	Small clarification.
9df0411	2017-07-21	Tim Newsome	Incorporate review feedback.
d67419c	2017-07-21	Tim Newsome	Clarify <code>dpc</code> contents.
9f50c05	2017-07-11	Tim Newsome	Use LL instead of L for 64-bit constant suffix.
23fd24a	2017-07-10	Megan Wachs	Cleaning up whitespaces
c5ab04c	2017-07-10	Megan Wachs	Update <code>abstract_commands.xml</code>
6e8cdf1	2017-07-10	Megan Wachs	Update <code>abstract_commands.xml</code>

cf6e3f2	2017-07-10	Megan Wachs	clarify DCSR.cause
79ffb9	2017-07-10	Megan Wachs	Clarify implications of CSR read, write, halt
013e191	2017-07-10	Megan Wachs	Clarify when you would get error halt/resume
231e457	2017-07-10	Megan Wachs	Quick Access error clarification
c54c2f2	2017-07-03	mwachs5	serial: add the XML file, not the TEX file
ac77477	2017-07-03	mwachs5	serial: Fix compile errors after moving serial port to appendix
6defcb8	2017-07-03	mwachs5	serial: Move serial ports out of main spec and into Future Work appendix
a28f639	2017-06-30	mwachs5	remove trace dependencies from Makefile
52a122b	2017-06-30	mwachs5	remove trace section
d9e166b	2017-06-30	mwachs5	remove trace registers
7caf4e5	2017-06-30	mwachs5	remove trace appendix
4688988	2017-06-29	mwachs5	DCSR: define a 'stepie' bit which may be hard-wired to 0.
9a0492c	2017-06-13	Megan Wachs	Add missing period and some other small text edits
13ccdbf	2017-06-13	Megan Wachs	fix typo in ProgBuf register macro
b01f989	2017-06-13	mwachs5	implementations: be a bit more concrete about the one example implementation we have.
a7b5f83	2017-06-13	mwachs5	jtagdtm: Move it out of the appendix as it is really part of the specification
87aceb0	2017-06-13	Megan Wachs	remove "spontaneous"
50b9950	2017-06-13	Megan Wachs	Forward reference for anynonexistent
adea3e2	2017-06-13	Megan Wachs	More clarifications on dret
1b8dd0e	2017-06-13	Megan Wachs	Define DRET instruction
617da4c	2017-06-08	Megan Wachs	Update description of R/W1C
de2c56b	2017-06-08	Megan Wachs	Clarify that DCSR is also not updated on ebreak
efa615d	2017-06-07	Tim Newsome	Increase xdebugver field size to 4 bits. (#92)
a0e147a	2017-06-07	Tim Newsome	Address some review comments.
89ffe50	2017-06-06	mwachs5	NDMRESET: Clarify what it may and may not do
1932da0	2017-06-06	mwachs5	DPC: Clarifications on its meaning
6470fdb	2017-06-06	mwachs5	ABSTRACTCS: Correct inconsistency on the number of data words.
3ca82b4	2017-06-06	Megan Wachs	More corrections for R vs R/W1C on SERCS
9705fb8	2017-06-06	Megan Wachs	Correct a bunch of W0 registers
1347371	2017-06-05	Tim Newsome	Add intdisable to dcsr.
989c60d	2017-06-05	Tim Newsome	Fix language. We can only halt harts, not cores.
517a08b	2017-06-05	Tim Newsome	Incorporate review feedback.
802be28	2017-06-05	Tim Newsome	Clarify/fix Quick Access example.
b8cc523	2017-06-02	Tim Newsome	Add included tex files as dependencies. (#78)
15f864a	2017-06-01	Tim Newsome	Language cleanups, consistency and typo fixes.
4ecae86	2017-06-01	Tim Newsome	Add page numbers to list-of-register tables.
59b3e4a	2017-05-19	Megan Wachs	Setting up a Travis regression to check for build errors (#72)
124bf44	2017-05-17	mwachs5	Debug Module: CMDERR is Write-1-to clear, not R/W0
bb6c7f0	2017-05-17	mwachs5	SW Registers file should be XML, not TEX



d360358	2017-05-10	Megan Wachs	Remove virtual register from core_registers.xml (Temporary Acct.)
bfc64fb	2017-05-10	Megan Wachs	Add missing sw_registers.tex file (Temporary Acct.)
0512f5d	2017-05-06	mwachs5	Move virtual 'prv' register to a separate section to make it more clear it is not a real register.
6b3c9d7	2017-05-06	mwachs5	Clarify haltreq/resumereq/resumack
0a487eb	2017-04-26	mwachs5	jtag: Change specified JTAG pinout from Coretex to AVR, to provide for TRSTn option.
93cdfaf	2017-04-26	mwachs5	DM : Clarify that DATA/PROGBUF can't be written while busy.
ef98f23	2017-04-19	mwachs5	jtag: Make it clear that a NOP is really a NOP.
a6f8efa	2017-04-17	mwachs5	single_step: Exceptions count as the 'step' completion.
bf11e9e	2017-04-17	mwachs5	resumeack: fix some LaTeX cross references
4afa081	2017-04-11	mwachs5	halt/resumereq: Clarify what setting them to 0 or 1 does
297a39b	2017-04-06	mwachs5	fix chisel build
082c499	2017-04-06	mwachs5	Rename resumed to resumeack, and add more text about what these bits mean.
909d617	2017-04-06	mwachs5	Correct some cross references after removing all the multiply listed registers
dd09914	2017-04-06	mwachs5	Add 'resumedall' and 'resumedany' bits to avoid race condition on about to resume and just halted
feb88fc	2017-04-05	mwachs5	JTAG DTM: Clarify that leading bits are 0 for more than 5-bit IR
75b96ea	2017-04-04	mwachs5	use renamed dm_registers file
9f3ec7e	2017-04-04	mwachs5	debugger_implementation: remove some old TODO and commentary.
45dd5b5	2017-04-04	mwachs5	Don't list out every single DM register for those that are just indexed versions
b8b3aa2	2017-04-04	mwachs5	remove core-side register definitions from Debug Module. Rename dm1 to dm
d979a13	2017-04-04	mwachs5	remove core-side serial port specification, as these should look like implementation-specific devices with appropriate drivers.
b56870b	2017-04-04	mwachs5	Remove the wording about 'debug exception', as it is called breakpoint exception in the RISC-V Spec.
1e9347d	2017-04-03	mwachs5	Add description of hasel
0dda84d	2017-04-03	mwachs5	JTAG DTM: Clean up TAP register descriptions
82ccde5	2017-04-03	mwachs5	JTAG DTM: Add a hard DMI bit which cancels the outstanding DMI transaction
bd2a3d1	2017-04-03	mwachs5	remove preexec
02c733a	2017-04-03	mwachs5	remove preexec from Abstract State diagram.
1e271d6	2017-04-03	mwachs5	Update Debugger implementation for DMI register access, and fix tex compile issues.

155dda4	2017-04-03	mwachs5	Rewrite HW Implementation examples to describe a pure abstract command approach, and to not rely on harts executing every instruction which is fetched from the Debug Module
556c2be	2017-04-03	mwachs5	minor wording edits about RISC-V core registers
523c64a	2017-04-03	mwachs5	Edits to the Debug Module section.
b9a371f	2017-04-03	mwachs5	add missing trace.tex file.
58b2396	2017-04-03	mwachs5	Re-order the JTAG DTM Sections
a8827e2	2017-04-03	mwachs5	Edits to the System Overview.
c5417ce	2017-04-03	mwachs5	add more sections as separate files.
287d5c6	2017-04-03	mwachs5	moving more files to separate tex files.
9e873f4	2017-04-03	mwachs5	move trigger info into separate file.
2c89a86	2017-04-03	mwachs5	move risc-v core debug info into separate file.
e676491	2017-04-03	mwachs5	Move System Overview to separate file
03df6ee	2017-04-03	mwachs5	Move Debug Module description to a separate file.
5faa430	2017-04-03	mwachs5	add back in JTAG DTM in appendix
7b28b11	2017-04-03	mwachs5	Move jtag DTM to appendix. Move some text to commentary.
cc183ba	2017-04-03	mwachs5	move introduction to a separate file. Comment out reading order.
f727d14	2017-04-03	mwachs5	Use Chapters vs Sections. Needs reorganization.
815951d	2017-04-03	mwachs5	Formatting updates. Make this look more like the RISC-V specs. Need to use chapter vs. section
69ffaf8	2017-03-31	mwachs5	Move XML files into a subdirectory.
b276384	2017-03-31	mwachs5	Remove debug_rom.S
112bbac	2017-03-31	mwachs5	figures: reorganize the figures into directories.
1e5c068	2017-03-27	Megan Wachs	Add LICENSE
fc17730	2017-03-22	Po-wei Huang	Change some halt mode into debug mode.
8ccf029	2017-03-22	Po-wei Huang	All halt mode changed to debug mode to synchronize with the priv spec.
f143d9e	2017-03-21	mwachs5	Correct duplicated progbuf register names
0797ec1	2017-03-17	mwachs5	autoexec: make autoexec bits match the number of data words there really are.
8e76d93	2017-03-17	mwachs5	dm1_registers: move a few more things around. Reduce abstract data words back to 12.
f8bf292	2017-03-17	mwachs5	dm1_registers: resolve some address conflicts and inconsistencies
a74dff9	2017-03-17	mwachs5	access_register: some small bit changes
2e6b0ca	2017-03-15	mwachs5	config string: Fix LaTeX compile errors.
f83260a	2017-03-10	mwachs5	Abstract Commands: clarify that 32-bit reads should always work. This allows reading MISA.
6f9347a	2017-03-10	mwachs5	Config String: change the Abstract Command to DMI registers. Allow the same registers to be used for unspecified identifier information.
4ea10ff	2017-03-10	mwachs5	abstract: Make autoexec apply to all data and progbuf words. Make a separate register which is optional.

5008436	2017-03-10	mwachs5	abstract: Allow up to 16 progbuf and/or data words. Inform debugger about dscratch registers available for its use.
aaa13e5	2017-03-06	mwachs5	Command: use the name 'cmdtype' not 'type' to allow easier auto-generation of Scala code.
e9bb72c	2017-03-06	mwachs5	Hart Array: Add registers for hart array.
5d17a35	2017-03-06	mwachs5	DM: Move addresses around for better separation of functionalities in HW
25ccaa8	2017-03-06	mwachs5	CONTROL: Rename control and status registers to ___CS for consistency and to accurately reflect their functionality.
45cf6c2	2017-03-06	mwachs5	Errors: fix up the bit assignments in SERSTATUS with the addition of error bit.
38cb5a0	2017-03-06	mwachs5	Errors: Make errors write-1-to-clear.
b436d77	2017-03-03	mwachs5	triggers: Clarify that matches are against virtual addresses.
793bb85	2017-03-03	mwachs5	triggers: Add suggested timings for best user experience.
2669866	2017-03-03	mwachs5	stoptime/stopcycle: Make their functionality match their name. Allow any reset value.
c85a1cf	2017-03-01	mwachs5	config_string: Simplify the Config String Address abstract command.
a303a6b	2017-03-02	Megan Wachs	Update README.md
92a4923	2017-03-01	mwachs5	serial: tweak addresses.
b09f460	2017-03-01	mwachs5	serial: tweak addresses.
6477837	2017-03-01	mwachs5	chisel: tweaks to class names.
be83e3e	2017-02-28	Tim Newsome	Clarify stoptime, stopcycle.
c17c17c	2017-02-27	Tim Newsome	Abstract command that returns config string addr.
096dfbc	2017-02-27	Tim Newsome	Acknowledge Alex.
c0253ab	2017-02-24	Tim Newsome	Explain tdata1 type a bit more.
e43ac2e	2017-02-24	Tim Newsome	Clarify how to enumerate triggers again.
c6e3e20	2017-02-23	Tim Newsome	Revert previous commit.
ef770bf	2017-02-23	Tim Newsome	mcontrol and icount mask tdata2, not tdata1.
27806f2	2017-02-23	mwachs5	rename 'type' to 'cmdtype' purely so my auto-generation scripts work.
e46798d	2017-02-22	mwachs5	Add Abstract Commands to automatic chisel
b3bb939	2017-02-21	mwachs5	Generate Chisel headers as well for Debug Module.
c9db98c	2017-02-22	Tim Newsome	Simplify description of op statuses.
bda39cc	2017-02-22	mwachs5	Add explicit type field to Abstract Command.
f83a1ca	2017-02-22	mwachs5	Finish up replacement of ibuf- $\bar{}$ progbuf
9666e51	2017-02-22	mwachs5	IBUF- $\bar{}$ PROGBUF
5308ecd	2017-02-22	mwachs5	Remove last references to "Instruction Supply"
f6ebde9	2017-02-22	Tim Newsome	Move authentication to a serial protocol.
0f079c8	2017-02-22	Tim Newsome	Reserve bit for per-hart reset.
f2c93ac	2017-02-22	Tim Newsome	Clarify that dmactive resets authentication.
f5e7b1c	2017-02-22	Alex Bradbury	Clarify that the halt state of all harts is maintained through reset
3dfe8fd	2017-02-22	Tim Newsome	More Debug Mode - $\bar{}$ Halt Mode.

d29fc1f	2017-02-22	Tim Newsome	Debug Mode - $\bar{i}$ Halt Mode
55d6030	2017-02-21	Tim Newsome	Generate debug_defines.h as part of normal make
b0e6a7f	2017-02-21	Tim Newsome	Minor clarifications.
0f9885c	2017-02-20	Tim Newsome	Various clarifications.
0802d5a	2017-02-15	mwachs5	Use consistent 'Control and Status' naming for CS registers.
5acc7d	2017-02-15	Tim Newsome	Change all the "other" JTAG IRs to just reserved.
bcbd7da	2017-02-15	mwachs5	sm_diagram: Show using resumereq bit to resume.
18f6e55	2017-02-14	Tim Newsome	Introduce resumereq command, similar to haltreq.
4b62c40	2017-02-14	mwachs5	SystemBus: Clean up some formatting and error specification notes.
bc97723	2017-02-14	mwachs5	quick-access: Update SM Diagram for Quick Access
d27066e	2017-02-14	Tim Newsome	Clarify haltreq bit.
6f8ec43	2017-02-14	Tim Newsome	Always generate long constants when required.
c6ac6bc	2017-02-13	Tim Newsome	Include field descriptions in C header file.
b849213	2017-02-13	Tim Newsome	Fix the build.
1cf8033	2017-02-12	mwachs5	jtag: More clarifications
6203bd6	2017-02-12	Megan Wachs	Update requirements- W GPRs Required
f2b43a7	2017-02-12	Megan Wachs	Remove double 'the'
2c64ef1	2017-02-12	Megan Wachs	Remove comma
f84abce	2017-02-12	Megan Wachs	Whitespace edits and address come comments
23c2648	2017-02-11	mwachs5	jtag-dtm: ask for clarification on TAP sharing.
7020d23	2017-02-11	mwachs5	jtag-dtm: Clarifications, DBUS- $\bar{i}$ DMI
292d49c	2017-02-11	Megan Wachs	fix indentation
b879b86	2017-02-11	Megan Wachs	Add missing period
bbe0521	2017-02-11	mwachs5	Make comments on program buffer size match the address map.
4ceaa37	2017-02-11	mwachs5	Flesh out and edit the introduction/background Add a description of use cases this spec has in mind, and what it doesn't cover.
cbf89d6	2017-02-11	Tim Newsome	Rewrite Quick Access.
170bff1	2017-02-10	Megan Wachs	Allow size 4 for the program buffer
c911e6e	2017-02-10	Tim Newsome	Clarify use of dmactive.
2ca296f	2017-02-09	Tim Newsome	Reserve command register space for custom use.
e49666e	2017-02-09	Tim Newsome	Clarify hart index change per Megan's comments.
84865e9	2017-02-09	Tim Newsome	Add header prefix for abstract commands.
2434f4f	2017-02-09	Tim Newsome	Select harts by index instead of hart ID.
7bf112a	2017-02-09	Tim Newsome	Generate correct headers for $\bar{i}$ 32-bit registers.
7f0f09a	2017-02-08	Tim Newsome	Reset dbus status to "failure" to avoid confusion.
8b1c6f0	2017-02-08	Megan Wachs	Fix line wrap issue
345c33f	2017-02-08	Megan Wachs	Call out "arg0" specifically.
9f080f5	2017-02-08	Megan Wachs	Clarify "arguments" to commands
259badd	2017-02-08	Tim Newsome	Make haltsum/halt registers mandatory.
eb0f1d3	2017-02-07	Tim Newsome	Allow for early abstract command failures.
bb49bd1	2017-02-07	Tim Newsome	Clarify error handling a little.
3fc0a97	2017-02-07	Tim Newsome	Explain when abstract data regs may be clobbered.
c37167e	2017-02-07	Tim Newsome	Fix old language in description of halt registers.
6943c96	2017-02-07	Tim Newsome	Generate more useful C header files from reg defs

98639df	2017-02-05	mwachs5	Include the SM Diagram as a figure. Also some minor capitalization fixes.
a95e4c3	2017-02-05	mwachs5	Update State Machine diagram to show uncertainty of halt bit during auto halt/resume.
ba76744	2017-02-05	Tim Newsome	Combine loabits and hiabits.
02b1d92	2017-02-05	Tim Newsome	DMI can get away with just 6 address bits.
35d6e33	2017-02-05	mwachs5	Update State machine diagram to show BUSY without HALTED
f511b05	2017-02-04	Tim Newsome	Clarify command busy bit.
d0f8961	2017-02-03	mwachs5	Update figures
e18a68d	2017-02-03	Tim Newsome	Clarify prehalt/postresume failure.
ac3e2a9	2017-02-02	Tim Newsome	Clarify abstract command failure behavior.
ce4baee	2017-02-02	Tim Newsome	Add Quick Access section.
0490377	2017-02-02	Tim Newsome	Add prehalt and postresume to reg command.
67515bd	2017-02-02	Tim Newsome	Deal with a few minor TODOs.
96456fc	2017-02-02	Tim Newsome	Turn register names into links.
317cd98	2017-02-02	Tim Newsome	Explain what register access is required.
f3ad2f2	2017-02-01	Tim Newsome	Revert Plain Exception implementation to be simple
a0ad281	2017-02-01	Tim Newsome	execb -i preexec, execa -i postexec
1d4a2c3	2017-02-01	Tim Newsome	Limit Program Buffer sizes to 0, 1, 8.
cc40815	2017-02-01	Tim Newsome	Incorporate Po-wei's feedback.
c8b45d6	2017-02-01	Tim Newsome	Clarify how all autoexec bits work.
dbb1deb	2017-02-01	Tim Newsome	Remove stale TODO.
c5f8f59	2017-02-01	Tim Newsome	Explain why cmderr inhibits starting new commands.
5c69194	2017-02-01	Tim Newsome	Fix editing error.
50f7c48	2017-02-01	Tim Newsome	Remove empty hart info register.
781c68e	2017-02-01	Megan Wachs	Update README.md
f46b32e	2017-02-01	mwachs5	Add a diagram of Abstract Command flow.
633bd63	2017-02-01	Tim Newsome	Move Reading Order into About This Document
51ec4d1	2017-02-01	Tim Newsome	Add reading order section.
03d20ad	2017-02-01	Tim Newsome	autoexec0 applies to data0, not inst0.
c302353	2017-01-31	Tim Newsome	Don't rely on hart fetching instructions once.
2558c25	2017-01-31	Tim Newsome	Change how exceptions in Halt Mode are handled.
a36ddce	2017-01-31	Tim Newsome	Add size to abstract register command.
64de458	2017-01-31	Tim Newsome	Detail bus master reads.
c08486f	2017-01-31	Megan Wachs	reset: Add some comments (#5)
1558049	2017-01-30	Tim Newsome	Automate Change Log.
51525a4	2017-01-29	Tim Newsome	Update System Overview
7d39ac0	2017-01-29	Tim Newsome	Update Supported Features.
9e7cbea	2017-01-29	Tim Newsome	Update RISC-V Core section.
515188d	2017-01-29	Tim Newsome	Update Hardware Implementations section.
4b19ed8	2017-01-29	mwachs5	system_bus: be consistent and always call it 'System Bus'. Even if some dislike the name, we should be consistent and clear in the spec.
9cccf3d	2017-01-29	Tim Newsome	Fleshed out some debugger implementation.
04b9176	2017-01-28	Tim Newsome	Rename debug exception to breakpoint exception.
5ac4ea1	2017-01-27	Tim Newsome	WIP on big update on instruction supply.
2d9c3e2	2017-01-27	Tim Newsome	Reorganize dm registers.

de50ba8	2017-01-27	Tim Newsome	Abstract command support is already addressed.
5085046	2017-01-26	mwachs5	Rename registers and fields like 'access' that were confusingly the same name.
10bbf6f	2017-01-26	Tim Newsome	Fix #2: DM address space table
a05c582	2017-01-26	Tim Newsome	Add debugger inspection as a feature.
4062681	2017-01-24	Tim Newsome	Add publish target.
5c8bb83	2017-01-24	Tim Newsome	Clarify use of data registers.
1504da6	2017-01-24	Tim Newsome	Replace manual date with automatic git hash/date.
997f2a0	2017-01-23	Tim Newsome	Deal with unsupported abstract commands.
cb6f2b8	2017-01-23	Tim Newsome	Reorder registers to prevent duplicates.
8b4db96	2017-01-23	Tim Newsome	Don't print out addresses if they're not provided.
b00cd21	2017-01-23	Tim Newsome	Add an abstract command.
675b556	2017-01-23	Tim Newsome	Reorganize DM bits into functional group regs.
5fc7512	2017-01-23	Tim Newsome	Remove bits 33:32 from sbdata[23].
ceb5d66	2017-01-20	Tim Newsome	Starting point for a comprehensive spec