



RISC-V Integer Conditional (Zicond) operations extension

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Preamble



This document is in the [Ratified state](#)

No changes are allowed. Any desired or needed changes can be the subject of a follow-on new extension. Ratified extensions are never revised

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Chapter 1. Introduction

The Zicond extension defines a simple solution that provides most of the benefit and all of the flexibility one would desire to support conditional arithmetic and conditional-select/move operations, while remaining true to the RISC-V design philosophy. The instructions follow the format for R-type instructions with 3 operands (i.e., 2 source operands and 1 destination operand). Using these instructions, branchless sequences can be implemented (typically in two-instruction sequences) without the need for instruction fusion, special provisions during the decoding of architectural instructions, or other microarchitectural provisions.

1.1. Motivation and use cases

One of the shortcomings of RISC-V, compared to competing instruction set architectures, is the absence of conditional operations to support branchless code-generation: this includes conditional arithmetic, conditional select and conditional move operations. The design principles of RISC-V (e.g. the absence of an instruction-format that supports 3 source registers and an output register) make it unlikely that direct equivalents of the competing instructions will be introduced.

Yet, low-cost conditional instructions are a desirable feature as they allow the replacement of branches in a broad range of suitable situations (whether the branch turns out to be unpredictable or predictable) so as to reduce the capacity and aliasing pressures on BTBs and branch predictors, and to allow for longer basic blocks (for both the hardware and the compiler to work with).

Chapter 2. Zicond specification

The "Conditional" operations extension provides a simple solution that provides most of the benefit and all of the flexibility one would desire to support conditional arithmetic and conditional-select/move operations, while remaining true to the RISC-V design philosophy. The instructions follow the format for R-type instructions with 3 operands (i.e., 2 source operands and 1 destination operand). Using these instructions, branchless sequences can be implemented (typically in two-instruction sequences) without the need for instruction fusion, special provisions during the decoding of architectural instructions, or other microarchitectural provisions.

The following instructions comprise the Zicond extension:

RV32	RV64	Mnemonic	Instruction
✓	✓	czero.eqz rd, rs1, rs2	Conditional zero, if condition is equal to zero
✓	✓	czero.nez rd, rs1, rs2	Conditional zero, if condition is nonzero



Architecture Comment: defining additional comparisons, in addition to equal-to-zero and not-equal-to-zero, does not offer a benefit due to the lack of immediates or an additional register operand that the comparison takes place against.

Based on these two instructions, synthetic instructions (i.e., short instruction sequences) for the following **conditional arithmetic** operations are supported:

- conditional add, if zero
- conditional add, if non-zero
- conditional subtract, if zero
- conditional subtract, if non-zero
- conditional bitwise-and, if zero
- conditional bitwise-and, if non-zero
- conditional bitwise-or, if zero
- conditional bitwise-or, if non-zero
- conditional bitwise-xor, if zero
- conditional bitwise-xor, if non-zero

Additionally, the following **conditional select** instructions are supported:

- conditional-select, if zero
- conditional-select, if non-zero

More complex conditions, such as comparisons against immediates, registers, single-bit tests, comparisons against ranges, etc. can be realized by composing these new instructions with existing instructions.

Chapter 3. Instructions (in alphabetical order)

3.1. czero.eqz

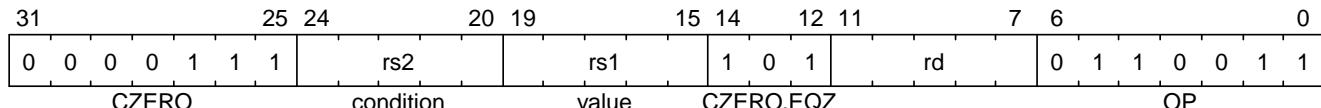
Synopsis

Moves zero to a register *rd*, if the condition *rs2* is equal to zero, otherwise moves *rs1* to *rd*.

Mnemonic

`czero.eqz rd, rs1, rs2`

Encoding



Description

If *rs2* contains the value zero, this instruction writes the value zero to *rd*. Otherwise, this instruction copies the contents of *rs1* to *rd*.

This instruction carries a syntactic dependency from both *rs1* and *rs2* to *rd*. Furthermore, if the Zkt extension is implemented, this instruction's timing is independent of the data values in *rs1* and *rs2*.

SAIL code

```
let condition = X(rs2);
result : xlenbits = if (condition == zeros()) then zeros()
                     else X(rs1);
X(rd) = result;
```

3.2. czero.nez

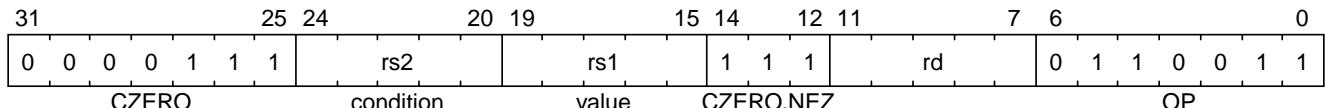
Synopsis

Moves zero to a register *rd*, if the condition *rs2* is nonzero, otherwise moves *rs1* to *rd*.

Mnemonic

`czero.nez rd, rs1, rs2`

Encoding



Description

If *rs2* contains a nonzero value, this instruction writes the value zero to *rd*. Otherwise, this instruction copies the contents of *rs1* to *rd*.

This instruction carries a syntactic dependency from both *rs1* and *rs2* to *rd*. Furthermore, if the Zkt extension is implemented, this instruction's timing is independent of the data values in *rs1* and *rs2*.

SAIL code

```
let condition = X(rs2);
result : xlenbits = if (condition != zeros()) then zeros()
                     else X(rs1);
X(rd) = result;
```

Chapter 4. Usage examples

The instructions from this extension can be used to construct sequences that perform conditional-arithmetic, conditional-bitwise-logical, and conditional-select operations.

4.1. Instruction sequences

Operation	Instruction sequence	Length
Conditional add, if zero $rd = (rc == 0) ? (rs1 + rs2) : rs1$	czero.nez rd, rs2, rc add rd, rs1, rd	
Conditional add, if non-zero $rd = (rc != 0) ? (rs1 + rs2) : rs1$	czero.eqz rd, rs2, rc add rd, rs1, rd	
Conditional subtract, if zero $rd = (rc == 0) ? (rs1 - rs2) : rs1$	czero.nez rd, rs2, rc sub rd, rs1, rd	
Conditional subtract, if non-zero $rd = (rc != 0) ? (rs1 - rs2) : rs1$	czero.eqz rd, rs2, rc sub rd, rs1, rd	
Conditional bitwise-or, if zero $rd = (rc == 0) ? (rs1 rs2) : rs1$	czero.nez rd, rs2, rc or rd, rs1, rd	2 insns
Conditional bitwise-or, if non-zero $rd = (rc != 0) ? (rs1 rs2) : rs1$	czero.eqz rd, rs2, rc or rd, rs1, rd	
Conditional bitwise-xor, if zero $rd = (rc == 0) ? (rs1 ^ rs2) : rs1$	czero.nez rd, rs2, rc xor rd, rs1, rd	
Conditional bitwise-xor, if non-zero $rd = (rc != 0) ? (rs1 ^ rs2) : rs1$	czero.eqz rd, rs2, rc xor rd, rs1, rd	
Conditional bitwise-and, if zero $rd = (rc == 0) ? (rs1 & rs2) : rs1$	and rd, rs1, rs2 czero.eqz rtmp, rs1, rc or rd, rd, rtmp	
Conditional bitwise-and, if non-zero $rd = (rc != 0) ? (rs1 & rs2) : rs1$	and rd, rs1, rs2 czero.nez rtmp, rs1, rc or rd, rd, rtmp	3 insns (requires 1 temporary)
Conditional select, if zero $rd = (rc == 0) ? rs1 : rs2$	czero.nez rd, rs1, rc czero.eqz rtmp, rs2, rc or rd, rd, rtmp	
Conditional select, if non-zero $rd = (rc != 0) ? rs1 : rs2$	czero.eqz rd, rs1, rc czero.nez rtmp, rs2, rc or rd, rd, rtmp	